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KIMath

SUBROUTINES PROGRAMMING MANUAL

**MICROCOMPUTER FAMILY
KIM MATHEMATICS SUBROUTINES
Programming Manual**

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INTRODUCTION

This manual contains instructions on the use of the KIMath subroutine package. KIMath provides the software for doing floating-point (decimal) arithmetic on the KIM microcomputer system.

Because of the wide availability of calculators and larger computers we take decimal arithmetic for granted. Perhaps you were surprised when you first read the 6502 microcomputer programming manual and discovered that the 6502 instruction set could do addition and subtraction of whole numbers only in the range of +128 to -127! If you tried writing a program to multiply two eight-bit binary words to produce a sixteen-bit result, you will appreciate the ability in KIMath to multiply two sixteen-digit BCD numbers just by calling a subroutine.

KIMath gives you the capability to add, subtract, multiply, divide and calculate square roots. In addition, KIMath can calculate logs and antilogs as well as tangents and arc tangents. Other common mathematical functions can be calculated by combining these basic operations. KIMath also has special subroutines to make it easy to evaluate polynomial expressions, which can be used to approximate most mathematical functions.

Since your KIM system, like any microcomputer, has to break large numbers into several successive words in memory and then work on them one word at a time, floating-point arithmetic is much more time- and memory-consuming than simple binary arithmetic. A single floating-point number may use as many as 18 words of memory in your KIM system, and the multiplication of two 16-digit numbers may take as long as 40 milliseconds--40,000 machine cycles! The same multiplication using only 4 digits of precision would take only 10 milliseconds, so if you will accept less precision you can speed up your calculations.

To make this possible KIMath allows you to specify the number of digits to be used in any calculation. You may use the same precision for your entire program or you may adjust the precision depending on the needs of each step in your calculations.

The KIMath subroutines are an ideal addition to the KIM resident assembler. Naturally, the assembler is not required in order to use KIMath.

Because each floating-point number is stored in several memory words, KIMath provide utility routines to move blocks of words corresponding to each number to and from the locations required to perform the calculations.

Finally, KIMath allows floating point values to be stored in several different representations or formats to improve computational efficiency while conserving memory. Routines are provided to convert between these formats.

Chapter 1 of this manual gives you an overview of the concepts used in KIMath. Chapter 2 shows how KIMath could be used to solve a simple equation. Chapter 3 examines the different storage formats and Chapters 4 and 5 cover the use of the individual subroutines.

Several appendices are provided for the user who desires a more rigorous explanation of the techniques used in designing KIMath. A storage map and a complete program listing are also included.

The KIMath subroutines were carefully designed and thoroughly tested to insure accuracy and correct operation. The final decision of their suitability, however, rests with the user and no warranty of fitness is implied by this document. If you suspect an error in these routines, contact the Manager of Product Support at MOS Technology Sales, Inc. 901 California Avenue, Palo Alto, California 94304.

This manual assumes familiarity with 650X assembly language. You should first study the KIM Resident Assembler if necessary.

CHAPTER 1

Overview

1.0 Basic Concepts

KIMath is a package of subroutines designed for use with any MCS650X microprocessor-based system. The subroutines assist the user in doing the floating-point arithmetic necessary for many business and scientific applications. In particular, KIMath is designed to work with the KIM resident assembler, although KIMath may be integrated into any MCS650X program.

1.1 Definitions

Before examining the operation of KIMath, here are definitions of some terms used throughout this manual:

WORD or BYTE - the basic unit of data used in the KIM system. It is composed of eight BITS or binary digits, each of which may only have values of 0 or 1. The microprocessor can only operate on one word at a time.

REGISTER - has a special meaning in this manual, where we define it to mean a group of adjacent words in memory. Since the numbers manipulated by KIMath are too big to store in a single word, we call the group of words used to hold a single floating-point number a register. This usage should not be confused with the registers A, X, Y, stack pointer and status which are single words stored within the microprocessor, not the memory.

POINTER - two adjacent words in memory containing the address of another memory location. The programming convention in the 6502 is to store the two low-order hex characters of the address in the first word and the two high-order characters in the second word. For example, given:

<u>Label</u>	<u>Location</u>	<u>Contents</u>
POINTL	17FA	00
POINTH	17FB	1C

we could say that the register (POINTL, POINTH) is a pointer which currently points to the address 1C00.

ROUTINE - a segment of machine-language code which performs a specified operation. In this document it is synonymous with a subroutine.

ARGUMENT - the KIMath routines expect to find the numbers they are to operate on in fixed locations in memory. When the KIMath routine ADD is called by a program, it will take data from two fixed locations, add them together and put the result in a third fixed location. The data in the input locations are the arguments for the ADD routine. It is usually necessary to transfer the data from other registers in memory into the argument registers and transfer the calculated result to another memory register for storage. KIMath provides routines for moving registers about in memory.

WHOLE NUMBER or INTEGER - a number which does not contain a decimal point.

FLOATING-POINT NUMBER - a number containing a decimal point. Note: 12 is an integer while 12.0 is a floating-point number..

1.2 Number Systems

All modern digital computers are based on the binary system, where numbers are represented as groups of signals which are on or off, 1 or 0. For instance, the decimal number 39 is represented in binary as 100111. For convenience, binary numbers are often converted to groups of four and represented in base 16 or hexadecimal. Decimal 39 is represented as:

2	7	hexadecimal
0010	0111	binary

Because our daily arithmetic is done in decimal, the hexadecimal format is still difficult to interpret. The MCS6502 microprocessor also has the capability to do arithmetic in a modified binary format known as BCD or binary coded decimal. In this system each eight-bit word is divided into two four-bit fields, each representing a decimal digit. Decimal 39 is now represented as:

3	9	decimal
0011	1001	BCD
memory words		

Although less efficient in use of storage, this representation is more readily understood. Appendix H of the 6502 Programming Manual contains a review of binary and BCD arithmetic.

We can represent any whole number as a series of BCD digits, with two digits packed in each memory word. The next problem arises when we wish to store numbers such as 3.14159 or 1276.3333. These numbers contain decimal fractions and BCD notation contains no provisions for decimal points. Very large or very small numbers such as 987000000000 or 0.000000625 often have large numbers of leading or trailing zeros which expand storage requirements. To overcome these problems, KIMath uses "scientific notation" which stores decimal numbers in two parts. The number is first reduced to a single whole number followed by a decimal fraction. This is called the mantissa. The second part is the number of decimal places we must shift the new decimal point to get back to the original number. This is called the exponent (scientific notation is also referred to as exponential notation). If the decimal point must be shifted to the right the exponent is a positive number; if a left shift is necessary the exponent is negative. For instance:

987000000000 becomes 9.87 E+11
0.000000625 becomes 6.25 E-7

(An E for exponent is commonly used to separate the two values.) Similarly:

-6.273411 E+3 is -6273.411

which is an example of a number with a negative mantissa and a positive exponent.

1.3 How Will I Use KIMath in My Program?

Your program will use data of two types: variables whose values will change from one program run to the next or within a single run, and constants which will always have the same value. For example, a simple program to compute the area of a circle given its radius will solve the equation:

$$\text{AREA} = 3.14 \text{ R}^2$$

The values for AREA and R will change each time you run the program; they are variables. The values 3.14 (pi) and 2 (the power to which R is raised) are constants.

When you write your program in assembly language you define memory locations to store this data. When you reserve storage for data used with KIMath, you will have to reserve several words of storage for each variable and constant and you will have to define a register for each data value. The size of each register will depend on the number of digits you wish to store--that is, how much precision you wish to maintain in your calculations.

When you wish to perform calculations using KIMath, you will have to transfer the value of the variables or constants from their memory storage registers into the argument registers (RX and RY) of KIMath, call the appropriate KIMath routine to do the calculation, and then transfer the answer or result from the KIMath result register (RZ) back to a variable storage register. KIMath provides routines for doing these register moves.

In addition, KIMath provides for a variety of data formats to minimize storage requirements and speed computation. KIMath provides routines for converting data between these different formats.

1.4 Summary of KIMath Calculation Subroutines

<u>Subroutine</u>	<u>Action</u>	<u>Description</u>
ADD	RX + RY → RZ	This routine allows the user to add two floating-point numbers. The user may select the number of digits in the mantissas of the arguments.
SUB	RX - RY → RZ	This routine allows the user to subtract two floating-point numbers. The user may select the number of digits in the mantissas of the arguments.
MULPLY	RX * RY → RZ	This routine allows the user to multiply two floating-point numbers. The user may select the number of digits in the mantissas of the arguments.
DIVIDE	RX ÷ RY → RZ	This routine allows the user to divide two floating-point numbers. The user may select the number of digits in the mantissas of the arguments.

1.4 Summary of KIMath Calculation Subroutines (Continued)

<u>Subroutine</u>	<u>Action</u>	<u>Argument Range</u>	<u>Description</u>
SQRT	$\sqrt{RX} \rightarrow RZ$	{1, 100}	This routine allows the user to form the square root of a floating-point number. The user may select the number of digits in the mantissa of the argument.
LOG	$\text{LOG}_{10}(RX) \rightarrow RZ$	$\left\{ \frac{1}{\sqrt{10}}, \sqrt{10} \right\}$	This routine allows the user to form the common logarithm of a floating-point number. The user may select the number of digits in the mantissa of the argument, but at most 14 will be used and at most 8 will be returned.
TENX	$10^{RX} \rightarrow RZ$	{0, 1}	This routine allows the user to form the common antilog of a floating-point number. The user may select the number of digits in the mantissa of the argument, but at most 12 will be used and at most 8 will be returned.
TANX	$\text{TAN}(RX) \rightarrow RZ$	{0, 1}	This routine allows the user to form the tangent of a floating-point number.. The user may select the number of digits in the mantissa of the argument, but at most 14 will be used and at most 8 will be returned. The argument and result are in radians.
ATANX	$\text{ARCTAN}(RX) \rightarrow RZ$	{0, 1}	This routine allows the user to form the arc tangent of a floating-point number. The user may select the number of digits in the mantissa, but at most 14 will be used and at most 8 will be returned. The argument and result are in radians.

1.5 KIMath Formats

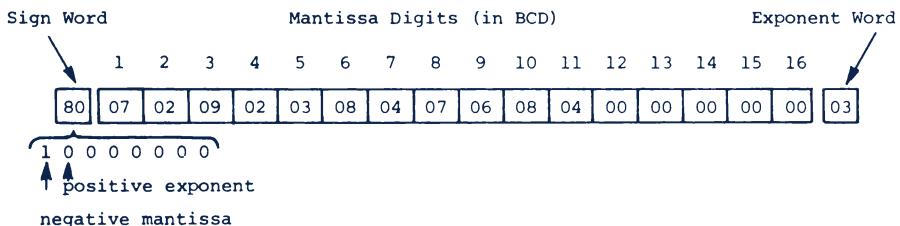
Several variations on exponential notation are used in KIMath.

1.5.1 Computational Format

Before actual computation can take place, every number transferred to a KIMath argument register must be converted to an eighteen-word format:

- a. The first word contains the sign of the mantissa in bit 7 (1 if minus, 0 if plus) and the sign of the exponent in bit 6.

- b. The next sixteen words contain the sixteen digits of the mantissa, one digit per word, using BCD notation.
- c. The last word contains the value of the exponent (which must be between 0 and 99) as two BCD numbers. For example: a value of -7292.3847684 would be stored in computational format as:



This format allows KIMath to calculate quickly since every floating-point number is contained in the same number of words and each word contains only a single digit (except for the exponent word).

This is a very inefficient use of storage since the left half of each word in the mantissa always contains a zero, and each mantissa contains sixteen digits even if many are just trailing zeros.

1.5.2 Packed Format

To avoid this waste of storage KIMath provides a more efficient format for storing data when it is not being used in a computation:

- a. The first word contains the signs of the mantissa and exponent as in computational format.
- b. Succeeding words each contain two digits of the mantissa in BCD. The number of words used is defined by the user.
- c. The final word contains the two digits of the exponent, as in computational format.

If the user had decided to use 12 digits of precision in storage, the value of -7292.3847684 would be stored as:



thus using eight words of storage instead of eighteen. KIMath contains routines to convert between computational and packed format. Packed format is the form

usually used for memory registers storing program variables. It is your responsibility to keep track of how many digits of precision are used for each variable, since KIMath will need that information when moving the memory register to the KIMath argument registers and converting to computational format. You may choose to use the same precision for all storage registers to simplify this task.

1.5.3 Unpacked (ASCII) Format

The ASCII code used by the KIM-1 serial communications interface represents each character typed in or out as two hexadecimal characters packed in a single word: the character 'A' is 41, a blank is 20, a carriage return is 0D. The digits 0 through 9 are coded as themselves plus 30: '4' is 34, '7' is 37, etc.

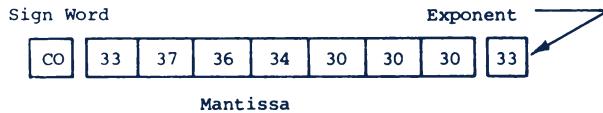
To make interfacing easier, KIMath has an unpacked format which uses ASCII codes rather than BCD.

- a. The first word contains the sign bits for the mantissa and exponent as used for computational format.
- b. Up to 16 words (depending on the precision desired) are used for the mantissa. Each word contains one digit of the mantissa in ASCII code.
- c. The last two words contain the two digits of the exponent, also in ASCII.

Using six digits of precision the value

-0.003764 (or -3.764 E-3)

would be stored as:



This format is usually used only when preparing to transfer data to or from a line of terminal input or output.

1.5.4 Constant Format

KIMath calculates functions by solving polynomial equations which closely approximate the desired function. Because the polynomials have many terms and the coefficients have differing precision requirements, a special format is used in KIMath to store constants. The format allows all constants to

be stored together in a list without having to specify elsewhere the precision of each constant. The constant format is composed of:

- a. A sign byte, as in the other formats.
- b. Up to eight bytes in packed BCD format (two BCD characters per byte) containing the mantissa.
- c. An exponent byte of two characters. The first (leftmost) character is always a hexadecimal 'F' and the second character is the one-digit BCD exponent. Thus, constants may only have exponents between -9 and +9.

Thus the constant 32767.413 (3.2761413 E+4) will be stored as:



Thus, KIMath can read a constant from memory by just knowing its starting location. It will continue to scan successive bytes until it locates one containing an 'F' and know it has reached the end of the constant. The next byte will contain the sign byte of the next constant.

The KIMath routine USRLKP is used to transfer a constant format value from memory to the argument register RY and convert it to computational format.

1.6 Summary of Formats

KIMath supports four different formats:

- a. Computational Format - used for the argument and result registers RX, RY and RE. Although inefficient in use of storage, they permit quick calculation.
- b. Packed BCD Format - most efficient in use of storage.
- c. Unpacked (ASCII) Format - less efficient in use of storage, but easier to convert to and from external devices.
- d. Constant Format - efficient in use of storage and a good way to store tables, but limited in the size of the exponents which can be used.

