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KIM-1 USER MANUAL

KIM-1
MICROCOMPUTER MODULE
USER MANUAL

JANUARY 1976

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CHAPTER 1

YOUR KIM-1 MICROCOMPUTER MODULE

Congratulations and welcome to the exciting new world of microcomputers! As the owner of a KIM-1 Microcomputer Module, you now have at your disposal a completely operational, fully tested, and very capable digital computer which incorporates the latest in microprocessor technology offered by MOS Technology, Inc. By selecting the KIM-1 module, you have eliminated all of the problems of constructing and debugging a microcomputer system. Your time is now available for learning the operation of the system and beginning immediately to apply it to your specific areas of interest. In fact, if you will follow a few simple procedures outlined in this manual, you should be able to achieve initial operation of your KIM-1 module within a few minutes after unpacking the shipping container.

Your KIM-1 module has been designed to provide you with a choice of operating features. You may choose to operate the system using only the keyboard and display included as part of the module. Next, you may add a low cost audio cassette tape recorder to allow storage and retrieval of your programs. Also, you may add a serial interfaced teleprinter to the system to provide keyboard commands, hard-copy printing, and paper tape read or punch capability.

At the heart of your KIM-1 system is an MCS 6502 Microprocessor Array operating in conjunction with two MCS 6530 arrays. Each MCS 6530 provides a total of 1024 bytes of Read-only Memory (ROM), 64 bytes of Random Access Memory (RAM), 15 Input/Output pins, and an Interval Timer. Stored permanently in the ROM's of the MCS 6530 arrays are the monitor and executive programs devised by MOS Technology, Inc. to control the various operating modes of the KIM-1 system.

The KIM-1 system is intended to provide you with a capable micro-computer for use in your "real-world" application. Accordingly, the system includes a full 1024 bytes of RAM to provide data and program storage for your application program. In addition, you are provided with 15 bidirectional input/output pins to allow interface control of your specific application. Finally, one of the interval timers included in the system is available for generation of time base signals required by your application.

Your KIM-1 system comes to you complete with all components mounted and tested as a system. You need not worry about timing signals (we've included a 1MHz crystal oscillator on the module), interface logic or levels between system components, or interface circuitry to peripheral devices. In fact, you need only apply the indicated power supply voltages to the designated pins to achieve full operation of your KIM-1 system.

We recommend that you read all of this manual before applying power to or attempting to operate your KIM-1 module. In the order presented, you will find:

- Chapter 2 - "hints and kinks" to help you achieve initial system operation
- Chapter 3 - a more detailed description of the KIM-1 system hardware and software
- Chapter 4 - operating procedures for all system modes
- Chapter 5 - an example of a typical application program using all of the features of the KIM-1 system.

At some future time, you may find it desirable to expand the KIM-1 system to incorporate more memory, different types of memory, or additional input/output capability. Again, we have tried to make system expansion as simple as possible with all required interface signals brought out to a special connector on the module. Watch for:

Chapter 6 - a guide to system expansion for increasing both memory and input/output capability

Despite our best efforts to provide you with a fully operable and reliable system, you might encounter some difficulties with your KIM-1 module. If so, refer to:

Chapter 7 - some guidance on warranty and service procedures for your KIM-1 module

Following the basic text of this manual, you will find a series of Appendices intended to provide you with detailed information on certain specialized subjects of interest to you in understanding the operation of the KIM-1 system.

Lastly, since this manual cannot presume to provide all of the technical information on the hardware or programming aspects of the MCS 6502 microprocessor array, we are including with your KIM-1 system two additional manuals for your reference. The Hardware Manual defines the various elements of the system, their electrical and interface characteristics, and the basic system architecture and timing. The Programming Manual provides the detailed information required to write effective programs using the MCS 6502 instruction program set.

So much for introductory comments! Now lets get started and see if we can get your KIM-1 Microcomputer Module doing some real work for you.

CHAPTER 2

GETTING STARTED

This chapter is intended to guide you through the first important steps in achieving initial operation of your KIM-1 Microcomputer Module. We will ask you to perform certain operations without explanation at this time as to why they are being done. In later sections of this manual, full explanations will be offered for every operating procedure.

2.1 PARTS COMPLEMENT

After unpacking the shipping container for your KIM-1, you should have located the following items:

- 3 Books - KIM-1 Users Manual
Hardware Manual
Programming Manual
- 1 Programming Card
- 1 System Schematic
- 1 KIM-1 Module
- 1 Connector (Already mounted on the Module)
- 1 Hardware Packet
- 1 Warranty Card

You may wish to save the shipping container and packing material should you need to return your KIM-1 module to us at some future date.

2.2 A FEW WORDS OF CAUTION

WARNING

Your KIM-1 module includes a number of MOS integrated circuits. All such circuits include protective devices to prevent damage resulting from inadvertant application of high voltage potentials to the pins of the device. However, normal precautions should be taken to prevent the application of high voltage static discharges to the pins of an MOS device.