CHAPTER 2
A Sample Program Using KIMath

2.0 Throughout the remainder of this manual we will develop the techniques required to create the following program:

An engineer wishes to create a program to find the resonant frequency of a tuned circuit, given values for L, the inductance, and C, the capacitance, in the circuit. The formula is:

$$f = \frac{1}{2\pi \sqrt{LC}}$$

The program will have to contain the following operations:

1. The starting address of the program must be defined and storage registers for L, C, and the constants 2π and 1 must be created.
2. The precision of the calculations must be defined to KIMath.
3. The value of L must be read in and transferred to a KIMath argument register.
4. The value of C must be read in and transferred to a KIMath argument register.
5. L must be multiplied by C.
6. Move the answer to step 5 to a KIMath argument register. Compute the square root. Since the square root routine in KIMath will only accept numbers between 1 and 100 we must adjust the exponent of the result of step 5.
7. Get the first constant (2π) and transfer it to an argument register.
8. Move the computed square root to another argument register and multiply.
9. Get the second constant (1) and put in an argument register.
10. Move the result of step 8 to an argument register and do the division.
11. Move the result of the calculation back to the register originally containing L and type it out.

To summarize as an assembly language program,

```
; (1) define storage  
;(2) define precision
```

```
; (3)  read L and transfer to RX
; (4)  read C and transfer to RY
; (5)  compute L times C--answer stored in RB
; (6)  move RB to RX and
;       compute square root--adjust exponent of result
; (7)  get  $2\pi$  to RX
; (8)  move RB to RY and multiply
; (9)  get 1 (constant) and put in RX
; (10) move RB to RY and divide
; (11) move RB back to L and print out
.END
```

To start with, our program only contains comments. As we learn more about KIMath we will begin to write the assembly language statements to implement the program.

To help you visualize what is happening in the program, the following table shows the contents of the two argument registers (RX, RY) and the result register (RB) after each program step:

At the end of step #	RX	RY	RB
1	--	--	--
2	--	--	--
3	L	--	--
4	L	C	--
5	L	C	LC (L times C)
6	LC	C	\sqrt{LC} (adjusted)
7	2π	$2\pi^*$	\sqrt{LC}
8	2π	\sqrt{LC}	$2\pi\sqrt{LC}$
9	1	1^*	$2\pi\sqrt{LC}$
10	1	$2\pi\sqrt{LC}$	$\frac{1}{2\pi\sqrt{LC}}$

*The constant appears in both RX and RY because it is moved from storage to RY and then to RX.

CHAPTER 3

Defining Storage for KIMath Variables and Constants

3.0 In the last chapter you learned that KIMath uses several different formats for memory storage registers used in KIMath calculation. Now let us see how we set aside memory to hold these variables and constants.

3.1 Defining Precision

If we define an even number NDIG which is the Number of DIGits of precision we wish to use, storing a number of that precision will take:

- a. NDIG + 3 words of storage in Unpacked format.
- b. NDIG/2 + 2 words of storage in Packed format.
- c. NDIG/2 + 2 words of storage in Constant format. (Constants may be stored in fewer digits if they have leading or trailing zeros.)
- d. Computational format always takes 18 words.

In our sample program (see Chapter 2) we defined two variables, L and C. Since we are reading them in from a terminal we will store them in Unpacked (ASCII) format. We also needed two constants (2π and 1) and will store them in constant format.

Since we may wish to change the precision of our program, we will define the variable NDIG to the assembler and use it to calculate our storage. The following code will define storage for our sample program:

```
; (1)  define storage
NDIG = 8; we will carry 8 digits of precision
* = $5000; start the program at location 5000 (hex)
L  * = * + NDIG + 3; reserve storage for L
C  * = * +NDIG +3; reserve storage for C
TWOPI .BYTE $00, $62, $83, $18, $54, $F0; define 2π
ONE   .BYTE $00, $10, $F0; define 1.0
```

Note that constant format allows us to eliminate the trailing zeros in the constant ONE. The symbol "*" stands for the current value of the program counter. Refer to the assembler documentation if you are not familiar with assembler directives.

Our program storage map will now be:

5000 - 500A	storage for L
500B - 5015	storage for C
5016 - 501B	storage for 2π constant
501C - 501E	storage for 1.0 constant
501E -	program storage

If we later re-assemble our program and change the value of NDIG to 4 or 10, to increase or decrease the precision of our calculations, the amount of storage allocated will automatically decrease or increase.

CHAPTER 4

Introduction to the KIMath Subroutines

4.0 Now that we have examined the data formats and storage requirements of KIMath we are ready to examine the subroutines themselves. The following table summarizes the KIMath subroutines, their arguments and their functions:

Summary of KIMath Subroutines

Subroutine	Arguments	Function
SAVXY	None	Save processor X and Y index registers.
RCLXY	None	Recall previously saved X and Y registers.
IPREC	PREC, EXTRA	PREC + EXTRA → N
PLOADX	ARGXL, ARGXH	Move user register in packed format to RX or RY and change to computational format.
PLOADY	ARGYL, ARGYH	
	PREC	
ULOADX	ARGXL, ARGXH	Move user register in unpacked format to RX or RY and change to computational format.
ULOADY	ARGYL, ARGYH	
	PREC	
PSTRES	RESL, RESH	Move contents of RB to user register in packed or unpacked format.
USTRES	PREC	
USRLKP	KONL, KONH, NKON	Transfer constant data to RY and convert to computational format.
POLY	KONL, KONH, NKON, DEG	Evaluates a polynomial.
CLRX	None	Clear RX, RY, or RB
CLRY		
CLRB		
DECHEX	CNT	Converts contents of CNT from decimal to hexadecimal notation.
XSY	RX, RY	Exchange contents of RX and RY.

4.1 Working Storage

The working storage area used by the KIMath subroutines is divided into three parts: the computational registers (RA, RB, RQ, RM and RN) which are not available to the user; the argument/result registers (RX, RY, RB); and the pointers, counters and indicators in page zero.

4.1.1 The Argument/Result Registers - RX, RY, RB

The registers RX and RY are the argument registers and RB is the result register. This means that when RX and RY are set to specific values and, for example, the MLTPLY routine is called, then the product is returned in RB. These registers are each 18 bytes long to permit a maximum of 16 digits for the mantissa. The registers RX, RY and RB are CLeaRed directly by means of the subroutines CLRX, CLRY and CLR~~B~~, respectively. The registers RX, RY, RB contain a sign byte (SX, SY, SB) and an exponent byte (EX, EY, EB). The sign byte is the first and the exponent byte the last. Remember that data in RX, RY or RB must be in computational format.

These registers may be transferred about by means of the subroutines whose names are of the form MVsd, where "s" stands for source and "d" for destination. For example, MVXY MoVes the contents of RX to RY.

4.1.2 The Pointers, Counters and Indicators

Before calling any KIMath routine, you may wish to SAVe the processor's X and Y registers in locations TMPX and TMPY by means of the routine SAVXY. The X and Y registers are ReCaLled by calling RCLXY.

Before calling an arithmetic or transcendental KIMath routine you must specify the number of digits of PRECision desired by loading locations PREC and EXTRA with two unsigned, binary values. The first value, PREC, specifies the number of digits of mantissa which the user expects in the RB register after completion of a calculation and PREC+EXTRA specifies the number of digits of precision used by the subroutine in computing the result. The limitations on these values are: $0 < \text{PREC} + \text{EXTRA} \leq 16$ and:

<u>Function</u>	<u>Internal Precision*</u>	<u>Precision* Returned</u>
ADD	PREC+EXTRA	PREC
SUB	PREC+EXTRA	PREC
MLTPLY	PREC+EXTRA	PREC
DIVIDE	PREC+EXTRA	PREC.
LOG	14	Lesser of PREC or 8
TENX	12	Lesser of PREC or 8
TANX	14	Lesser of PREC or 8
ATANX	14	Lesser of PREC or 8
SQRT	PREC+EXTRA	PREC

*Note: Here the word "precision" refers not to error but to the number of digits used in calculation.

The sum PREC+EXTRA may be computed by calling the routine named IPREC. This routine computes the sum and stores the result in N. N is the byte which the arithmetic and transcendental routines use to decide the precision of internal calculation. If the user so desires, the use of PREC and EXTRA may be avoided and N used directly. However, the routines LOG, TENX, TANX and ATANX may change any value placed in N.

The bytes in page Ø named ARGXL, ARGXH, ARGYL, and ARGYH are used to contain the addresses of user-defined registers in the RAM memory area. By using these address pointers and the subroutines PLOADX, PLOADY, ULOADX and ULOADY, the user loads the KIMath argument registers from any place in memory.

The bytes RESL and RESH are used in conjunction with the KIMath routines PSTRES (Packed SToring of RESults) and USTRES (Unpacked SToring of RESults) to place the contents of R₈ in a user-defined register, which is specified by the pointer (RESL, RESH).

The bytes KONL, KONH are used together with DEG and KIMath routines USRLKP and POLY to evaluate any polynomial, subject, of course, to the constraints that the coefficients have fewer than 16 digits of mantissa and that the exponent has only one digit. More shall be said about these techniques in the section on Advanced Applications.

CHAPTER 5
The KIMath Subroutines

5.1 The KIMath subroutines may be divided into four groups, Functions, Conversions, User Utilities and System Utilities. This section will deal with the first three classes.

5.1.1 The Function Subroutines are called ADD, SUB, MLTPLY, DIVIDE, SQRT, LOG, TENX, TANX, and ATANX. These routines may be divided into the Arithmetic Routines and the Transcendental Routines. Each of the routines requires that the argument registers (RX and RY) be in the Computational Format.

The following tables give some relevant information about the Arithmetic and Transcendental Routines:

<u>Arithmetic Routines</u>			
Subroutine	Precision ⁽¹⁾	Error ⁽²⁾	Symbolic Action
ADD	$0 < N \leq 16$	1 Count	$RZ \leftarrow RX + RY$ ⁽³⁾
SUB	$0 < N \leq 16$	1 Count	$RZ \leftarrow RX - RY$ ⁽³⁾
MLTPLY	$0 < N \leq 16$	1 Count	$RZ \leftarrow RX * RY$
DIVIDE	$0 < N \leq 16$	1 Count	$RZ \leftarrow RX \div RY$
SQRT	$0 < N \leq 16$	5 Counts	$RZ \leftarrow \sqrt{RX}$

(1) Note - Precision refers to the number of digits to which the calculation is to be carried out.

(2) Note - Error refers to the error in the least significant digit of the mantissa and so is the relative error.

(3) Note - In the case of ADD, at return time RX contains the argument with the largest absolute value. Thus one may use this routine to sort a table of numbers. In the case of SUB the sign of the mantissa of RY is changed and then the arguments are added. Thus the arguments are intact except for a sign.

Transcendental Routines

Subroutines	Interval	Precision	Error	Symbolic Action
LOG	$\left\{ \frac{1}{\sqrt{10}}, \sqrt{10} \right\}$	14	$<10^{-8}* \quad$	$RZ \leftarrow \text{LOG}_{10}(RX)$
TENX	{0, 1}	12	$10^{-8} \quad$	$RZ \leftarrow 10^{RX}$
TANX	{0, 1}	14	$10^{-8} \quad$	$RZ \leftarrow \text{TAN}(RX)$
ATANX	{0, 1}	14	$10^{-8} \quad$	$RZ \leftarrow \text{ARCTAN}(RX)$

*Note: The error in the case of LOG is the absolute error; the others are relative error.

The user may use the bytes PREC and EXTRA together with the subroutine IPREC to calculate N, which is the value used by these routines. Note that the Transcendental Routines supply their own value for N and will overwrite any value supplied by the user.

Now that we know the names of the calculation routines and how to indicate to KIMath the desired precision of calculation we can add some statements to our sample program. Now it will look like (the > sign indicates new statements we have added) :

```
; (1) define storage
NDIG = 8
* = $5000
L * = *+NDIG+3
C * = *+NDIG+3
TWOP1    •BYTE $00, $62, $83, $18, $54, $F0
ONE      •BYTE $00, $10, $F0
; (2) define precision
>      EX = 2; use 2 extra digits in calculations
>      LDA # NDIG
>      STA PREC
>      LDA # EX
>      STA EXTRA
>      JSR IPREC ; calculate N
; (3) read L and transfer to RX
```

```
; (4) read C and transfer to RY
; (5) compute L times C - answer stored in R8
>      JSR MLTPLY ; do the multiplication
; (6) move R8 to RX and compute square root - adj. exp.
>      JSR SQRT
; (7) get 2π to RX
; (8) move R8 to RY and multiply
>      JSR MLTPLY
; (9) get 1 (constant) and put in RX
; (10) move R8 to RY and divide
>      JSR DIVIDE
; (11) move R8 to L and print out
      END
```

Now by simply changing the assigned value for NDIG and EX and re-assembling the program, both calculated precision and storage allocation can be automatically changed. This allows re-assembly of the program with changed precision (and thus changed locations for registers L and C). The pointers will automatically be changed to the correct locations.

5.1.2 The Conversion Routines permit the user to transform data from the packed BCD, constant, or unpacked formats to the computational format and back again. They are PLOADX, PLOADY, ULOADX, 'LOADY, PSTRES, and USTRES. These routines all use PREC to determine the number of bytes to be converted.

- a) PLOADX, PLOADY: These routines load RX and RY with the packed format data found at the addresses {ARGYL, ARGXH} and {ARGYL, ARGYH}. The data is unpacked into the computational format and the value in PREC determines how many digits will be transferred. Note that PREC should be an even number to prevent the loss of a digit in transfer. This is natural when it is noted that the packed format requires an even number of digits in the mantissa bytes.
- b) ULOADX, ULOADY: These routines load RX and RY with the unpacked format data found at the addresses {ARGXL, ARGYL} and {ARGYL, ARGYH} respectively. The data is converted into computational format and the number of digits transferred is determined by PREC. In this case PREC need not be even to effect an accurate transfer.

The use of ULOAD and PLOAD allows us to move data to the computational registers from memory registers. We can now add them to our program at the following locations:

```
; (3) Read L and transfer to RX
> JSR GETNL ; a user-provided subroutine to get a
> ; number from the terminal and place
> ; it in L in unpacked format.
> LDA #<L ; low-order byte of address of L
> STA ARGXL
> LDA #>L ; high-order byte
> STA ARGXH
> JSR ULOADX ; move L to RX
; (4) read C and transfer to RY
> LDA #<C ; low-order byte of address of C
> STA ARGYL
> LDA #>C ; high-order byte of address of L
> STA ARGYH
> JSR GETNC ; user-supplied routine to read value of C into
register C
> JSR ULOADY
```

Note: The assembler interprets the symbols "<L" to mean the low-order byte of the address of the variable named L. ">L" implies the high-order byte of the address of L.

- c) PSTRES: This routine moves the contents of RB into the location specified by {RESL, RESH} as a starting address and converts it to packed format. The byte PREC is used to determine the number of bytes to be converted.
- d) USTRES: This routine converts the contents of RB from computational format as it moves them to the destination specified by {RESL, RESH} and converts the contents to unpacked format. PREC, again, determines the number of bytes to be moved.