Immediately before removal of the packing material from your KIM-1 module, you should develop the following precautionary habits:

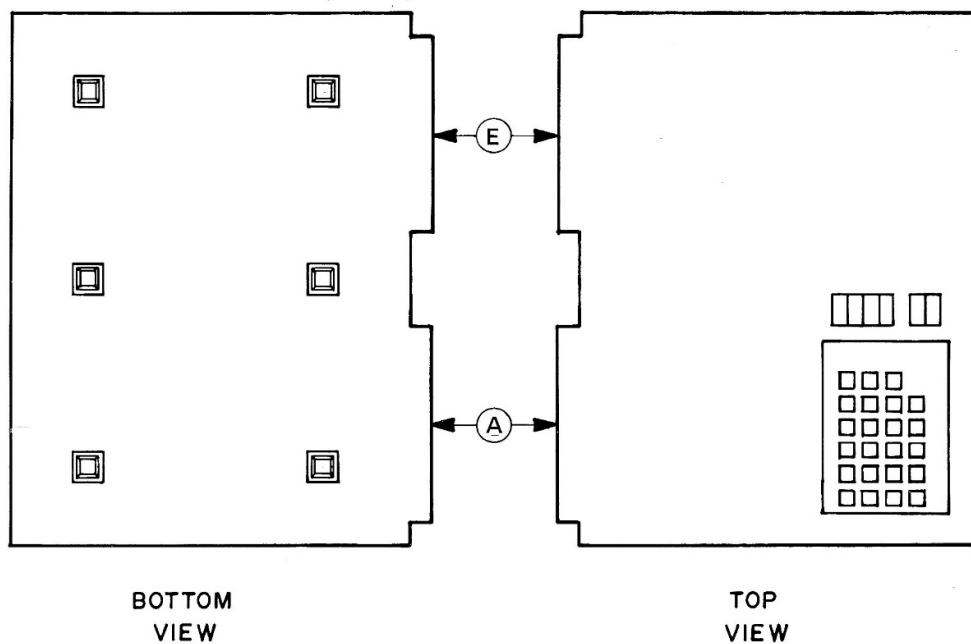
1. Discharge any static charge build up on your body by touching a ground connection before touching any part of your KIM-1 module. (This precaution is especially important if you are working in a carpeted area)
2. Be certain that soldering irons or test equipment used on the KIM-1 module are properly grounded and not the source of dangerously high voltage levels.

On a different subject, after unpacking your KIM-1 module, you will note the presence of a potentiometer. This adjustment has been set at the factory to insure correct operation of the audio cassette interface circuits. It should never be necessary for you to change the position of this potentiometer.

2.3 FIRST STEPS

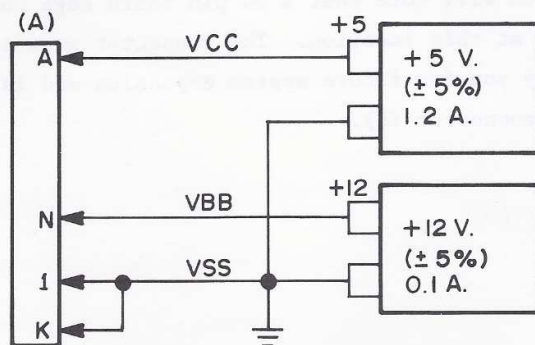
After unpacking the KIM-1 module, locate the small hardware packet and install the rubber pads provided. The rubber pads are located at the bottom of the module (see attached sketch) and act both to lift the card off your work surface and to provide mechanical support for the module while you depress keys.

Place the module such that the keyboard is to your lower right and observe that two connector locations extend from the module to your left. The connector area on the lower left is referred to as the Application connector (A). You will note that a 44 pin board edge connector is already installed at this location. The connector area to the upper left is for use by you for future system expansion and is referred to as the Expansion connector (E).



KIM-1 Module
FIGURE 2.1

Remove the (A) connector from the module and connect the pins as shown in the sketch.



Power Supply Connections
FIGURE 2.2

Reinstall the (A) connector making certain that the orientation is correct.

- Note 1: The +12 volt power supply is required only if you will be using an audio cassette recorder in your system.
- Note 2: The jumper from pin A-K to Vss (Pin A-1) is essential for system operation. If you expand your system later, this jumper will be removed and we'll tell you what to do to pin A-K.
- Note 3: If you don't have the proper power supplies already available, you may wish to construct the low cost version shown with schematic and parts list in Appendix D. In any event, your power supply must be regulated to insure correct system operation and must be capable of supplying the required current levels indicated in the sketch.

Now, recheck your connections, turn on your power supplies, and depress **RS** (reset). You should see the LED display digits light as your first check that the system is operational. If not, recheck your hookup or refer to Appendix C (In Case of Trouble).

2.4 LETS TRY A SIMPLE PROGRAM

Assuming that you have completed successfully all of the steps thus far, a simple program now can be tried to demonstrate the operation of the system and increase your confidence that everything works properly. We'll be using only the keyboard and display on the module for this example. (In the next two sections we'll worry about the teleprinter and the audio cassette).

For our first example, we will add two 8 bit binary numbers together and display the result. We presume that you are familiar with the hexadecimal representation of numbers and the general rules for binary arithmetic.

First check and be sure that the slide switch in the upper right corner of the keyboard is pushed to the left (SST Mode is OFF). Now proceed with the following key sequence:

<u>Press Keys</u>	<u>See On Display</u>	<u>Step #</u>
AD	xxxx xx	1
0 0 0 2	0002 xx	2
DA	0002 xx	3
1 8	0002 18	4
+ A 5	0003 A5	5
+ 0 0	0004 00	6
+ 6 5	0005 65	7
+ 0 1	0006 01	8
+ 8 5	0007 85	9
+ F A	0008 FA	10
+ A 9	0009 A9	11
+ 0 0	000A 00	12
+ 8 5	000B 85	13
+ F B	000C FB	14
+ 4 C	000D 4C	15
+ 4 F	000E 4F	16
+ 1 C	000F 1C	17

What you have just done is entered a program and stored it in the RAM at locations 0002 through 000F. You should have noticed the purpose of several special keys on your keyboard:

- AD - selects the address entry mode
- DA - selects the data entry mode
- + - increments the address without changing the entry mode
- 0 to F - 16 entry keys defining the hex code for address or data entry

You've noticed as well that your display contains 6 digits. The four on the left are used to display the hex code for an address. The two on the right show the hex code for the data stored at the address shown. Therefore, when you pressed AD (step 1) and 0 0 0 2 (step 2), you defined the address entry mode, selected the address 0002, and displayed the address 0002 in the four left-most display digits. Incidentally, when we show an "x" in the display chart, we mean that we don't know what will be displayed and we "don't care."

Next you pressed DA (step 3) followed by 1 8 (step 4). Here, you have defined the data entry mode and entered the value 18 to be stored at your selected address 0002. Of course, the 18 then was displayed in the two right-most digits of your display.

You remained in the data entry mode but began to press + followed by a two digit number (steps 5 to 17). Note that each depression of the + key caused the address displayed to increase by one. The hex keys following the + key continued to enter the data field of the display. This procedure is merely a convenience when a number of successive address locations are to be filled.

If you made any mistakes in pressing the keys, you should have noticed that correcting an error is simply a matter of reentering the data until the correct numbers show on the display.

The program you have entered is a simple loop to add two 8 bit binary numbers together and present the result on the display. For a programmer, the listing of the program entered might appear as follows:

```

POINTL      = $FA
POINTH      = $FB
START       = $1C4F
0000        VAL1
0001        VAL2
0002  18          PROG    CLC
0003  A5 00          LDA   VAL1
0005  65 01          ADC   VAL2
0007  85 FA          STA   POINTL
0009  A9 00          LDA   #000
000B  85 FB          STA   POINTH
000D  4C 4F 1C       JMP   START

```

Stated in simple terms, the program will clear the carry flag (CLC), load VAL1 into the accumulator (LDA VAL1), add with carry VAL2 to the accumulator (ADC VAL2), and store the result in a location POINTL (STA POINTL). A zero value is stored in a location POINTH (LDA #000 and STA POINTH) and the program jumps to a point labelled START (JMP START). This pre-stored program will cause the display to be activated and will cause the address field of your display to show the numbers stored in locations POINTH and POINTL. Note that the result of the addition has already been stored in location POINTL.

The hex codes appearing next to the address field of the listing are exactly the numbers you entered to store the program. We refer to these as machine language codes. For example, 4C is the hex code for the JMP instruction of the microprocessor. The next two bytes of the program define 1C4F (START) as the jump address.

As yet, you are not able to run the program because you have not yet entered the two variables (VAL1 and VAL2). Lets try an actual example:

<u>Press Keys</u>	<u>See On Display</u>	<u>Step #</u>
AD	000F 1C	18
0 0 0 0	0000 xx	19
DA	0000 xx	20
0 2	0000 02	21
+ 0 3	0001 03	22
GO	0005 xx	23

Steps 18 to 21 store the hex value 02 in location 0000 (VAL1). Step 22 stores the hex value 03 in location 0001 (VAL2). Now we are ready to run the program. In step 23, the GO key causes the program to execute and the result, 05, appears in the right two digits of the address display. Although the problem appears trivial, it illustrates the basic principles of entering and executing any program as well as providing a fairly high assurance level that your KIM-1 module is operating properly.

You should try one more example using your stored program. Repeat steps 18 to 23 but substitute the value FF for VAL1 and VAL2 at locations 0000 and 0001. Now when you press the GO key, your display should read:

00FE xx

The answer is correct because:

$$\begin{array}{r}
 \text{FF} = 1111\ 1111 \\
 + \text{FF} = \underline{1111\ 1111} \\
 \text{FE} \quad 1111\ 1110
 \end{array}$$

Try some more examples if you wish and then let's move on to the rest of the system.

2.5 ADDING A TAPE RECORDER

In the previous section, you entered and executed a program. If you turn off the power supplies to the system, your program is lost since the memory into which you stored your program is volatile. If you require the same program again, you would have to repower the system and reenter the program as in the previous example.