These routines allow transfer of the final answer to our example problem back to register L for printout:

```
; (11) move RZ to L and printout
>      LDA <L; set up
>      STA RESL; pointer
>      LDA >L; to locate
>      STA RESH; destination register
>      JSR USTRES; move it
>      JSR PRINTL; user-supplied routine to print the value
      in L
```

- e) USRLKP: This routine loads RY with the constant format data beginning at the address {KONL, KОН}. The data is converted from constant format to computational format. The number of digits transferred is determined by the location of the byte containing an 'F.' This routine also uses the byte named NKON to point at the constant. Thus if the user sets NKON equal to zero and the pair {KONL, KОН} equal to some address prior to calling USRLKP then at return time RY contains the constant. {KONL, KОН, and NKON} point at the next constant, that is, {KONL, KОН} + NKON is the address of the byte following the byte containing an F (presumably the first byte of the next constant). Using this routine and POLY one may step through a table of coefficients and evaluate a polynomial. Use of POLY is covered in Appendix A. Clearly, one may also use this routine for other purposes such as a simple table lookup of constants needed frequently in a calculation.

Now we have the capability to load our constants into the program. Now we can fill out steps (7) and (9):

```
; (7) get 2π to RX
>      LDA <TWOPI; get low-order
>      STA KONL ; pointer byte
>      LDA >TWOPI; get high-order
>      STA KОН ; pointer byte
>      LDA #00
>      STA NKON ; initialize offset
>      JSR USRLKP ; lookup constant and put in RY
      ; still have to get RY to RX
```

```
> ; (9) get 1 (constant) and put in RX
> JSR USRLKP ; a lot easier the second time
> ; still must transfer RX from RY
```

Up to 256 bytes of constants may be indexed through by successive calls to USRLKP.

- f) DECHECK: This routine converts the contents of the byte CNT from BCD to HEX. The converted value may be found at return time in CNT.

The User Utilities permit the user to move and clear various registers.

- a) The routines whose names are of the form MV s, d, where "s" stands for source and "d" destination, allow the user to move the registers RX, RY, R_Z, RM and RN about with ease. The following table shows the supported moves. RM and RN are working registers and will not normally be available to the user.

Source	Destinations	Corresponding Routine Names
RX	RY, R _Z , RM, RN	MVXY, MVXB, MVXM, MVXN
RY	RX, R _Z , RM, RN	MVYX, MVYB, MVYM, MVYN
R _Z	RX, RY, RM, RN	MVBX, MVBY, MVBM, MVBN
RM	RX, RY, R _Z , RN	MVMX, MVMY, MVMB, MVMN
RN	RX, RY, R _Z , RM	MVNX, MVNY, MVNB, MVNM

- b) The routines CLRX, CLRY and CLR_Z may be used to set RX, RY and R_Z equal to zero, respectively.
- c) The routine XSY exchanges RX and RY.

The move subroutines allow us to transfer computational data between working registers. We need them in steps (6), (7), (8), (9), and (10):

```
> ; (6) Move RZ to RX and compute square root - adj. exp.
> JSR MVBX; move RZ to RX
JSR SQRT
; (7) get 2π to RX
LDA #<TWOPI
```

```
        STA KONL
        LDA #>TWOP1
        STA KONH
        LDA #00
        STA NKON
        JSR USRLKP ; constant to RY
>      JSR MVYX ; move it to RX
; (8) move RZ to RY and multiply
>      JSR MVZY
        JSR MLTPLY
; (9) get 1 (constant) and put in RX
        JSR USRLKP ; get next constant
>      JSR MVYX ; move it to RX
; (10) move RZ to RY and divide
>      JSR MVZY ; move RZ to RY
        JSR DIVIDE
```

The sample program is now complete except for the fact that SQRT will only operate on numbers between 1 and 100, so we must adjust the exponent of the argument, take the square root and adjust the results. The two subroutines required are listed in Appendix B and are called SQRIN & SQROUT. They must be added to the program (they are not included in KIMath). Step (6) now becomes:

```
; (6) Move RZ to RX and compute square root
        JSR MVZX ; move RZ to RX
>      JSR SQRIN ; adjust exponent
        JSR SQRT ; compute root
>      JSR SQROUT ; adjust exponent back
```

5.2 Completed KIMath Program

```
; (1) define storage
NDIG = 8
* = $5000
L * = *+NDIG+3
C * = *+NDIG+3
TWOP1 •BYTE #00, $62, $83, $18, $54, $F0; TWOP1
ONE    •BYTE $00, $10, $F0
; (2) define precision
```

```
EX = 2; use 2 extra digits in calculation
LDA #NDIG
STA PREC
LDA #EX
STA EXTRA
JSR PREC ; calculate N
; (3) Read L and transfer to RX
JSR GETNL; user-provided subroutine to input value for L
           ; to L register in unpacked format.
LDA #>L; low-order byte of address of L
STA ARGXL
LDA #<L ; high-order byte
STA ARGXH
JSR ULOADX ; Move L to RX
; (4) Read C and transfer to RY
LDA #<C ; low-order byte of address of C
STA ARGYL
LDA #>C ; high-order byte of address of C
STA ARGYH
JSR GETNC ; user-provided subroutine to input value
           ; for C to C register in unpacked format.
JSR ULOADY ; transfer C to RY
; (5) Compute L times C
JSR MLTPLY
; (6) Move RZ to RX and compute square root
JSR MVEX ; move RZ to RX
JSR SQRIN ; adjust exponent
JSR SQRT ; compute root
JSR SQROUT ; adjust exponent back
; (7) get 2pi to RX
LDA #<TWOPI ; address low of First constant
STA KONL
LDA #>TWOPI
STA KONH
LDA #00
```

```
STA NKON
JSR USRLKP ; constant to RY
JSR MVYX; move it to RX
; (8) Move RB to RY and multiply
JSR MVBY
JSR MLTPLY
; (9) get 1 (constant) and put in RX
JSR USRLKP ; get next constant
JSR MVYX ; move it to RX
; (10) Move RB to RY and divide
JSR MVBY ; move RB to RY
JSR DIVIDE ;
; (11) Move RB to L and print it out
LDA #<L ; set up
STA RES ; pointer
LDA #>L ; to locate
STA RES+1 ; destination register
JSR USTRES ; move it
JSR PRINTL ; user-supplied routine to print out L.
.END
```

Note: The above program would also require that the starting address of all KIMath subroutines be defined. See Appendix D for the correct addresses.

Appendix A
EVALUATING POLYNOMIALS

KIMath uses polynomial approximations to generate several of its functions, and these subroutines are also available for calculation of user-defined polynomials.

The user must first define the constants to be used in the calculation, then define the degree of the polynomial and the starting address of the stored constant. Finally, the argument for the polynomial is transferred to RX and the routine POLY is called. The value of the polynomial is returned in RB.

Specifically, given a polynomial of the form:

$$y = c_0 + c_1x + c_2x^2 + c_3x^3 + \cdots + c_Nx^N$$

the coefficients ($c_0, c_1, c_2, c_3, \dots, c_N$) are stored in constant format in contiguous memory. The highest order coefficient is stored in the lowest memory address. NKON is set to zero, DEG is set to a value of $N - 1$ where N is the degree of the polynomial. The address pair (KONL, KONH) must point to the first byte of the constant storage area. RX is loaded with the value of x and the subroutine POLY is called. Upon return, y (the value of the polynomial) is in RB.

A Sample Program

The sine function can be expressed as:

$$\sum_{N=1}^{\infty} (-1)^{N+1} \frac{z^{2N-1}}{(2N-1)!}$$

Expanded, this yields the polynomial:

$$\sin z = z - \frac{z^3}{3!} + \frac{z^5}{5!} - \frac{z^7}{7!} + \frac{z^9}{9!} - \frac{z^{11}}{11!} \dots$$

For our example we will use the first six non-zero terms for our application. Our coefficients are:

$$c_0 = 0$$

$$c_1 = 1.0$$

$$c_2 = 0$$

$$\begin{aligned}c_3 &= -\frac{1}{3!} = -1.6666667 \text{ E-1} \\c_4 &= 0 \\c_5 &= +\frac{1}{5!} = +8.3333333 \text{ E-3} \\c_6 &= 0 \\c_7 &= -\frac{1}{7!} = -1.9841270 \text{ E-4} \\c_8 &= 0 \\c_9 &= +\frac{1}{9!} = 2.7557319 \text{ E-6} \\c_{10} &= 0 \\c_{11} &= -\frac{1}{11!} = -2.5052108 \text{ E-8}\end{aligned}$$

A program to evaluate the sine function would look like:

```
; define parameter and pointers
PREC * = $10
        .BYTE 04 ; value for PREC
EXTRA * = $11
        .BYTE 04 ; value for EXTRA
KON * = $0E
        .WORD $3000
DEG * = $05
        .BYTE 10
NCON * = $01
        .BYTE $00
* = $3000 ; starting address
NIN = * + 7 ; reserve 7 bytes for input register
; define constants C0 - C11 - note that they are stored in
; reverse order
        .BYTE $C0,$25,$05,$21,$08,$F8 ; C11
        .BYTE 0,0,$F0;C10
        .BYTE $40,$27,$55,$73,$19,$F6 ; C9
        .BYTE 0,0,$F0 ; C8
        .BYTE $C0,$19,$84,$12,$70,$F4 ; C7
        .BYTE 0,0,$F0 ; C6
        .BYTE 04,$B3,$33,$33,$33,$F3 ; C5
        .BYTE 0,0,$F0 ; C4
```

```
.BYTE $C0,$16,$66,$66,$67,$F1 ; C3
.BYTE 0,0$FO ; C2
.BYTE 0,01,$FO ; C1
CONST .BYTE 0,0,$FO ; CØ
; call user-written subroutine to load some
; value into NIN register.
JSR GETVAL

LDA KON ; set up
STA ARGXL ; pointer to
LDA KON+1 ; input
STA ARGXH ; buffer

JSR ULOADX ; move NIN to RX and convert
JSR POLY ; evaluate polynomial
LDA KON ; set up
STA RES ; pointer to
LDA KON+1 ; output
STA RES+1 ; results
JSR USTRES ; move RZ to NIN and convert

; call user-written subroutine to printout
; value in NIN
JSR PUTVAL
.END ; all done
```

Appendix B
APPLICATIONS

In this section we shall deal with extension of the range of the functions, and extension of the function set.

- a) Common Logarithm: Let $x \cdot 10^r$ denote an arbitrary positive floating-point number. Then

$$\begin{aligned}x \cdot 10^r &= x \cdot 10^{-1/2} \cdot 10^{1/2} \cdot 10^r \\&= (x \cdot 10^{-1/2}) \cdot 10^{r+1/2}\end{aligned}$$

Now $1 \leq x < 10$ and so $10^{-1/2} \leq x \cdot 10^{-1/2} < 10^{1/2}$. But $\sqrt{0.1} = 10^{-1/2}$ and so LOG may be used to evaluate $\log_{10}(x \cdot 10^{-1/2})$. By using the following equation, one may compute the common log for any positive argument.

$$\begin{aligned}\log_{10}(x \cdot 10^r) &= \log_{10}(x \cdot 10^{-1/2} \cdot 10^{r+1/2}) \\&= \log_{10}(x \cdot 10^{-1/2}) + r + 0.5\end{aligned}$$

- b) Common Antilog: Let x be an arbitrary floating-point number satisfying the inequality

$$|x| < 100,$$

and let $I = [x]$ and $F = <x>$, where $[x]$ stands for the largest integer less than or equal to x and $<x> = x - [x]$. With these definitions in mind we find that

$$x = I + F$$

Now $0 \leq F < 1$ and I is an integer and so

$$10^x = 10^F \cdot 10^I$$

Since $0 \leq F < 1$ we may use TENX to find 10^F , after which we have a number between 1 and 10. Therefore I is the floating point exponent of 10^x .

- c) Tangent: Let x be a floating-point number which satisfies the inequality $0 \leq x < \frac{\pi}{2}$. This implies that $0 \leq x \cdot \frac{2}{\pi} < 1$. Now $\tan\left(\frac{\pi}{4} \cdot \frac{2}{\pi} x\right) = \tan\left(\frac{x}{2}\right)$ but

$$\tan x = \frac{2 \tan\left(\frac{x}{2}\right)}{1 - \tan^2\left(\frac{x}{2}\right)}$$

Thus by applying TANX to the value of $\frac{2}{\pi} x$ and applying the above trigonometric identity one may extend the tangent function to the interval $[0, \frac{\pi}{2}]$.

Note also that

$$\cos x = \frac{1 - \tan^2\left(\frac{x}{2}\right)}{1 + \tan^2\left(\frac{x}{2}\right)} \quad \text{and}$$

$$\sin x = \frac{2 \tan\left(\frac{x}{2}\right)}{1 + \tan^2\left(\frac{x}{2}\right)}$$

and so one may use TANX to evaluate the other trig functions.

- d) Arc tangent: This function may be extended from the interval $[0, 1]$ by means of the following identities:

$$\arctan(-x) = -\arctan x$$

$$\arctan(|x|) = \frac{\pi}{2} - \arctan\left(\frac{1}{|x|}\right), \quad x \neq 0$$

With these identities and the ATANX subroutine one can evaluate the arc tangent function for any argument. The other inverse trigonometric functions may then be computed by means of the following identities:

$$\arcsin x = \arctan\left(\frac{x}{(1 - x^2)^{1/2}}\right)$$

$$\arccos x = \arctan\frac{(1 - x^2)^{1/2}}{x}$$

- e) Square Root: Since every floating-point number can be expressed as $x \cdot 10^r$ where $1 \leq |x| < 100$ and $|r|$ is even or zero, one can form the square root as follows:

$$\sqrt{x \cdot 10^r} = 10^{r/2} \cdot \sqrt{x}$$

where $x \geq 0$ and \sqrt{x} is evaluated by using SQRT. Note that $r/2$ is an integer and $1 \leq \sqrt{x} < 10$ and so $r/2$ is the floating-point exponent of $\sqrt{x \cdot 10^r}$.

Chapter 5 mentions that the sample program must be provided with two subroutines to extend the range of the square root function. These subroutines are provided below. The subroutine SQRIN must be called immediately prior to calling SQRT and SROUT must be called immediately afterwards. SQRIN and SROUT are not part of KIMath and must be supplied by the user:

Exponent Adjustment for SQRT

	VAL * = *+1
SQRIN	SED
	SEC
	LDA EX
	SBC #2
	STA EX
	BCC OUT
	CLC
	LDA VAL
	ADC #1
	STA VAL
	BNE SQ1
OUT	ADC #2
	BIT SX
	BVC SQ2
	STA EX
	CLC
	ADC VAL
	STA VAL
	RTS

SQ2 STA EX
RTS

SQROUT LDA VAL
STA EZ
RTS

Appendix C
THE APPROXIMATIONS

This section is included for those users who are interested in how KIMath generates its functions. Understanding of this section is not required for use of KIMath.