The KIM-1 system is designed to work with an audio cassette tape recorder/player to provide you with a medium for permanent storage of your programs or data. The cassette with recorded data may be reread by the system as often as you wish. In this section, you will connect the audio cassette unit to the system and verify its operation.

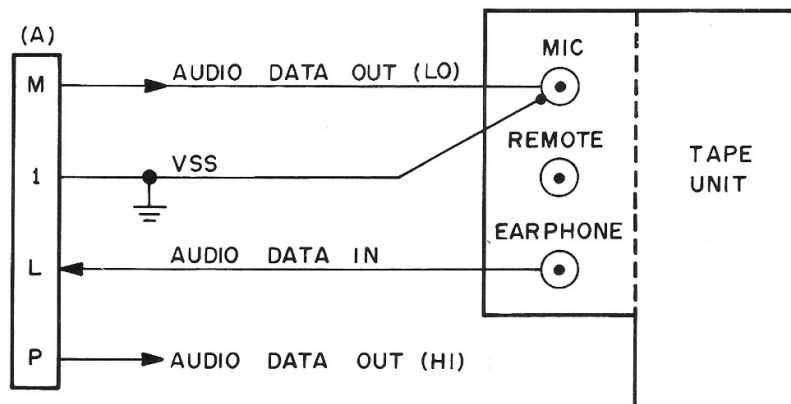
The recording technique used by the KIM-1 system and the interface circuits provided have been selected to insure trouble-free operation with virtually any type and any quality level audio cassette unit. (We have demonstrated correct operation with a tape unit purchased for less than \$20.00 from a local discount outlet). In addition, tapes recorded on one unit may be played back to the system on a different unit if desired. We recommend, of course, that you make use of the best equipment and best quality tapes you have available.

In selecting a tape unit for use with your KIM-1 system, you should verify that it comes equipped with the following features:

1. An earphone jack to provide a source of recorded tape data to the KIM-1 system.
2. A microphone jack to allow recording of data from the KIM-1 system on the tape.
3. Standard controls for Play, Record, Rewind, and Stop.

Note: You should avoid certain miniaturized tape equipment intended for dictating applications where the microphone and speaker are enclosed within the unit and no connections are provided to external jacks. If such equipment is used, you will have to make internal modifications to reach the desired connection points.

To connect your tape unit to the KIM-1 module, turn off the power supplies and remove the connector (A) from the module. Add the wires shown in the sketch:



*Audio Tape Unit Connections
FIGURE 2.3*

Keep the leads as short as possible and avoid running the leads near sources of electrical interference. The connections shown are for typical, portable type units. The Audio Data Out (LO) signal has a level of approximately 15 mv. (peak) at pin M. Should you desire to use more expensive and elaborate audio tape equipment, you may prefer to connect the high level (1 volt peak) audio signal available at pin P to the "LINE" input of your equipment.

Return the connector (A) to its correct position on the KIM-1 module and turn on the power supplies. To verify the operation of your audio cassette equipment, try the following procedures:

1. Reenter the sample program following the procedures outlined in the previous section (2.4). Try the sample problem again to be sure the system is working correctly.
2. Install a cassette in your tape equipment and REWIND to the limit position.
3. Define the starting and ending address of the program to be stored and assign an identification number (ID) to the program:

Press Keys	See On Display	Step #
AD	xxxx xx	1
1 7 F 4	17F4 xx	2
DA 0 0	17F4 00	3
+ 0 0	17F5 00	4
+ 1 0	17F6 10	5
+ 0 0	17F7 00	6
+ 0 1	17F8 01	7
AD	17F8 01	8
1 8 0 0	1800 xx	9

You will recall that the program we wish to store on tape was loaded into locations 0000 to 000F of the memory. Therefore, we define a starting address for recording as 0000 and store this in locations 17F4 and 17F5 (Steps 1 to 4). We define an ending address for recording as one more than the last step of our program and stored the value 0010 (= 000F + 1) in locations 17F6 and 17F7 (Steps 5, 6). Finally we pick an arbitrary ID as 01 and store this value at location 17F8 (Step 7).

The starting address of the tape recording program is 1800. In Steps 8 and 9 we set this address value into the system. If we were to press **GO**, the system would proceed to load data on to the magnetic tape. But first, we'd better start the tape!

4. Select the Record/Play mode of the tape recorder. Wait a few seconds for the tape to start moving and now:

Press **GO**

5. The display will go dark for a short time and then will relight showing:

0000 xx

6. As soon as the display relights, the recording is finished and you should STOP the tape recorder.

Now, you should verify that the recording has taken place correctly. This can be proven by reading the tape you have just recorded. Proceed as follows:

1. Rewind the tape cassette to its starting position.
2. Turn off the system power supplies and then later, turn them back on.

This has the effect of destroying your previously stored program which you already have recorded on tape.

3. Prepare the system for reading the tape as follows:

<u>Press Keys</u>	<u>See On Display</u>	<u>Step #</u>
AD	xxxx xx	1
1 7 F 8	17F8 xx	2
DA	17F8 xx	3
0 1	17F8 01	4
AD	17F8 01	5
1 8 7 3	1873 xx	6
GO	(Dark)	7

The KIM-1 system is now looking for tape input data with the ID label 01. Recall that this is the same ID label we assigned when we recorded the program.