- 1) LOG: This subroutine has been built around the rational function given by the following equation:

$$R(x) = t(x) p(t(x)),$$

where $t(x) = \frac{x - 1}{x + 1}$, $x \in [\sqrt{0.1}, \sqrt{10}]$ and $p(x)$ is a polynomial given by

$$p(x) = a_0 + a_1x + a_2x^2 + a_3x^3 + a_4x^4 + a_5x^5$$

where

$$\begin{aligned} a_0 &= 8.685887483405 \times 10^{-1} \\ a_1 &= 2.89551130267 \times 10^{-1} \\ a_2 &= 1.731095517 \times 10^{-1} \\ a_3 &= 1.3136901121 \times 10^{-1} \\ a_4 &= 5.53427387 \times 10^{-2} \\ a_5 &= 1.820912997 \times 10^{-1} \end{aligned}$$

The absolute error on the interval $[\sqrt{0.1}, \sqrt{10}]$ for $\log_{10}(x)$ as approximated by $R(x)$ is less than 1×10^{-8} .

- 2) TENX: This subroutine has been built around the polynomial given by the following equation:

$$P(x) = (a_0 + a_1x + a_2x^2 + a_3x^3 + a_4x^4 + a_5x^5 + a_6x^6 + a_7x^7)^2,$$

where $x \in [0, 1]$ and

$$\begin{aligned} a_0 &= 1 \\ a_1 &= 1.15129277603 \\ a_2 &= 6.6273088429 \times 10^{-1} \\ a_3 &= 2.5439357484 \times 10^{-1} \end{aligned}$$

$$\begin{aligned}a_4 &= 7.295173666 \times 10^{-2} \\a_5 &= 1.742111988 \times 10^{-2} \\a_6 &= 2.55491796 \times 10^{-3} \\a_7 &= 9.3264267 \times 10^{-4}\end{aligned}$$

The relative error on the interval $[0, 1]$ for 10^x as approximated by $P(x)$ is less than 5×10^{-9} .

- 3) TANX: This subroutine has been built around the polynomial given by

$$P(x) = x(a_0 + a_1(x^2)^1 + a_2(x^2)^2 + a_3(x^2)^3 + a_4(x^2)^4 + a_5(x^2)^5 + a_6(x^2)^6),$$

where

$$\begin{aligned}a_0 &= 7.853981762291 \times 10^{-1} \\a_1 &= 1.614897776174 \times 10^{-1} \\a_2 &= 3.98659104705 \times 10^{-2} \\a_3 &= 9.8345945393 \times 10^{-3} \\a_4 &= 2.7974335037 \times 10^{-3} \\a_5 &= 2.031171084 \times 10^{-4} \\a_6 &= 4.109741948 \times 10^{-4}\end{aligned}$$

The relative error on the interval $[0, 1]$ for $\tan\left(\frac{\pi}{4}x\right)$ as approximated by $P(x)$ is less than 1×10^{-8} .

- 4) ATANX: This subroutine has been built around the polynomial given by

$$P(x) = x(a_0 + a_1(x^2)^1 + a_2(x^2)^2 + a_3(x^2)^3 + a_4(x^2)^4 + a_5(x^2)^5 + a_6(x^2)^6 + a_7(x^2)^7 + a_8(x^2)^8),$$

where

$$\begin{aligned}a_0 &= 9.999999847657 \times 10^{-1} \\a_1 &= -3.333307334505 \times 10^{-1} \\a_2 &= 1.999261939166 \times 10^{-1} \\a_3 &= -1.420364446652 \times 10^{-1} \\a_4 &= 1.06409340253 \times 10^{-1} \\a_5 &= -7.50429453889 \times 10^{-2} \\a_6 &= 4.26915192711 \times 10^{-2} \\a_7 &= -1.60686289604 \times 10^{-2} \\a_8 &= 2.8498896208 \times 10^{-3}\end{aligned}$$

The relative error on the interval $[0, 1]$ for $\arctan x$ as approximated by $P(x)$ is less than 1×10^{-8} .

APPENDIX D

KIMath Addresses

CALCULATION SUBROUTINES

<u>ADDRESS</u>	<u>NAME</u>
F808	ADD
F800	SUB
F90B	MULPLY
FA16	DIVIDE
FA9E	SQRT
FAE7	LOG
FB41	TENX
FB5C	TANX
FB78	ATANX

UTILITY ROUTINES

FEFO	SAVXY
FEF5	RCLXY
FEE8	IPREC
FE0A	PLOADX
FE23	PLOADY
FE8A	ULOADX
FEA2	ULOADY
FE3C	PSTRES
FEBA	USTRES
FD92	USRCLKP
FDC1	POLY
FD71	CLRX
FD7C	CLRY
FD87	CLRZ
FBC3	DECHEX
FCBF	XSY

WORKING STORAGE

PAGE ZERO <u>STARTING ADDRESS</u>	<u>NAME</u>
10	PREC
11	EXTRA
06	ARGXL
07	ARGXH
08	ARGYL
09	ARGYH
0A	RESL
0B	RESH
0E	KONL
0F	KONH
01	NKON
05	DEG
03	CNT

COMPUTATION REGISTERS

<u>STARTING ADDRESS</u>	<u>NAME</u>
0235	RX
0235	SX (sign)
0246	EX (exponent)
0247	RY
0247	SY
0258	EY
0259	RZ
0259	SY
026A	EY

CARD	LOC	CODE	CARD
1	0000		*=\$0000
2			COPYRIGHT MUS TECHNOLOGY, INC. SEPT., 1976
3			ALL RIGHTS RESERVED. REV. 0
4			
5			
6			
7			
8			
9			
10			
11			
12			
13			
14			
15			
16			
17			
18			
19			
20			
21			
22			
23			
24			
25			
26			
27	0000		N =*=*+1
28	0001		NKON *==*+1
29	0002		J *==*+1
30	0003		CNT *==*+1
31	0004		LENGTH *==*
32	0004		CNTA *==*+1
33	0005		DEG *==*+1
34	0006		ARGXL *==*+1
35	0007		ARGXH *==*+1
36	0008		ARGYL *==*+1
37	0009		ARGYH *==*+1
38	000A		KES *==*+2
39	000C		PTR *==*+2
40	000E		KUN *==*+1
41	000F		KONH *==*+1
42	0010		PREC *==*+1
43	0011		EXTRA *==*+1
44	0012		TEMP *==*+1
45	0013		TEMP1 *==*+1
46	0014		OVERR *==*+1
47	0015		TMPX *==*+1
48	0016		TMPY *==*+1
49	0017		*=\$0200
50	0200		KA *==*+LEN+1
51	0212		KB *==*+LEN+1
52	0224		RQ *==*+LEN

CARD = LOC	CODE	CARD	
53 0235		RX *=*	
54 0235		SX *=*+LEN	
55 0246		EX *=*+1	
56 0247		KY *=*	
57 0247		SY *=*+LEN	
58 0258		EY *=*+1	
59 0259		RZ *=*+0	
60 0259		SZ *=*+LEN	
61 026A		EZ *=*+1	
62 026B		RM *=*+LEN+1	
63 027D		RN *=*+LEN+1	
64 028F		RAMCOD *=*+3	
65 0292		RAMA *=*+3	
66 0295		RAMB *=*+5	
67 029A		*=\$F800	
68		FLUTING POINT ADD-SUBTRACT ROUTINE.	
69			
70			
71 F800 AD 47 02	SUB	LDA SY	
72 F803 49 80		EOR =\$80	
73 F805 80 47 02		STA SY	
74 F808 AD 35 02	ADD	LDA SX	
75 F80B 4D 47 02		EOR SY	
76 F80E 85 12 02		STA TEMP	
77 F810 F8		SED	
78		CLEAR WORKING STORAGE.	
79			
80			
81 F811 20 B8 FB		JSR CLEAR	
82			
83		TEST RX FOR ZERO.	
84			
85 F814 20 A6 FC		JSR XZTST	
86			
87		TEST RY FOR ZERO.	
88			
89 F817 F0 0E		BEQ ADD2	
90 F819 20 B3 FC		JSR YZTST	
91 F81C F0 0C		BEQ ADD3	
92 F81E 24 12		BIT TEMP	
93 F820 50 70		BVC ADD6	
94		IF THE SIGNS OF THE EXPONENTS	
95		DIFFER THEN SWAP RX AND RY.	
96			
97			
98 F822 2C 35 02	ADD1	BIT SX	
99 F825 50 03		BVC ADD3	
100 F827 20 BF FC	ADD2	JSR XSY	
101 F82A F8	ADD3	SED	
102 F82B 24 12		BIT TEMP	
103 F82D 70 03		BVS ADD31	
104 F82F 4C AC F8		JMP ADD9	

CARD	LOC	CODE	CARD	
105	F832	AD 46 02	ADD31	LDA EX
106	F835	18		CLC
107	F836	6D 58 02		ADC EY
108	F839	80 12		BCS ADD5
109	F83B	85 03	ADD4	STA CNT
110				COMPUTE THE HEX VALUE OF THE
111				BCD DIFFERENCE OF THE EXPONENTS.
114	F83D	20 C3 FB		JSR DECHEX
115	F840	C5 00		CMP N
116	F842	B0 09		BCS ADD5
117				MOVE RY TO KB.
120	F844	20 1C FC		JSR RBERY
121				ALIGN DECIMAL POINTS.
124	F847	20 EE FB		JSR RSBCNT
125				ROUND RB OFF.
128	F84A	20 07 FC		JSR RBOFF
129	F84D	AU 46 02	ADD5	LDA EX
130	F850	8D 6A 02		STA EZ
132				MOVE RX TO RA.
134	F853	20 29 FC		JSR RAERX
135	F856	24 12		BIT TEMP
136	F858	30 6A		BMI ADD13
137				ADD KB TO RA.
139	F85A	20 5C FC		JSK RAPRB
140	F85D	AD 00 02		LDA RA
141	F860	F0 18		BEC ADD120
142	F862	20 A6 FB		JSR RSRA
143				CORRECT SIGN AND EXPONENT.
145				
146	F865	AD 46 J2		LDA EX
147	F868	38		SEC
148	F869	2C 35 02		BIT SX
149	F86C	50 19		BVC ADD110
150	F86E	E9 01		SBC =1
151	F870	80 6A 02		STA EZ
152	F873	D0 08		BNE ADD120
153	F875	A9 8F		LDA =\\$8F
154	F877	2D 35 02		AND SX
155	F87A	4C 80 F8		JMP ADD12
156	F87D	AD 35 02	ADD120	LDA SX

CARD	LOC	CODE	CARD
157	F880	80 59 02	ADD12 STA SZ
158			MOVE RA TO RZ.
159			
160			
161	F883	20 4F FC	ADD121 JSR RZERA
162	F886	60	RTS
163	F887	69 00	ADD110 ADC =0
164	F889	80 6A 02	STA EZ
165	F88C	90 EF	BCC ADD120
166			
167			SET RZ=9.9...9E99
168			
169	F88E	20 D2 FC	JSR INFIN
170	F891	60	RTS
171			COMPARE ABS(RX) TO ABS(RY).
172			
173	F892	20 82 FC	ADD6 JSR COMPXY
175	F895	A5 04	LDA CNTA
176	F897	F0 03	BEC ADD8
177			
178			SWAP RX AND RY,
179			SO THAT RX HAS THE
180			LARGEST ABS. VALUE.
181			
182	F899	20 BF FC	ADD7 JSR XSY
183	F89C	A0 46 02	ADD8 LDA EX
184	F89E	CD 58 02	CMP EY
185	F8A2	F0 05	BEQ ADD81
186	F8A4	90 F3	BCC ADD7
187	F8A6	40 22 F8	JMP ADD1
188	F8A9	40 2A F8	ADD81 JMP ADD3
189			
190			COMPUTE THE ABSOLUTE VALUE
191			OF THE SIGNED DIFFERENCE OF
192			THE EXPONENTS.
193			
194	F8AC	38	ADD9 SEC
195	F8AD	20 35 02	BIT SX
196	F8B0	70 09	BVS ADD10
197	F8B2	AD 46 02	LDA EX
198	F8B5	ED 58 02	SBC EY
199	F8B8	40 38 F8	JMP ADD4
200	F8BB	AD 58 02	ADD10 LDA EY
201	F8BE	ED 46 02	SBC EX
202	F8C1	40 38 F8	JMP ADD4
203			
204			SUBTRACT RB FROM RA.
205			
206	F8C4	20 70 FC	ADD13 JSR RAMRB
207	F8C7	AD 46 02	LDA EX
208	F8CA	80 6A 02	STA EZ

CARD	LOC	CODE	CARD
209	F8CD	AD 35 02	LDA SX
210	F8D0	8D 59 02	STA SZ
211			TEST RA FOR ZERO.
212			
213			
214	F8D3	20 99 FB	JSR AZTST
215	F8D6	F0 18	BEQ ADD18
216	F8D8	AD 01 02	ADD15 LDA RA+1
217	F8DB	8D A6	BNE ADD121
218			
219			IF RA+1 IS ZERO THEN
220			LEFT SHIFT RA ONE DIGIT.
221			
222	F8DD	20 91 FB	JSR LSRA
223	F8E0	2C 59 02	ADD17 BIT SZ
224	F8E3	38	SEC
225	F8E4	AD 6A 02	LDA EZ
226	F8E7	50 0B	BVC ADD20
227	F8E9	69 00	ADC =0
228	F8EB	8D 6A 02	STA EZ
229	F8EE	90 EB	BCC ADD15
230			SET RZ EQUAL TO ZERO.
231			
232	F8F0	20 87 FD	ADD18 JSK CLRZ
233	F8F3	80	ADD19 RTS
234			AJUST SIGN AND EXPONENT
235			OF THE ANSWER.
236			
237			
238			
239	F8F4	E9 01	ADD20 SBC =1
240	F8F6	8D 6A 02	STA EZ
241	F8F9	8D DD	BCS ADD15
242	F8FB	A9 01	LDA =1
243	F8FD	8D 6A 02	STA EZ
244	F900	A9 40	LDA =\\$40
245	F902	8D 59 02	CKA SZ
246	F905	8D 59 02	STA SZ
247	F908	4C D8 F8	JMP ADD15
248			FLOATING POINT PRODUCT ROUTINE.
249			
250			
251	F908	F8	MLTPLY SED
252			CLEAR WORKING STORAGE.
253			
254			
255	F90C	20 B8 FB	JSR CLEAR
256	F90F	A9 00	LDA =0
257	F911	85 03	STA CNT
258	F913	85 13	STA TEMP1
259			
260			TEST RA FOR ZERO.

CARD	LOC	CODE	CARD
261			
262	F915	20 A6 FC	JSR XZTST
263	F918	F0 05	BEQ MULT1
264			
265			TEST RY FOR ZERO.
266			
267	F91A	20 B3 FC	JSR YZTST
268	F91D	D0 08	BNE MULT3
269			
270			SET RZ EQUAL TO ZERO.
271			
272	F91F	20 87 FD	MULT1 JSR CLRZ
273	F922	60	RTS
274			
275			MOVE RA TO RZ.
276			
277	F923	20 4F FC	MULT2 JSR RZERA
278	F926	60	RTS
279			
280			MOVE RY TO RB.
281			
282	F927	20 1C FC	MULT3 JSR RBERY
283			
284			MOVE RX TO RG.
285			
286	F92A	20 36 FC	JSR RGERX
287			
288			FORM PRODUCT OF MANTISSAS.
289			
290	F92D	20 53 FA	JSR MLT
291			
292			FIGURE THE SIGN AND EXPONENT OF
293			THE ANSWER FOR THE MULTIPLY
294			AND DIVIDE ROUTINES.
295			
296	F930	AD 47 02	MULT4 LDA SY
297	F933	40 35 02	EOR SX
298	F936	85 12	STA TEMP
299	F938	24 12	BIT TEMP
300	F93A	AU 46 02	LDA EX
301	F93D	70 1E	BVS MD100
302	F93F	18	MD1 CLC
303	F940	60 58 02	ADC EY
304	F943	90 2B	BCC MD2
305	F945	D0 61	BNE MD59
306	F947	A5 13	LDA TEMP1
307	F949	F0 15	BEQ MDUV2
308	F94B	20 35 02	BIT SX
309	F94F	70 67	BVS MD7
310	F950	A5 04	LDA CNTA
311	F952	F0 5F	BEQ MD61
312	F954	A9 00	LDA =0

CARD	LOC	CODE	CARD	STA	CNTA
313	F956	85 04		LDA	=\$99
314	F958	A9 99	MDOV1	JMP	MD2
315	F95A	4C 70	F9	MDO100	JMP MD10
316	F95D	4C C3	F9	MDOV2	BIT SX
317	F960	2C 35	02	BVC	MD61
318	F963	50 4E		LDA	RA
319	F965	AD 00	02	BEQ	MDOV1
320	F968	F0 EE		JSR	RSRA
321	F96A	20 A6	F8		MDOV1
322	F96D	4C 58	F9	JMP	
323	F970	80 6A	02	MD2	STA EZ
324	F973	D0 64		BNE	MD11
325	F975	AD 35	02	LDA	SX
326	F978	29 BF		AND	=\$BF
327	F97A	80 59	02	MD3	STA SZ
328	F97D	A5 12		LDA	TEMP
329	F97F	30 3A		BMI	MD8
330	F981	A9 7F		LDA	=\$7F
331	F983	2D 59	02	AND	SZ
332	F986	8D 59	02	MD5	STA S7
333	F989	A5 13		LDA	TEMP1
334	F98E	D0 52		BNE	DIVEXT
335	F98D	AU 00	02	LDA	RA
336	F990	F0 13		BEG	M051
337	F992	20 A6	F8	JSR	RSRA
338	F995	AU 6A	02	LDA	E7
339	F998	2C 59	02	BIT	SZ
340	F99B	70 6B		BVS	MD9
341	F99D	18		CLC	
342	F99E	69 U1		ADC	=1
343	F9A0	F0 0C		BEC	MD6
344	F9A2	80 6A	02	STA	E7
345	F9A5	4C 23	F9	MD51	JMP MULT2
346	F9A8	AU 35	02	MD59	LDA SX
347	F9AB	80 59	02		STA SZ
348	F9AE	2C 59	02	MD6	BIT SZ
349	F9B1	70 04		BVS	MD7
350	F9B3	20 U2	FC	MD61	JSR INFIN
351	F9B6	60		RTS	
352	F9B7	20 87	F0	MD7	JSR CLRZ
353	F9BA	60		RTS	
354	F9BB	A9 80		MD8	LDA =\$A0
355	F9BD	0D 59	02		CRA SZ
356	F9C0	4C 86	F9		JMP MD5
357	F9C3	38		MU10	SEC
358	F9C4	ED 58	02		SBC EY
359	F9C7	B0 A7			BCS MD2
360	F9C9	38			SEC
361	F9CA	AD 58	02		LDA EY
362	F9CC	ED 46	02		SBC EX
363	F9D0	80 6A	02		STA EZ
364	F9D3	AU 47	02		LDA SY

CARD	LOC	CODE	CARD	JMP	MD3
365	F9D6	4C 7A F9		LDA	SX
366	F9D9	A0 35 02	MD11	JMP	MD3
367	F9DC	4C 7A F9		LDA	CNTA
368	F9DF	A5 04	DVEXT	BEQ	MD51
369	F9E1	F0 C2		BIT	SZ
370	F9E3	2C 59 02	DVEXT0	LDA	EZ
371	F9E6	A0 6A 02		SEC	
372	F9E9	38		BVC	DVEXT2
373	F9FA	50 0A		ADC	=0
374	F9EC	69 00		BEQ	MD6
375	F9EE	F0 BE	DVEXT1	STA	EZ
376	F9F0	8D 6A 02		JMP	MULT2
377	F9F3	4C 23 F9	DVEXT2	BEQ	DVEXT3
378	F9F6	F0 05		SBC	=1
379	F9F8	E9 01		JMP	DVEXT1
380	F9FA	4C F0 F9	DVEXT3	LDA	SZ
381	F9FB	A0 59 02		ORA	=\\$40
382	FA00	09 40		STA	SZ
383	FA02	8D 59 02		JMP	DVEXT0
384	FA05	4C E3 F9	MD9	SEC	
385	FA08	38		SBC	=1
386	FA09	E9 01		BEQ	MD22
387	FA0B	F0 06		STA	EZ
388	FA0D	8D 6A 02		JMP	MULT2
389	FA10	4C 23 F9	MD22	JMP	MD2
390	FA13	4C 70 F9			
391				FLOATING POINT DIVIDE ROUTINE.	
392					
393					
394	FA16	F8		DIVIDE SED	
395					
396					
397					
398	FA17	20 B3 FC		JSR	YZTST
399	FA1A	F0 97		BEQ	MD61
400					
401					
402					
403	FA1C	20 A6 FC		JSR	XZTST
404	FA1F	F0 96		BEQ	MD7
405					
406					
407					
408	FA21	20 B8 FB		JSR	CLEAR
409					
410					
411					
412	FA24	20 29 FC		JSR	RAERX
413					
414					
415					
416	FA27	20 1C FC		JSR	RBERY

CARD	=	LCC	CODE	CARD
417				COMPARE RX TO RY.
418				JSR COMPXY
419				FORM QUOTIENT.
420	FA2A	20 82 FC		JSR DIV
421				
422				
423				
424	FA2D	20 75 FA		COMPUTE SIGN AND EXPONENT OF
425				ANSWER.
426				
427				
428				
429	FA30	A9 01	DIV6	LDA =1
430	FA32	85 13		STA TEMP1
431	FA34	AD 47 02		LDA SY
432	FA37	49 40		EOR =\\$40
433	FA39	80 47 02		STA SY
434	FA3C	20 43 FC		JSR RAERQ
435	FA3F	AD 01 02		LDA RA+1
436	FA42	00 03		BNE DIV7
437	FA44	20 91 FB	DIV7	JSR LSKA
438	FA47	20 30 F9		JSR MULT4
439	FA4A	AD 47 02		LDA SY
440	FA4D	49 40		EOR =\\$40
441	FA4F	8D 47 02		STA SY
442	FA52	60		RTS
443				
444				THIS ROUTINE COMPUTES THE
445				PRODUCT OF THE MANTISSAS
446				OF THE ARGUMENTS BY REPEATED
447				ADDITION. THE RESULT IS BUILT
448				IN RA.
449				
450	FA53	A5 00	MLT	LDA N
451	FA55	85 02		STA J
452	FA57	C6 02		DEC J
453	FA59	A6 02	MLT0	LDX J
454	FA58	BD 24 02		LDA RQ,X
455	FA5E	85 03		STA CNT
456	FA60	C6 03	MLT1	DFC CNT
457	FA62	30 06		BMI MLT2
458	FA64	20 5C FC		JSR RAPKB
459	FA67	4C 60 FA		JMP MLT1
460	FA6A	20 A6 FB	MLT2	JSR RSRA
461	FA6D	C6 02		DEC J
462	FA6F	10 E8		BPL MLT0
463	FA71	20 91 FB		JSR LSRA
464	FA74	60		RTS
465				
466				THIS ROUTINE COMPUTES THE
467				QUOTIENT OF RA AND RB BY
468				REPEATED SUBTRACTION. THE

CARD	=	LOC	CODE	CARD	RESULT IS BUILT IN R0.
469				DIV	LDA =0
470					STA J
471	FA	75	A9 00	DIVO	LDA =0
472	FA	77	85 02		STA CNT
473	FA	79	A9 00		INC CNT
474	FA	7B	85 03		BNE DIV1
475	FA	7D	20 75	FC	JSR RAMRB
476	FA	80	90 04		BCC DIV2
477	FA	82	E6 03		INC CNT
478	FA	84	DU F7		JSR RAPRB
479	FA	86	20 5C	FC	JSR LSRA
480	FA	89	20 91	FC	LDX J
481	FARC	A6	02		LDA CNT
482	FA	8E	A9 03		STA R0,X
483	FA	90	9U 24	02	INC J
484	FA	93	E6 02		LDA J
485	FA	95	A5 02		CMP N
486	FA	97	C5 0J		BEQ DIV0
487	FA	99	F0 DE		BCC DIV0
488	FA	9B	90 DC		RTS
489	FA	9D	60		
490					
491					THIS ROUTINE COMPUTES THE
492					SQUARE ROOT OF A FLOATING POINT
493					NUMBER BETWEEN 1 AND 100 BY
494					HERONS METHOD.
495					
496	FA	9E	A9 07	SQRT	LDA =7
497	FA	A0	85 U1		STA NKUN
498	FA	A2	20 F8	FC	JSR MVXN
499	FA	A5	20 87	FD	JSR CLRZ
500	FA	A8	AJ 07		LDA =7
501	FA	AA	8U 5A	02	STA R7+1
502	FA	AD	A9 08		LDA =8
503	FA	AF	8D 5B	02	STA RZ+2
504	FA	AH	20 14	FD	JSR MVZN
505	FA	AB	20 20	FD	JSR MVMY
506	FA	BB	20 2C	FD	JSR MVNX
507	FA	BB	20 16	FA	JSR DIVIDE
508	FA	BF	20 10	FD	JSR MVZY
509	FA	C1	20 1C	FD	JSR MVMX
510	FA	C4	20 08	F8	JSR ADD
511	FA	C7	20 0C	FD	JSR MVZX
512	FA	CA	20 7C	FD	JSR CLRY
513	FA	CD	A9 40		LDA =\$40
514	FA	CF	8D 47	02	STA RY
515	FA	D2	A9 05		LDA =5
516	FA	D4	8D 48	02	STA RY+1
517	FA	D7	A9 01		LDA =1
518	FA	DC	8D 58	02	STA EY
519	FA	DF	20 08	F9	JSR MLTPLY
520	FA	EF	20 14	FD	JSR MVZN

CARD = LOC CODE

521 FAE2 C6 01
 522 FAE4 10 CF
 523 FAE6 60

CARD DEC NKON
 BPL SQRT0
 RTS

THIS ROUTINE COMPUTES THE
 COMMON LOG OF A FLOATING POINT
 NUMBER BETWEEN SQRT(-1) AND SQRT(10).

524		LOG	LDA =14
525	FAE7 A9 0E		STA N
526	FAE9 85 00		JSR SETKCN
527	FAEB 20 FC FE		JSR MVXN
528	FAEE 20 F8 FC		JSR CLR Y
529	FAF1 20 7C FD		LDA =1
530	FAF4 A9 01		STA RY+1
531	FAF6 8D 48 02		JSR SUB
532	FAF9 20 00 F8		JSR MVNX
533	FAFC 20 2C FD		JSR CLR Y
534	FAFF 20 7C FD		LDA =1
535	F802 A9 01		STA RY+1
536	F804 8D 48 02		JSR MVZN
537	F807 20 18 FD		JSR ADD
538	F80A 20 08 F8		JSR MVZY
539	F80D 20 10 FD		JSR MVNX
540	F810 20 2C FD		JSR DIVIDE
541	F813 20 16 FA		JSR MVZN
542	F816 20 18 FD		JSR MVZX
543	F819 20 0C FD		JSR MVZY
544	F81C 20 10 FD		JSR MLTPLY
545	F81F 20 08 F9		LDA =4
546	F822 A9 04		STA DEG
547	F824 85 05		LDA =0
548	F826 A9 00		LOGEND STA NKON
549	F828 85 01		JSR POLY
550	F82A 20 C1 FD		JSR MVNY
551	F82D 20 30 FD		LUNDO JSR MVZX
552	F830 20 0C FD		MLTPLY
553	F833 20 08 F9		CHCP LDA =J
554	F836 A9 00		LDX =LEN/2-1
555	F838 A2 07		CHUPO STA RZ+y,X
556	F83A 90 62 02		DEX
557	F83D CA		BPL CHOPO
558	F83E 10 FA		RTS

THIS ROUTINE COMPUTES THE
 COMMON ANTI-LOG OF A FLOATING
 POINT NUMBER BETWEEN 0 AND 1.

564		TENX	LDA #12
565	F841 A9 0C		STA N
566	F843 85 00		JSR SETKCN
567	F845 20 FC FE		JSR MVXZ
568	F848 20 F0 FC		

CARD	LOC	CODE	CARD	
573	FB4B	A9 06	LDA	=6
574	FB4D	85 05	STA	DEG
575	FB4F	A9 2E	LDA	=46
576	FB51	85 01	STA	NKON
577	FB53	20 C1 FD	JSR	POLY
578	FB56	20 10 FD	JSR	MVZY
579	FB59	4C 30 FB	JMP	LGN00
580				
581			THIS ROUTINE COMPUTES THE	
582			TANGENT OF A FLOATING POINT NUMBER	
583			BETWEEN 0 AND PI/4.	
584				
585	FB5C	A9 0E	TANX	LDA =14
586	FB5E	85 00		STA N
587	FB60	20 FC FE		JSR SETK0N
588	FB63	20 F8 FC		JSR MVXN
589	FB66	20 EC FC		JSR MVXY
590	FB69	20 0B F9		JSR MLTPLY
591	FB6C	20 36 FB		JSR CHOP
592	FB6F	A9 05		LDA =5
593	FB71	85 05		STA DEG
594	FB73	A9 64		LDA =100
595	FB75	4C 28 FB		JMP LOGEND
596				
597			THIS ROUTINE COMPUTES THE	
598			ARCTANGENT OF A FLOATING POINT NUMBER	
599			BETWEEN 0 AND 1.	
600				
601	FB78	A9 0E	ATANX	LDA =14
602	FB7A	85 00		STA N
603	FB7C	20 FC FE		JSR SETK0N
604	FB7F	20 F8 FC		JSR MVXN
605	FB82	20 EC FC		JSR MVXY
606	FB85	20 0B F9		JSR MLTPLY
607	FB88	A9 07		LDA =7
608	FB8A	85 05		STA DEG
609	FB8C	A9 9C		LDA =156
610	FB8E	4C 28 FB		JMP LOGEND
611			LEFT SHIFT RA ONE DIGIT.	
612				
613				
614	FB91	A2 00	LSRA	LDX =0
615	FB93	BD 01 02	LSRAO	LDA RA+1,X
616	FB96	9D 00 02		STA RA,X
617	FB99	E8		INX
618	FB9A	E4 00		CPX N
619	FB9C	90 F5		BCC LSKAO
620	FB9E	F0 F3		BEC LSRAO
621	FBA0	A9 00		LDA =?
622	FBA2	9D 00 02		STA RA,X
623	FBA5	60		RTS
624				