4. If your tape unit has a volume control, set the control at approximately the half way point.
5. Now, turn on the tape using the PLAY mode. The tape will move forward and the system will accept the recorded data. As soon as the data record (ID=01) has been read, the display should relight showing:

0000 xx

You may now stop the tape unit. If the display relights and shows;

FFFF xx

this means that the selected record has been located and read but that an error has occurred during the reading of the data. In this case, press the RS key and repeat the read tape procedures from the beginning. If the FFFF still shows on the display, repeat the entire recording and playback procedures checking each step carefully. If the problem persists, refer to Appendix C, (In Case of Trouble).

If the tape continues to run and the display does not relight, this means that the system has been unsuccessful in reading any data back from the tape. In this case, repeat the entire recording and playback procedures checking each step carefully. If the problem persists, refer to Appendix C, (In Case of Trouble).

6. Assuming that you have read the tape successfully, you now may verify that the program has been restored to memory by trying a sample problem. ($02 + 03 = 05$)

NOTE: The KIM-1 interface circuits for the audio tape system are designed so that you do not require special test equipment to set up correct operating levels. If you have followed the procedures indicated, the tape system should work without the need of any adjustments by you.

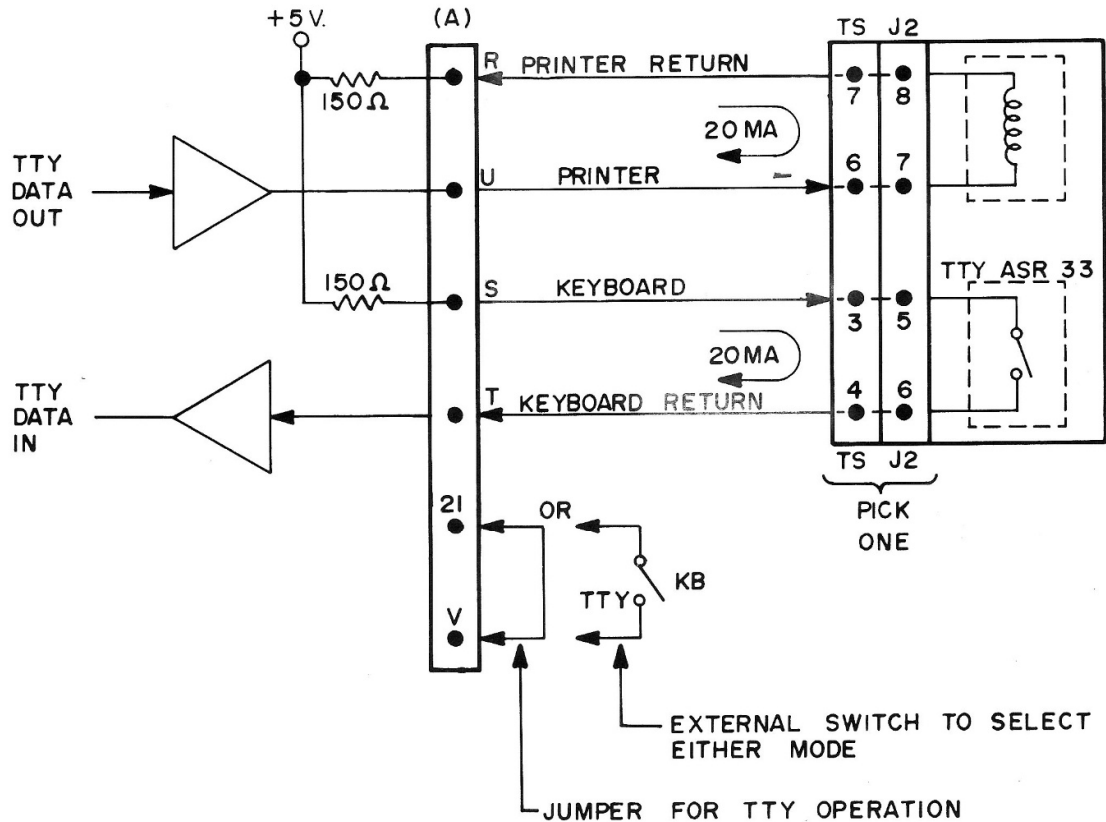
2.6 ADDING A TELEPRINTER

If you have access to a serial teleprinter, you may add such a unit to the KIM-1 system with very little effort. One of the more commonly available units of this type is the Teletype Model 33ASR which we will use for the purposes of illustration in this section. However, if you have available different equipment, you may use the information presented here as a guide to connecting your specific unit. In any case, we recommend you follow the directions offered by the equipment manufacturer in his instruction manual to effect the desired wiring and connection options.