CARD	LOC	CODE	CARD	RIGHT SHIFT RA ONE DIGIT.
625			RSRA	LDX N
626				DEX
627	FBA6	A6 00	RSRAO	LDA RA,X
628	FBA8	CA		STA RA+1,X
629	FBA9	BD 00 02		
630	FBAE	9D 01 02		
631	FBAF	CA		DEX
632	FBB0	10 F7		BPL RSRAO
633	FBB2	A9 00		LDA =?
634	FBB4	8D 00 02		STA RA
635	FBB7	60		RTS
636				CLEAR WORKING STORAGE.
638				
639	FBB8	A2 34	CLEAR	LDX =LEN*3+1
640	FBB9	A9 00		LDA =0
641	FBC0	9D 00 02	AZ0	STA RA,X
642	FBCF	CA		DEX
643	FBC0	10 FA		BPL AZ0
644	FBC2	60		RTS
645				
646				CONVERT THE CONTENTS OF CNT
647				FROM BCD TO HEX AND STORE THE
648				RESULT IN CNT.
649				
650	FBC3	F8	DECHEX	SED
651	FBC4	A2 00		LDX =0
652	FBC6	38		SEC
653	FBC7	A5 03	DHCNV1	LDA CNT
654	FBC9	E9 16		SBC =\$16
655	FBCB	90 06		BCC DHCNV2
656	FBCD	85 03		STA CNT
657	FBCF	E8		INX
658	FBD0	4C C7 FB		JMP DHCNV1
659	FBD3	D8	DHCNV2	CLD
660	FBD4	A5 03		LDA CNT
661	FBD6	C9 0A		CMP =\$0A
662	FBD8	90 04		BCC DHCNV3
663	FBD9	29 0F		AND =\$0F
664	FBDc	69 09		ADC =\$09
665	FBDc	86 03	DHCNV3	STX CNT
666	FBE0	06 03		ASL CNT
667	FBE2	06 03		ASL CNT
668	FBE4	06 03		ASL CNT
669	FBE6	06 03		ASL CNT
670	FBE8	05 03		ORA CNT
671	FBEA	85 03		STA CNT
672	FBEc	F8		SED
673	FBED	60	DHCNVE	RTS
674				
675				RIGHT SHIFT RB CNT TIMES.
676				

CARD	LOC	CODE	CARD	LOC	CNT
677	FBEE	A5 03	KSBCNT	LDA	CNT
678	FBFO	F0 29		BEC	RBOFE
679	FBF2	A6 00		LDX	N
680	FBF4	BD 12 02	KSBC	LDA	RB,X
681	FBF7	9D 13 02		STA	RB+1,X
682	FBFA	CA		DEX	
683	FBFB	10 F7		BPL	RSBC
684	FRFD	A9 00		LDA	=)
685	FBFF	8D 12 02		STA	RB
686	FC02	C6 03		DEC	CNT
687	FC04	00 E8		BNE	RSBCNT
688	FC06	60		RTS	
689					
690			ROUND	RB	OFF.
691					
692	FC07	A6 00	RBCFF	LDX	N
693	FC09	9D 13 02		LDA	RB+1,X
694	FC0C	C5 05		CMP	=5
695	FC0E	BD 12 02	RBOF	LDA	RB,X
696	FC11	69 90		ADC	=\\$90
697	FC13	29 0F		AND	=\\$0F
698	FC15	9D 12 02		STA	RB,X
699	FC18	CA		DEX	
700	FC19	10 F3		BPL	RBCF
701	FC1B	60	RBOFE	RTS	
702					
703			MOVE	RY	TO RB.
704					
705	FC1C	A6 00	RBERY	LDX	N
706	FC1E	CA		DEX	
707	FC1F	BD 48 02	RBRY	LDA	RY+1,X
708	FC22	90 13 02		STA	RB+1,X
709	FC25	CA		DEX	
710	FC26	10 F7		BPL	RBRY
711	FC28	60		RTS	
712					
713			MOVE	RX	TO RA.
714					
715	FC29	A6 00	RAERX	LDX	N
716	FC2B	CA		DEX	
717	FC2C	BD 36 02	RAKXO	LDA	RX+1,X
718	FC2F	9D 01 02		STA	RA+1,X
719	FC32	CA		DEX	
720	FC33	10 F7		BPL	RAKXO
721	FC35	60	RAKXE	RTS	
722					
723			MOVE	RX	TO RQ.
724					
725	FC36	A6 00	RQERX	LDX	N
726	FC38	CA		DEX	
727	FC39	BD 36 02	RQRX	LDA	RX+1,X
728	FC3C	9D 24 02		STA	RQ,X

CARD	LUC	CODE	CARD
729	FC3F	CA	DEX
730	FC40	10 F7	BPL RQRX
731	FC42	60	RTS
732			MOVE RQ TO RA.
734			
735	FC43	A6 00	RAERQ LDX N
736	FC45	BD 24 02	RARQ LDA RQ,X
737	FC48	9D 01 02	STA RA+1,X
738	FC4B	CA	DEX
739	FC4C	10 F7	BPL RARQ
740	FC4E	60	RTS
741			MOVE RA TO RZ.
743			
744	FC4F	A6 00	RZERA LDX N
745	FC51	CA	DEX
746	FC52	BD 01 02	RZRAO LDA RA+1,X
747	FC55	9D 5A 02	STA RZ+1,X
748	FC58	CA	DEX
749	FC59	10 F7	BPL RZRAO
750	FC5B	60	RTS
751			ADD RB TO RA.
753			
754	FC5C	A6 00	RAPRB LDX N
755	FC5E	18	CLC
756	FC5F	BD 00 02	APB LDA RA,X
757	FC62	7D 12 02	ADC RB,X
758	FC65	69 90	ADC =\$90
759	FC67	29 0F	AND =\$0F
760	FC69	9D 00 02	STA RA,X
761	FC6C	CA	DEX
762	FC6D	10 F0	BPL APB
763	FC6F	60	RTS
764			SUBTRACT RB FROM RA.
766			
767	FC70	A6 00	RAMRB LDX N
768	FC72	38	SEC
769	FC73	BD 00 02	AMB LDA RA,X
770	FC76	FD 12 02	SBC RB,X
771	FC79	29 0F	AND =\$0F
772	FC7B	9D 00 02	STA RA,X
773	FC7E	CA	DEX
774	FC7F	10 F2	BPL AMB
775	FC81	60	RTS
776			COMPARE RX TO RY.
778			
779	FC82	A9 00	COMPXY LDA =0
780	FC84	85 04	STA CNTA

CARU =	LOC	CODE	CARD
	781	FC86 A6 00	LDX N
	782	FC88 CA	DEX
	783	FC89 38	SEC
	784	FC8A BD 36 02	COML LDA RX+1,X
	785	FC8D FD 48 02	SBC RY+1,X
	786	FC90 CA	DEX
	787	FC91 10 F7	BPL COM1
	788	FC93 90 01	BCC COM2
	789	FC95 60	RTS
	790	FC96 E6 04	COMP INC CNIA
	791	FC98 60	RTS 1
	792		TEST RA FOR ZERO.
	793		X
	794		AZTST LDX N
	795	FC99 A6 00	INX
	796	FC9B E8	AZTSTO LDA RA,X
	797	FC9C BD 00 02	BNE XZTST1
	798	FC9F D0 11	DEX
	799	FCA1 CA	BPL AZTST0
	800	FCA2 10 F8	BMI XZTST2
	801	FCA4 30 0A	H02
	802		TEST RX FOR ZERO.
	803		XZTST LDX N
	804	FCA6 A6 00	XZTSTO LDA RX,X
	805	FCA8 BD 35 02	BNE XZTST1
	806	FCA9 D0 05	DEX
	807	FCAE CA	BPL XZTST0
	808	FCAE 10 F8	XZTST2
	809	FCAE A9 00	LDA =0
	810	FCAE 60	XZTST1 RTS
	811	FCA2	P12
	812		TEST RY FOR ZERO.
	813		YZTST LDX N
	814		YZTSTO LDA RY,X
	815	FCE3 A6 00	BNE XZTST1
	816	FCB5 BD 47 02	DEX
	817	FCB8 D0 F8	BPL YZTST0
	818	FCB8 CA	BMI XZTST2
	819	FCCB 10 F8	P21
	820	FCCB 30 F1	SWAP RX AND RY.
	821		822
	823		XSY LDX =LEN
	824	FCCF A2 11	XSY1 LDA RX,X
	825	FCC1 BD 35 02	LDY RY,X
	826	FCC4 BC 47 02	STA RY,X
	827	FCC7 9D 47 02	TYA
	828	FCCA 98	STA RX,X
	829	FCCB 9D 35 02	DEX
	830	FCCF CA	BPL XSY1
	831	FCCF 10 F0	RTS
	832	FCD1 60	R32

CARD = LOC CODE

CARD

833			
834			SET RZ=9.9...9E99 AND OVERR=1.
835			
836	FCD2	A6 00	INFIN LDX N
837	FCD4	CA	DEX
838	FCD5	A9 09	LDA =9
839	FCD7	9D 5A 02	STA RZ+1,X
840	FCDA	CA	DEX
841	FCDB	10 FA	BPL INFO
842	FCCD	A9 99	LDA =\\$99
843	FCDF	8D 6A 02	STA EZ
844	FCE2	A9 00	LDA =0
845	FCE4	8D 59 02	STA SZ
846	FCE7	A9 01	LDA =1
847	FCE9	85 14	STA OVERR
848	FCEB	60	RTS

849
 850 THE FOLLOWING ROUTINES ARE USED
 851 TO MOVE THE CONTENTS FROM ONE
 852 REGISTER TO ANOTHER. THE NAMES ARE
 853 OF THE FORM MVSD, WHERE S STANDS
 854 FOR SOURCE AND D FOR DESTINATION.
 855

856	FCEC	A9 01	MVXY LDA =XY
857	FCEE	D0 4A	BNE MVTR
858	FCFO	A9 02	MVXZ LDA =XZ
859	FCF2	D0 46	BNE MVTR
860	FCF4	A9 03	MVXM LDA =XM
861	FCF6	D0 42	BNE MVTR
862	FCF8	A9 04	MVXN LDA =XN
863	FCFA	D0 3E	BNE MVTR
864	FCFC	A9 10	MVYX LDA =YX
865	FCFE	D0 3A	BNE MVTR
866	FD00	A9 12	MVYZ LDA =YZ
867	FD02	D0 36	BNE MVTR
868	FD04	A9 13	MVYM LDA =YM
869	FD06	D0 32	BNE MVTR
870	FD08	A9 14	MVYN LDA =YN
871	FD0A	D0 2E	BNE MVTR
872	FD0C	A9 20	MVZX LDA =ZX
873	FD0E	D0 2A	BNE MVTR
874	FD10	A9 21	MVZY LDA =ZY
875	FD12	D0 26	BNE MVTR
876	FD14	A9 23	MVZM LDA =ZM
877	FD16	D0 22	BNE MVTR
878	FD18	A9 24	MVZN LDA =ZN
879	FD1A	D0 1E	BNE MVTR
880	FD1C	A9 30	MVMX LDA =MX
881	FD1E	D0 1A	BNE MVTR
882	FD20	A9 31	MVMY LDA =MY
883	FD22	D0 16	BNE MVTR
884	FD24	A9 32	LDA =MZ

CARD	LOC	CODE	CARD
865	FD26	DJ 12	BNE MVTR
866	FC28	A9 34	LDA =MN
887	FC2A	D0 0E	BNE MVTR
888	FD2C	A9 40	MVNX LDA =NX
889	FD2F	D0 0A	BNE MVTR
890	FC30	A9 41	MVNY LDA =NY
891	FD32	D0 06	BNE MVTR
892	FD34	A9 42	MVNZ LDA =NZ
893	FD36	D0 02	BNE MVTR
894	FD38	A9 43	MVNW LDA =NM
895	FD3A	48	MVTR PHA
896	FD3B	A2 0B	LDX =11
897	FD3D	BD 65 FD	MVTRO LDA MOVR,X
898	FD40	9D 8F 02	STA RAMCUD,X
899	FD43	CA	DEX
900	FD44	10 F7	BPL MVTRO
901	FD46	68	PLA
902	FD47	48	PHA
903	FD48	29 0F	AND =\$0F
904	FD4A	AA	TAX
905	FD4B	BD 60 FD	LDA TAB,X
906	FD4E	8D 95 02	STA RAMB
907	FD51	68	PLA
908	FD52	4A	LSR A
909	FD53	4A	LSR A
910	FD54	4A	LSR A
911	FD55	4A	LSR A
912	FD56	AA	TAX
913	FD57	BD 60 FD	LDA TAB,X
914	FD5A	8D 92 02	STA RAMA
915	FD5D	4C 8F 02	JMP RAMCUD
916	FD60	35	TAB .BYTE \$35,\$47,\$59,\$68,\$7D
916	FD61	47	
916	FD62	59	
916	FD63	68	
916	FD64	7D	
917	FD65	A2 11	MOVR LDX =LEN
918	FD67	BD 35 02	MOVRO LDA RX,X
919	FD6A	9D 47 02	STA RY,X
920	FD6D	CA	DEX
921	FD6E	10 F7	BPL MOVRO
922	FD70	60	RTS
923			SET RX EQUAL TO ZERO.
924			
925			
926	FD71	A2 11	CLRX LDX =LEN
927	FD73	A9 00	LDA =J
928	FD75	9D 35 02	CLRRO STA RX,X
929	FD78	CA	DEX
930	FD79	10 FA	BPL CLRRO
931	FD7B	60	RTS
932			

CARD	LOC	CODE	CARD
933			SET RY EQUAL TO ZERO.
934			
935	FD7C	A2 11	CLRY LDX =LEN
936	F07E	A9 00	LDA =0
937	F080	9D 47 02	CLRYO STA RY,X
938	F083	C4	DEX
939	F084	10 FA	BPL CLRYO
940	F086	60	RTS
941			SET RZ EQUAL TO ZERO.
942			
943	F087	A2 11	CLRZ LDX =LEN
945	F089	A9 00	LDA =0
946	F088	9D 59 02	CLRZO STA RZ,X
947	F08F	C4	DEX
948	F08F	10 FA	BPL CLRZO
949	F091	60	RTS
950			
951			THIS ROUTINE IS USED TO LOOK UP
952			THE COEFFICIENTS OF THE POLY-
953			NOMIALS USED IN THE APPROXIMATIONS
954			OF THE TRANSCENDENTAL FUNCTIONS.
955			
956	F092	20 7C FD	LOOKUP JSR CLRY
957	F095	A2 00	LDX =0
958	F097	A4 01	LDY NKON
959	F099	B1 0E	LDA (KON),Y
960	F09B	8D 47 U2	STA SY
961	F09E	C8	LKPO INY
962	F09F	B1 0E	LDA (KON),Y
963	F0A1	C9 F0	CMP =\$F0
964	F0A3	80 13	BCS LKPI
965	F0A5	48	PHA
966	F0A6	29 0F	AND =\$0F
967	F0A8	9D 49 02	STA RY+2,X
968	F0AB	68	PLA
969	F0AC	4A	LSR A
970	F0AD	4A	LSR A
971	F0AE	4A	LSR A
972	F0AF	4A	LSR A
973	F0B0	9D 48 02	STA RY+1,X
974	F0B3	E8	INX
975	F0B4	E8	INX
976	F0B5	4C 9E FD	JMP LKPO
977	F0B8	29 0F	LKPI AND =\$0F
978	F0BA	8D 58 02	STA EY
979	F0BD	C8	INY
980	F0BE	84 01	STY NKON
981	F0C0	60	RTS
982			
983			THIS ROUTINE EVALUATES POLYNOMIALS
984			BY MEANS OF THE NESTED MULTIPLICATION

CARD	=	LOC	CODE	CARD	
985				POLY	JSR MVZN
986					JSR MVZX
987	FDC1	20 14	FD		JSR LOOKUP
988	FDC4	20 0C	FD		JSR MLTPLY
989	FDC7	20 92	FD		JSR LOOKUP
990	FDCA	20 06	F9	POLYO	JSR MVZY
991	FDCD	20 92	FD		DEC DEG
992	FDC0	20 0C	FD		BPL POLYO
993	FDD3	20 08	F8		RTS
994	FDC6	20 1C	FD		
995	FDC9	20 10	FD		
996	FDC	C6 05			
997	FDEE	10 EA			
998	FDEJ	60			
999					
1000				THIS ROUTINE UNPACKS AN ARGUMENT	
1001				AND STORES THE RESULT IN RZ.	
1002					
1003	FDE1	A2 00		PGTARG	LDX =0
1004	FDE3	A0 00			LDY =0
1005	FDE5	B1 0C			LDA (PTR),Y
1006	FDE7	80 59	02		STA SZ
1007	FDEA	C8		PGTRGO	INY
1008	FDEB	C4 04			CPY LENGTH
1009	FDEF	F0 15			BEG PGTRG1
1010	FDEF	B1 0C			LDA (PTR),Y
1011	FDF1	48			PHA
1012	FDF2	29 0F			AND =\$0F
1013	FDF4	90 5B	02		STA RZ+2,X
1014	FDF7	68			PLA
1015	FDF8	4A			LSR A
1016	FDF9	4A			LSR A
1017	F DFA	4A			LSR A
1018	FDFB	4A			LSR A
1019	FDFC	90 5A	02		STA RZ+1,X
1020	- FDFD	E8			INX
1021	FEO0	E8			INX
1022	FEO1	4C EA FD			JMP PGTRGO
1023	FEO4	B1 0C		PGTRG1	LDA (PTR),Y
1024	FEO6	8D 6A	02		STA EZ
1025	FEO9	60			RTS
1026					
1027				THIS ROUTINE UNPACKS AN ARGUMENT	
1028				LOCATED AT (ARGXL,ARGXH) AND STORES	
1029				THE RESULTS IN RZ AND RX.	
1030					
1031	FE0A	A5 06		PLOADX	LDA ARGXL
1032	FE0C	85 0C			STA PTR
1033	FE0E	A5 07			LDA ARGXH
1034	FE10	85 0D			STA PTR+1
1035	FE12	A5 10			LDA PREC
1036	FE14	4A			LSR A

CARD	LOC	CODE	CARD	LOC	CODE
1037	FE15	69 01			
1038	FE17	85 04	ADC =1		
1C39	FE19	20 87 FD	STA LENGTH		
1040	FE1C	20 E1 FD	JSR CLRZ		
1041	FE1F	20 0C FD	JSR PGTARG		
1042	FE22	60	JSR MVZX		
1C43			RTS		
1044					
1045			THIS ROUTINE UNPACKS AN ARGUMENT		
1046			LOCATED AT (ARGYL,ARGYH) AND STORES		
1C47			THE RESULT IN RY AND RZ.		
1048	FE23	A5 08	PLOADY LDA ARGYL		
1049	FE25	85 0C	STA PTR		
1050	FE27	A5 09	LDA ARGYH		
1C51	FE29	85 0D	STA PTR+1		
1052	FE2B	A5 10	LDA PREC		
1053	FE2D	4A	LSR A		
1054	FE2E	69 01	ADC =1		
1055	FE30	85 04	STA LENGTH		
1056	FE32	20 87 FD	JSR CLRZ		
1057	FE35	20 E1 FD	JSR PGTARG		
1058	FE38	20 10 FD	JSR MVZY		
1059	FE3B	60	RTS		
1060					
1061			THIS ROUTINE PACKS THE CONTENTS		
1062			OF RZ INTO THE LOCATIONS STARTING		
1063			WITH ADDRESS (RES,RES+1).		
1064					
1065	FE3C	A2 00	PSTRES LDX =0		
1066	FE3E	A0 00	LDY =0		
1067	FE40	AD 59 02	LDA SZ		
1068	FF43	91 0A	STA (RES),Y		
1069	FE45	C8	INY		
1070	FE46	BD 5A 02	PTRES LDA RZ+1,X		
1071	FE49	0A	ASL A		
1072	FE4A	0A	ASL A		
1073	FE4B	0A	ASL A		
1074	FE4C	0A	ASL A		
1075	FE4D	1D 5B 02	ORA RZ+2,X		
1076	FE50	91 0A	STA (RES),Y		
1077	FE52	C8	INY		
1078	FE53	E8	INX		
1079	FE54	E8	INX		
1080	FE55	E4 10	CPX PREC		
1081	FE57	90 ED	BCC PTRES		
1082	FE59	AD 6A 02	LDA EZ		
1083	FE5C	91 0A	STA (RES),Y		
1084	FE5E	60	RTS		
1085					
1086			THIS ROUTINE CONVERTS AN ARGUMENT		
1087			FROM ASCII FORMAT TO COMPUTATIONAL		
1088			FORMAT AND STORES THE RESULT IN RZ.		

CARD	=	LOC	CODE	CARD
1089		FE5F	A0 00	UGTARG LDY =0
1090		FE61	B1 0C	LDA (PTR),Y
1091		FE63	8D 59 02	STA SZ
1092		FE66	C8	UGTARO INY
1093		FE67	C4 04	CPY LENGTH
1094		FE69	F0 0A	BEQ UGTARI
1095		FE6B	B1 0C	LDA (PTR),Y
1096		FE6D	29 0F	AND =\$0F
1097		FE6F	99 59 02	STA RZ,Y
1098		FE72	4C 66 FE	JMP UGTARO
1099		FE75	B1 0C	UGTARI LDA (PTR),Y
1100		FE77	0A	ASL A
1101		FE78	0A	ASL A
1102		FE79	0A	ASL A
1103		FE7A	0A	ASL A
1104		FE7B	8D 6A 02	STA EZ
1105		FE7E	C8	INY
1106		FE7F	B1 0C	LDA (PTR),Y
1107		FE81	29 0F	AND =\$0F
1108		FE83	0D 6A 02	ORA EZ
1109		FE86	8D 6A 02	STA EZ
1110		FE89	60	RTS
1111				
1112				
1113				THIS ROUTINE CONVERTS AN ARGUMENT
1114				FROM ASCII FORMAT TO COMP. FORMAT.
1115				THE ADDRESS OF THE ARG. IS FOUND IN
1116				(ARGXL,ARGXH) AND THE RESULT IS STORED
1117				IN RZ AND RX.
1118				
1119		FE8A	A5 06	ULOADX LDA ARGXL
1120		FE8C	85 0C	STA PTR
1121		FE8E	A5 07	LDA ARGXH
1122		FE90	85 0D	STA PTR+1
1123		FE92	A5 10	LDA PREC
1124		FE94	85 04	STA LENGTH
1125		FE96	E6 04	INC LENGTH
1126		FE98	20 87 FD	JSR CLRZ
1127		FE9B	20 5F FE	JSR UGTARG
1128		FE9E	20 0C FD	JSR MVZX
1129		FEA1	60	RTS
1130				
1131				THIS ROUTINE CONVERTS AN ARGUMENT
1132				FROM ASCII FORMAT TO COMP. FORMAT.
1133				THE ADDRESS OF THE ARG. IS FOUND IN
1134				(ARGYL,ARGYH) AND THE RESULT IS
1135				STORED IN RZ AND KY.
1136				
1137		FEA2	A5 08	ULCADY LUA ARGYL
1138		FEA4	85 0C	STA PTR
1139		FEA6	A5 09	LDA ARGYH
1140		FEA8	85 0D	STA PTR+1

CARD	LUC	CODE	CARD	
1141	FEAA	A5 10	LDA PREC	
1142	FEAC	B5 04	STA LENGTH	
1143	FEAE	E6 04	INC LENGTH	
1144	FE80	20 87 FD	JSR CLRZ	
1145	FEB3	20 5F FE	JSR UGTARG	
1146	FEB6	20 10 FD	JSR MVZY	
1147	FEB9	60	RTS	
1148				
1149			THIS ROUTINE CONVERTS THE CONTENTS	
1150			OF RZ TO ASCII FORMAT WHILE MOVING	
1151			THEM TO THE ADDRESS SPECIFIED BY	
1152			(RES,RES+1).	
1153				
1154	FFBA	A0 00	USTRES LDY =0	
1155	FEB8C	AD 59 02	LDA S7	
1156	FEBF	91 0A	STA (RES),Y	
1157	FEC1	C8	USTRSU INY	
1158	FEC2	C4 10	CPY PREC	
1159	FEC4	F0 02	BEQ USTRS1	
1160	FEC6	80 09	BCS USTRS2	
1161	FEC8	89 59 02	LDA RZ,Y	
1162	FECB	09 30	GRA =\$30	
1163	FECF	91 0A	STA (RES),Y	
1164	FECF	D0 F0	BNE USTRS0	
1165	FEU1	C8	USTRS2 INY	
1166	FED2	AD 6A 02	LDA EZ	
1167	FED5	4A	LSR A	
1168	FED6	4A	LSR A	
1169	FED7	4A	LSR A	
1170	FED8	4A	LSR A	
1171	FED9	09 30	GRA =\$30	
1172	FED9	91 0A	STA (RES),Y	
1173	FEDD	C8	INY	
1174	FEOF	AD 6A 02	LDA EZ	
1175	FEE1	29 OF	AND =\$0F	
1176	FEE3	09 30	ORA =\$30	
1177	FEE5	91 0A	STA (RES),Y	
1178	FEE7	60	RTS	
1179				
1180			THIS ROUTINE COMPUTES THE	
1181			INTERNAL PRECISION N FROM	
1182			PREC AND EXTRA. THE ACU IS	
1183			A BINARY ADD (UNSIGNED).	
1184				
1185	FEE8	18	IPREC CLC	
1186	FEE9	A5 10	LDA PREC	
1187	FEEB	65 11	ADC EXTRA	
1188	FED0	85 00	STA N	
1189	FEEF	60	RTS	
1190				
1191			SAVE THE PROCESSOR INDEX REGISTERS.	
1192				

CARD	LOC	CODE	CARD
1193	FEF0	86 15	SAVXY STX TMPX
1194	FEF2	84 16	STY TMPY
1195	FEF4	60	RTS
1196			RECALL THE PROCESSOR INDEX REGISTERS.
1197			
1198	FEF5	A6 15	RCLXY LDX TMPX
1199	FEF7	A4 16	LDY TMPY
1200	FEF9	60	RTS
1201	FEFA	07 FF	KADUR .WORD KONST
1202	FEFC	AD FA FE	SETKON LDA KADOR
1203	FEFF	85 0E	STA KON
1204	FF01	AD FB FE	LDA KACDR+1
1205	FF04	85 0F	STA KONH
1206	FF06	60	RTS
1207			THESE ARE THE COEFFICIENTS USED
1208			IN THE EVALUATION OF THE TRANSCENDENTAL
1209			FUNCTIONS.
1210			
1211			
1212			
1213	FF07	40	KONST .BYTE \$40,\$18,\$20,\$91,\$29,\$97,\$F1
1213	FF08	18	
1213	FF09	20	
1213	FF0A	91	
1213	FF0B	29	
1213	FF0C	97	
1213	FF0D	F1	
1214	FF0E	40	.BYTE \$40,\$55,\$34,\$27,\$38,\$70,\$F2
1214	FF0F	55	
1214	FF10	34	
1214	FF11	27	
1214	FF12	38	
1214	FF13	79	
1214	FF14	F2	
1215	FF15	40	.BYTE \$40,\$13,\$13,\$69,\$01,\$12,\$10,\$F1
1215	FF16	13	
1215	FF17	13	
1215	FF18	69	
1215	FF19	01	
1215	FF1A	12	
1215	FF1B	10	
1215	FF1C	F1	
1216	FF1D	40	.BYTE \$40,\$17,\$31,\$09,\$55,\$17,\$F1
1216	FF1E	17	
1216	FF1F	31	
1216	FF20	09	
1216	FF21	55	
1216	FF22	17	
1216	FF23	F1	
1217	FF24	40	.BYTE \$40,\$28,\$95,\$51,\$13,\$02,\$67,\$F1
1217	FF25	28	
1217	FF26	95	

CARD	=	LOC	CUDE	CARD
1217		FF27	51	
1217		FF28	13	
1217		FF29	02	
1217		FF2A	67	
1217		FF2B	F1	
1218		FF2C	40	.BYTE \$40,\$86,\$85,\$88,\$74,\$83,\$40,\$50,\$F1
1218		FF2D	86	
1218		FF2E	85	
1218		FF2F	88	
1218		FF30	74	
1218		FF31	83	
1218		FF32	40	
1218		FF33	50	
1218		FF34	F1	
1219		FF35	40	.BYTE \$40,\$93,\$26,\$42,\$67,\$F4
1219		FF36	93	
1219		FF37	26	
1219		FF38	42	
1219		FF39	67	
1219		FF3A	F4	
1220		FF3B	40	.BYTE \$40,\$25,\$54,\$91,\$79,\$60,\$F3
1220		FF3C	25	
1220		FF3D	54	
1220		FF3E	91	
1220		FF3F	79	
1220		FF40	6U	
1220		FF41	F3	
1221		FF42	40	.BYTE \$40,\$17,\$42,\$11,\$19,\$88,\$F2
1221		FF43	17	
1221		FF44	42	
1221		FF45	11	
1221		FF46	19	
1221		FF47	88	
1221		FF48	F2	
1222		FF49	40	.BYTE \$40,\$72,\$95,\$17,\$36,\$66,\$F2
1222		FF4A	72	
1222		FF4B	95	
1222		FF4C	17	
1222		FF4D	36	
1222		FF4E	66	
1222		FF4F	F2	
1223		FF50	40	.BYTE \$40,\$25,\$43,\$93,\$57,\$48,\$40,\$F1
1223		FF51	25	
1223		FF52	43	
1223		FF53	93	
1223		FF54	57	
1223		FF55	48	
1223		FF56	40	
1223		FF57	F1	
1224		FF58	40	.BYTE \$40,\$66,\$27,\$30,\$88,\$42,\$90,\$F1
1224		FF59	66	
1224		FF5A	27	