The KIM-1 provides for a 4 wire interface to the TTY. Specifically, the "20 MA loop" configuration should be used and you should check that your TTY has been wired for this configuration. If not, you may easily change from "60 MA loop" to "20 MA loop" configurations following the manufacturers directions. You need not concern yourself with full or half-duplex configurations since the KIM-1 system will work with either. Also, you are not restricted to units with specific bit rates (10 CPS for TTY) since the KIM-1 system automatically adjusts for a wide variety of data rates (10CPS, 15CPS, 30CPS, etc.).

To connect the TTY to the system, proceed as follows:

1. Turn off system power and remove connector (A) from the module.
2. Add the wires shown in the sketch to connector (A) and to the appropriate connector on the TTY unit.



TTY Connections
FIGURE 2.4

3. The jumper wire from A-2I to A-V is used to define for the KIM-1 system that a teleprinter will be used as the only input/display device for the system. If you expect to use both TTY and the KIM-1 keyboard/display, you should install the switch shown instead of the jumper. Now, the switch, when open, will allow use of the keyboard and display on the KIM-1 module and, when closed, will select the teleprinter as the input/display device. (Of course, you may use a clip-lead instead of the switch if you desire).
4. Be sure pins A-2I and A-V are connected. Reinstall connector (A) and return power to the system. Turn-on the TTY.
5. Press the RS key on the KIM-1 module then press the RUB
OUT key on the TTY. This step is most important since the KIM-1 system adjusts automatically to the bit rate of the serial teleprinter and requires this first key depression to establish this rate.

If everything is working properly you should immediately observe a message being typed as follows:

KIM

This is a prompting message telling you that the TTY is on-line and the KIM-1 system is ready to accept commands from the TTY keyboard.

Should the prompting message not be typed press the RS key on the KIM-1 keyboard and then the RUB
OUT key on the TTY. If the "KIM" message still is not typed, recheck all connections and the TTY itself and try again. If the problem persists, refer to Appendix C, (In Case of Trouble).

6. Assuming that the TTY is operable, you may now try a simple group of operations to verify correct system operation:

<u>Press Keys</u>	<u>See Printed</u>	<u>Step #</u>
	KIM	
0 0 0 2	xxxx xx	1
SPACE	0002	2
1 8 ●	0002 xx	3
	18.	4
A 5 ●	0003 xx	5
	A5.	6
LF	0004 xx	7
RUB OUT	0003 A5	8
	KIM	
	xxxx xx	9

Step 1 shows the "KIM" prompting message. In Step 2, an address (0002) is selected followed by a space key in Step 3. The address cell 0002 together with the data stored at that location (xx) is printed. Step 4 shows the "modify cell" operation using the ● key and the hex data keys preceding. Step 5 shows the incrementing to the next address cell (0003) after the ● key. Note that the modification of cell 0002 also occurs. Steps 6 and 7 show the modification of data in cell 0003 and the incrementing to cell 0004. Step 8 shows the action of the LF key in backing up one cell to 0003 where we can see from the printout that the correct data (A5) has been stored at that location. Step 9 shows the reaction to the RUB
OUT key in resetting the system and producing a new "KIM" prompting message. Note, by the way, that in this example you have repeated a portion of the program entry exactly as you did in Section 2.4 but this time using the TTY.

So much for now! If all of the operations have occurred properly, you may be certain that your TTY and KIM-1 module are working together correctly. We will describe in detail all of the other operations possible with the TTY in a later section of the manual.

If you have reached this point without problems, you now have completed all of the required system tests and may be confident that the KIM-1 module and your peripheral units are all working correctly. Our next task is to learn more about the KIM-1 system and its operating programs.

CHAPTER 3

THE KIM-1 SYSTEM

Up to this point you have been engaged in bringing up your KIM-1 system and verifying its correct operation. Now it's time to learn more about the various parts of the KIM-1, how the parts work together as a system, and how the operating programs control the various activities of the system. The diagrams included in this section together with your full sized system schematic will be helpful in understanding the elements of your KIM-1 module.

3.1 KIM-1 SYSTEM DESCRIPTION

Figure 3-1 shows a complete block diagram of the KIM-1 system. You should note first the presence of the MCS 6502 Microprocessor Array which acts as the central control element for the system. This unit is an 8 bit microprocessor which communicates with other system elements on three separate buses. First, a 16 bit address bus permits the 6502 to address directly up to 65,536 memory locations in the system. Next, an 8 bit, bidirectional data bus carries data from the 6502 array to any memory location or from any memory location back to the 6502 array. Lastly, a control bus carries various timing and control signals between the 6502 array and other system elements.

Associated with the 6502 array is a 1 MHz crystal which operates with an oscillator circuit contained on the 6502 array. This crystal controlled oscillator is the basic timing source from which all other system timing signals are derived. In particular, the ϕ_2 signal generated by the 6502 array and used either alone, or gated with other control signals, is used as the system time base by all other system elements.

The 6502 microprocessor is structured to work in conjunction with various types of memory. In the KIM-1 system, all memory may be considered to be of the Read-only (ROM) or Read/Write (RAM) variety. The ROM portion of the memory provides permanent storage for the operating programs essential to the control of the KIM-1 system. You will note the inclusion of two devices, labelled 6530-002 and 6530-003. Each of these devices include a 1024 byte (8 bits per byte) ROM with different portions of the operating program stored permanently in each ROM.

RAM type memory is available at three locations in the system. Again, each of the 6530 arrays include 64 bytes of RAM primarily used for temporary data storage in support of the operating program. In addition, a separate 1024 byte RAM is included in the KIM-1 system and provides memory storage for user defined application programs and data.

Input/output controls for the system also are included within the 6530 arrays. Each 6530 array provides 15 I/O pins with the microprocessor and operating program defining whether each pin is an input pin or output pin, what data is to appear on the output pins, and reading the data appearing on input pins. The I/O pins provided on the 6530-002 are dedicated to interfacing with specific elements of the KIM-1 system including the keyboard, display, TTY interface circuit, and audio tape interface circuit. The 15 I/O pins on the 6530-003 are brought to a connector and are available for the user to control a specific application.

Finally, each 6530 array includes an interval timer capable of counting a specific number of system clocks to generate precise timing gates. The exact time interval is preset under program control. The interval timer on the 6530-003 array is available for a user defined application program and is not required by the operating programs.

Figure 3-1 shows a major block labelled Control Logic. Included under this category are an address decoder used for generation of chip select signals for the 6530 arrays and the static RAM. Also included is the logic required to debounce the keys for system reset (RS key) and program stop (ST key). Lastly, special logic is included to allow operation of the system in a "single instruction" mode to facilitate program debugging.

Figure 3-1 shows the keyboard/display logic interfacing with the I/O pins of the 6530-002. Also shown are the interface circuits for transmission of data to and reception of data from the TTY and audio tape units.

Figure 3-2 shows the detailed interconnections between the MCS 6502 and the two MCS 6530 arrays.

Figure 3-3 shows detailed logic and schematics for the control logic.

Figure 3-4 shows a detailed schematic of the static RAM.

Figure 3-5 and 3-6 show the detailed schematic of the keyboard and display logic and circuits.

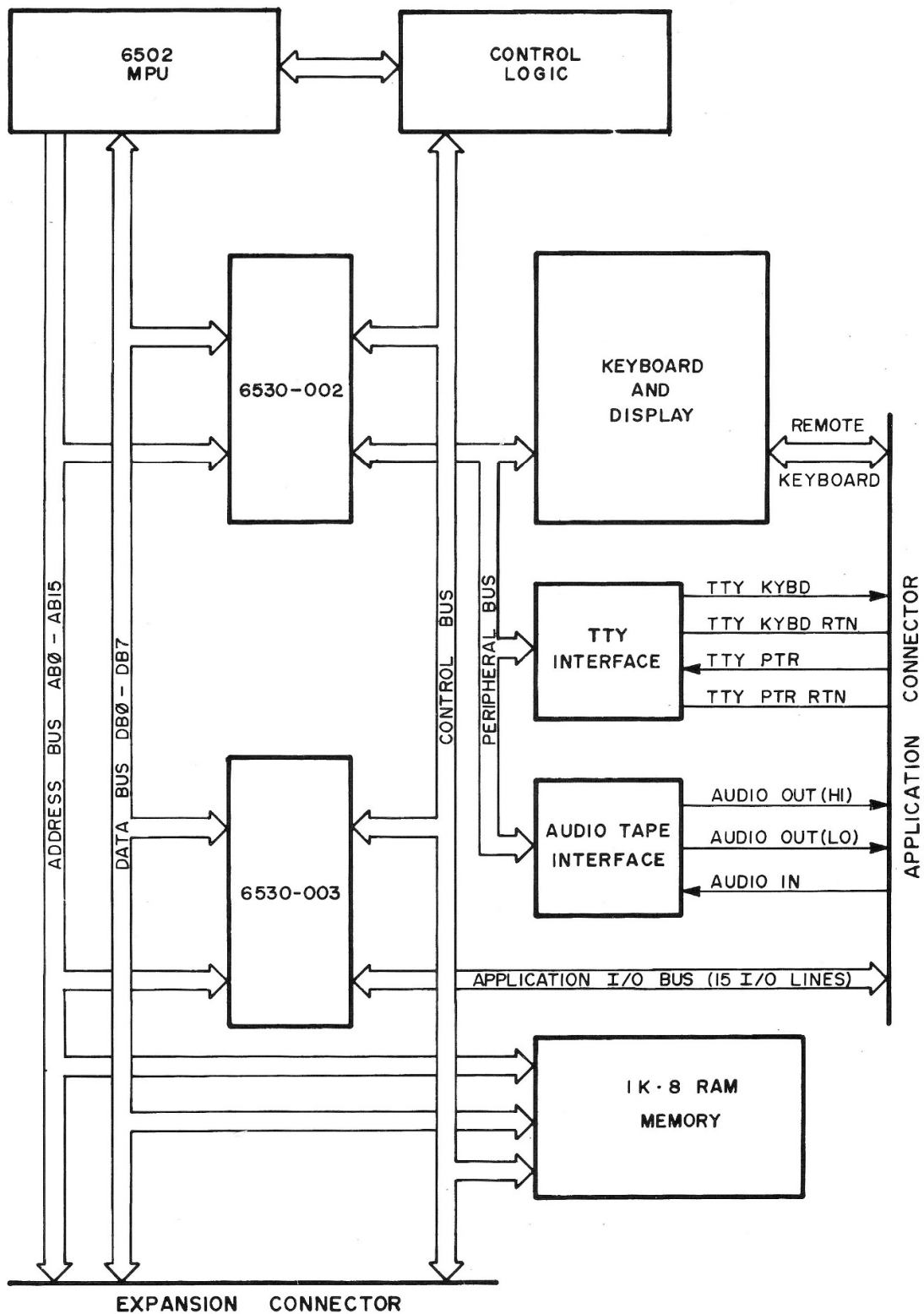
Figure 3-7 details the schematic of the TTY interface circuits.