CARD	=	LOC	CODE	CARD
1224		FF5B	30	
1224		FF5C	88	
1224		FF5D	42	
1224		FF5E	90	
1224		FF5F	F1	
1225		FF60	00	.BYTE \$00,\$11,\$51,\$29,\$27,\$76,\$03,\$F0
1225		FF61	11	
1225		FF62	51	
1225		FF63	29	
1225		FF64	27	
1225		FF65	76	
1225		FF66	03	
1225		FF67	F0	
1226		FF68	00	.BYTE \$00,\$10,\$F0
1226		FF69	10	
1226		FF6A	F0	
1227		FF6B	40	.BYTE \$40,\$41,\$09,\$74,\$19,\$48,\$F4
1227		FF6C	41	
1227		FF6D	09	
1227		FF6E	74	
1227		FF6F	19	
1227		FF70	48	
1227		FF71	F4	
1228		FF72	40	.BYTE \$40,\$20,\$31,\$17,\$10,\$84,\$F4
1228		FF73	20	
1228		FF74	31	
1228		FF75	17	
1228		FF76	10	
1228		FF77	84	
1228		FF78	F4	
1229		FF79	40	.BYTE \$40,\$27,\$97,\$43,\$35,\$03,\$70,\$F3
1229		FF7A	27	
1229		FF7B	97	
1229		FF7C	43	
1229		FF7D	35	
1229		FF7E	03	
1229		FF7F	70	
1229		FF80	F3	
1230		FF81	40	.BYTE \$40,\$98,\$34,\$59,\$45,\$39,\$30,\$F3
1230		FF82	98	
1230		FF83	34	
1230		FF84	59	
1230		FF85	45	
1230		FF86	39	
1230		FF87	30	
1230		FF88	F3	
1231		FF89	40	.BYTE \$40,\$34,\$86,\$59,\$10,\$47,\$05,\$F2
1231		FF8A	39	
1231		FF8B	86	
1231		FF8C	59	
1231		FF8D	10	
1231		FF8E	47	

CARD	LOC	CODE	CARD
1231	FF8F	05	
1231	FF90	F2	
1232	FF91	40	.BYTE \$40,\$16,\$14,\$89,\$77,\$76,\$17,\$40,\$F1
1232	FF92	16	
1232	FF93	14	
1232	FF94	89	
1232	FF95	77	
1232	FF96	76	
1232	FF97	17	
1232	FF98	40	
1232	FF99	F1	
1233	FF9A	40	.BYTE \$40,\$78,\$53,\$98,\$17,\$62,\$29,\$10,\$F1
1233	FF9B	78	
1233	FF9C	53	
1233	FF9D	98	
1233	FF9E	17	
1233	FF9F	62	
1233	FFA0	29	
1233	FFA1	10	
1233	FFA2	F1	
1234	FFA3	40	.BYTE \$40,\$28,\$49,\$88,\$96,\$20,\$80,\$F3
1234	FFA4	28	
1234	FFA5	49	
1234	FFA6	88	
1234	FFA7	96	
1234	FFA8	20	
1234	FFA9	80	
1234	FFAA	F3	
1235	FFAB	C0	.BYTE \$C0,\$16,\$06,\$86,\$28,\$96,\$04,\$F2
1235	FFAC	16	
1235	FFAD	06	
1235	FFAE	86	
1235	FFAF	28	
1235	FFB0	96	
1235	FFB1	04	
1235	FFB2	F2	
1236	FFB3	40	.BYTE \$40,\$42,\$69,\$15,\$19,\$27,\$11,\$F2
1236	FFB4	42	
1236	FFB5	69	
1236	FFB6	15	
1236	FFB7	19	
1236	FFB8	27	
1236	FFB9	11	
1236	FFBA	F2	
1237	FFB0	C0	.BYTE \$C0,\$75,\$04,\$29,\$45,\$38,\$89,\$F2
1237	FFBC	75	
1237	FFBD	04	
1237	FFBE	29	
1237	FFBF	45	
1237	FFC0	38	
1237	FFC1	89	
1237	FFC2	F2	

CARD =	LUC	CODE	CARD	.BYTE
1238	FFC3	40		\$40,\$10,\$64,\$09,\$34,\$02,\$53,\$F1
1238	FFC4	10		
1238	FFC5	64		
1238	FFC6	09		
1238	FFC7	34		
1238	FFC8	02		
1238	FFC9	53		
1238	FFCA	F1		
1239	FFCB	C0		
1239	FFCC	14		
1239	FFCD	20		
1239	FFCE	36		
1239	FFCF	44		
1239	FFD0	46		
1239	FFD1	65		
1239	FFD2	20		
1239	FFD3	F1		
1240	FFD4	40		
1240	FFD5	19		
1240	FFD6	99		
1240	FFD7	26		
1240	FFD8	19		
1240	FFD9	39		
1240	FFDA	16		
1240	FFDB	60		
1240	FFDC	F1		
1241	FFDD	C0		
1241	FFDE	33		
1241	FFDF	33		
1241	FFE0	30		
1241	FFE1	73		
1241	FFE2	34		
1241	FFE3	50		
1241	FFE4	50		
1241	FFE5	F1		
1242	FFE6	40		
1242	FFE7	99		
1242	FFE8	99		
1242	FFE9	99		
1242	FFEA	98		
1242	FFEB	47		
1242	FFEC	65		
1242	FFED	70		
1242	FFEE	F1		
1243			.	END

END OF MOS/TECHNOLOGY 650X ASSEMBLY VERSION 4
 NUMBER OF ERRORS = 0, NUMBER OF WARNINGS = 0

SYMBOL TABLE

SYMBOL	VALUE	LINE	DEFINED	CROSS-REFERENCES			
ADD	F808	74	510	542	993		
ADD1	F822	98	187				
ADD10	F88B	200	196				
ADD11	F887	163	149				
ADD12	F88U	157	155				
ADD120	F87U	156	141	152	165		
ADD121	F883	161	217				
ADD13	F8C4	206	136				
ADD15	F8D8	216	229	241	247		
ADD17	F8E0	223	***				
ADD18	F8F0	233	215				
ADD19	F8F3	234	***				
ADD2	F827	100	89				
ADD20	F8F4	239	226				
ADD3	F82A	101	91	99	188		
ADD31	F832	105	103				
ADD4	F838	109	199	202			
ADD5	F84D	129	108	116			
ADD6	F892	174	93				
ADD7	F899	182	186				
ADD8	F89C	183	176				
ADD81	F8A9	188	185				
ADD9	F8AC	194	104				
AM8	FC73	769	774				
APB	FC5F	756	762				
ARGXH	0007	35	1033	1121			
ARGXL	0006	34	1031	1119			
ARGYH	0009	37	1050	1139			
ARGYL	0008	36	1048	1137			
ATANX	F878	601	***				
AZTST	FC99	795	214				
AZTSTL	FC9C	797	800				
AZJ	FBBC	641	643				
CHOP	FB36	558	591				
LHOPD0	FB3A	560	562				
CLEAK	FB88	639	81	255	408		
CLRX	FD71	926	***				
CLRXG	FD75	928	930				
CLRY	FD7C	935	512	533	538	956	
CLRYC	FD80	937	939				
CLRZ	FD87	944	233	272	352	499	1039
CLRZU	FD88	946	948				1056
CNT	0003	30	109	257	455	456	474
			665	666	667	668	669
LNTA	0004	32	175	310	313	368	670
COMPXY	FC82	779	174	420			671
CUM1	FC8A	784	787				677
CUM2	FC96	790	788				686
DECHEX	FBC3	650	114				
DEG	0005	33	551	574	593	608	996
DHCNVT	FEE6	673	***				
DHCNV1	FBC7	653	658				

SYMBOL	VALUE	LINE	DEFINED	CROSS-REFERENCES							
MLT	FA53	450	290								
MLTPLY	F908	251	519	549	557	590	606	990			
MLTO	FA59	453	462								
MLT1	FA60	456	459								
MLT2	FA6A	460	457								
MN	0034	22	886								
MOVR	FD65	917	897								
MOVRO	FD67	918	921								
MULT1	F91F	272	263								
MULT2	F923	277	345	377	389						
MULT3	F927	282	268								
MULT4	F930	296	438								
MVFN	FD28	886	****								
MVFX	FD1C	880	509	994							
MVFY	FD20	882	505								
MVHZ	FD24	884	****								
MVNFM	FD38	894	****								
MVNFX	FD2C	888	506	537	544						
MVNFX	FD30	890	555								
MVNFX	FD34	892	****								
MVTR	FD3A	895	857	859	861	863	865	867	869	871	873
			877	879	881	883	885	887	889	891	893
MVTRO	FD3U	897	900								
MVXM	FCF4	860	****								
MVXN	FCF8	862	498	532	588	604					
MVXY	FCFL	856	589	605							
MVXZ	FCFJ	858	572								
MVYM	F004	868	****								
MVYN	F008	870	****								
MVYX	FCFC	864	****								
MVYZ	FD00	866	****								
MVZM	FD14	876	504	520	987						
MVZN	FD18	878	541	546							
MVZX	FD0C	872	511	547	556	988	992	1041	1128		
MVZY	FD10	874	508	543	548	578	995	1058	1146		
MX	0030	19	880								
MY	0031	20	882								
MZ	0032	21	884								
N	0000	27	115	450	486	530	570	586	602	618	627
			692	705	715	725	735	744	754	767	781
			805	815	836	1188					
NKUN	0001	28	497	521	553	576	958	980			
NM	0043	26	894								
NX	0040	23	888								
NY	0041	24	890								
NZ	0042	25	892								
OVERR	0014	46	847								
PGTARG	FDE1	1003	1040	1057							
PGTRG2	FDEA	1007	1022								
PGTRG1	FED4	1023	1009								
PLUADX	FEOA	1031	****								
PLOADY	FE23	1048	****								
POLY	FDC1	987	554	577							
POLYO	FDCA	990	997								
PREC	0110	42	1035	1052	1080	1123	1141	1158	1186		
PSTRES	FE3C	1065	****								
PTK	000C	39	1005	1010	1023	1032	1036	1049	1051	1091	1096
			1107	1120	1122	1138	1140				

SYMBOL	VALUE	LINE	DEFINED	CROSS-REFERENCES
ULLUADY	FEA2	1137	****	
USTRES	FEBA	1154	****	
USTRS0	FEC1	1157	1164	
USTRS1	FEC8	1161	1159	
USTRS2	FED1	1165	1160	
XM	0003	9	860	
XN	0004	10	862	
XSY	FCBF	824	100	182
XSY1	FCC1	825	831	
XY	0001	7	856	
XZ	0002	8	858	
XZTST	FCA6	805	85	262 403
XZTST0	FCA8	806	809	
XZTST1	FCB2	811	798	807 817
XZTST2	FCB0	810	801	820
YM	0013	13	868	
YN	0014	14	870	
YY	0010	11	864	
YZ	0012	12	866	
YZTST	FCB3	815	90	267 398
YZTST0	FCB5	816	819	
ZM	0023	17	876	
ZN	0024	18	878	
ZX	0020	19	872	
ZY	0021	16	874	

INSTRUCTION COUNT

ADC	13
AND	14
ASL	12
BCC	11
BCS	6
BEQ	24
BIT	13
BMI	5
BNE	32
BPL	26
BRK	0
BVC	6
BVS	6
CLC	5
CLD	1
CLI	0
CLV	0
CMP	6
CPX	2
CPY	3
DEC	6
DEX	300
DEY	5
EOR	5
INC	5
INX	9
INY	10
JMP	28
JSR	114
LDA	159
LDX	31
LDY	7
LSR	18
NOP	0
ORA	9
PHA	4
PHP	0
PLA	4
PLP	0
ROL	0
RTI	0
RTS	48
SBC	11
SEC	10
SED	6
SEI	6
STA	118
STX	22
STY	22
TAX	2
TAY	2
TSX	00
TXA	00
TXS	00
TYA	1

= SYMBOLS = 238 (LIMIT = 800) = BYTES = 2032 (LIMIT = 4096)
= LINES = 1449 (LIMIT = 4000) = XREFS = 743 (LIMIT = 1800)

Notizen

Notizen

ALLES ÜBER MOS -

MICRO'S!

VOM DATENBLATT
ÜBER HANDBÜCHER
BIS ZUM 4 - FARB - POSTER

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KIM-1 HANDBUCH

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PROGRAMMER-HANDBUCH

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KIM-1
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USER MANUAL

KIM
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KIM-2/3/4 USER MANUAL
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MEMORY-EJECTA KEYBOARD - MOTHERBOARD



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microcomputer
datensysteme gmbh

Luikenplatz 4 · d 61 darmstadt ☎ 0 61 51/23959

MICROCOMPUTER-HANDBÜCHER (Netto + MWSt. + Versand)

Wir bestellen zu Ihren bekannten Bedingungen:

1.10 KIM-1 User Manual	englisch	DM	19,80	Stck:	DM:
1.11 KIM-1 Handbuch	deutsch	DM	19,80	Stck:	DM:
1.20 KIM-2/3/4 User Manual Expansion Modules Memory/Errata Sheet/Motherboard	englisch	DM	9,25	Stck:	DM:
1.30 KIMath Subroutines Programming Manual	englisch	DM	15,45	Stck:	DM:
1.40 KIM Text Editor User Manual	englisch	DM	8,90	Stck:	DM:
1.50 KIM Assembler Manual Preliminary	englisch	DM	10,20	Stck:	DM:
1.60 Programming Manual	englisch	DM	28,60	Stck:	DM:
1.61 Programmierhandbuch	deutsch	DM	28,60	Stck:	DM:
1.62 THE FIRST BOOK OF KIM-1 DAS KIM-REZEPTBUCH	englisch	DM	19,80	Stck:	DM:
1.70 Hardware Manual	englisch	DM	24,90	Stck:	DM:
1.71 Hardware Handbuch	deutsch	DM	24,90	Stck:	DM:
1.80 TIM Terminal Interface Monitor Manual	englisch	DM	12,90	Stck:	DM:
1.90 Cross Assembler Manual Preliminary	englisch	DM	12,70	Stck:	DM:
1.99 Bücherliste über weitere ca. 60 Bücher mit Kurzbeschreibung		DM	--,--	Stck:	DM:

MICROPROCESSOR DATENBLÄTTER (Netto + MWSt. + Versand)

2.10 Spectrum of Products	englisch	DM	0,95	Stck:	DM:
2.20 MCS 6500 Microprocessors	englisch	DM	0,95	Stck:	DM:
2.30 MCS 6520 Microprocessors	englisch	DM	0,95	Stck:	DM:
2.31 MCS 6522 Microprocessors	englisch	DM	0,95	Stck:	DM:
2.40 MCS 6530 Microprocessors	englisch	DM	0,95	Stck:	DM:
2.60 MCS 6500 X Instruction Set 6502-6515	englisch	DM	1,10	Stck:	DM:
2.80 MCS 6532 Microprocessors	englisch	DM	0,95	Stck:	DM:

MICROCOMPUTER-LOSEBLATT-BIBLIOTHEK (Netto + MWSt. + Versand)

3.10 Loseblattsammelordner aus Kunstleder für alle Bücher und Datenblätter		DM	4,95	Stck:	DM:
3.20 MICRO-Info Hardware ca. 1400 Seiten (FMI)	deutsch	DM	42,00*	Stck:	DM:
3.23 MICRO-Info Software ca. 400 Seiten (FMI)	deutsch	DM	38,00*	Stck:	DM:
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