Figure 3-8 details the schematic of the audio tape cassette interface circuits.

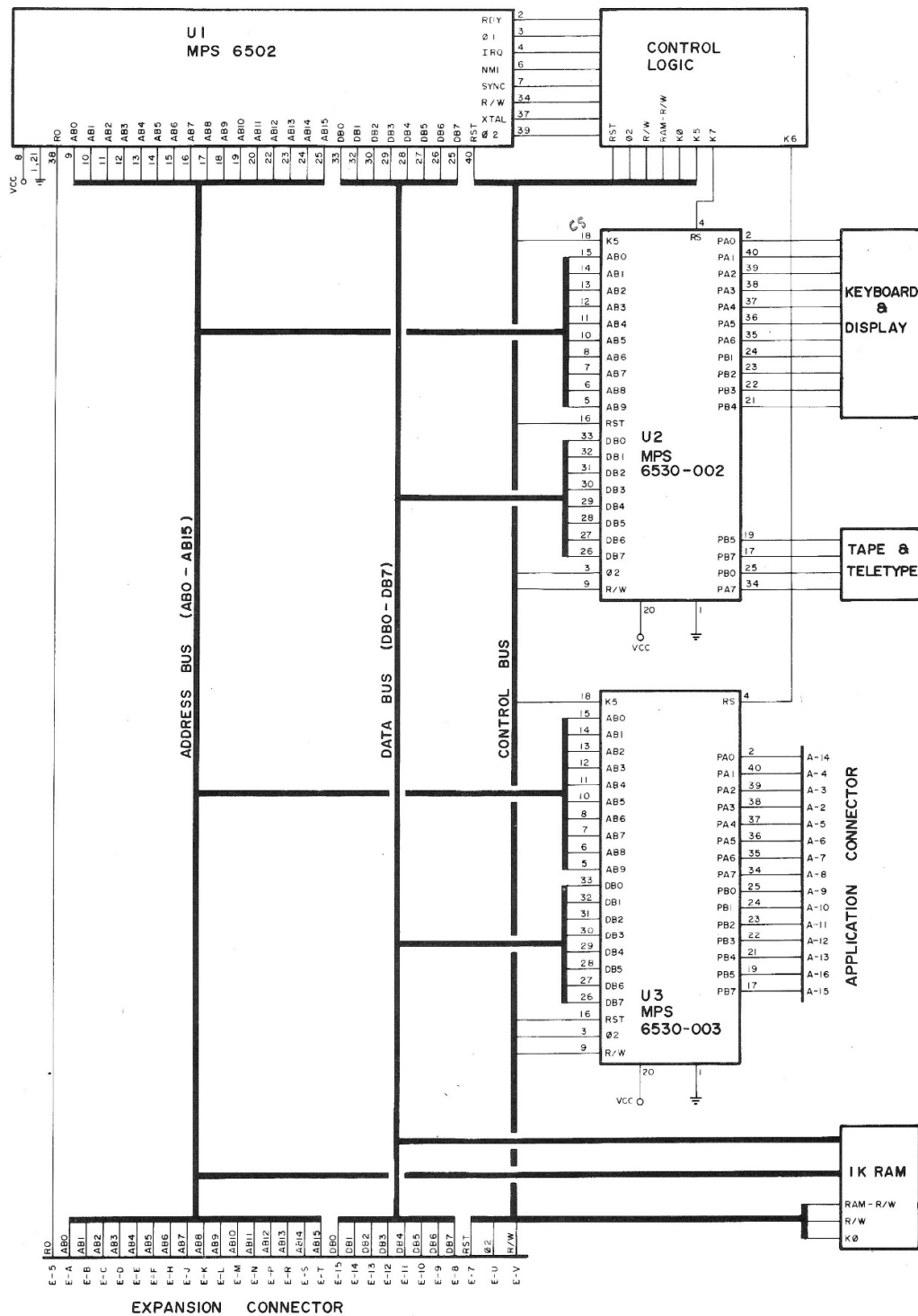
Figures 3-9 and 3-10 provide a summary of all signals available on either the Application connector or the Expansion Connector.

The fold-out system schematic shows all of the elements of the system connected together and all signals appearing on the module connectors.

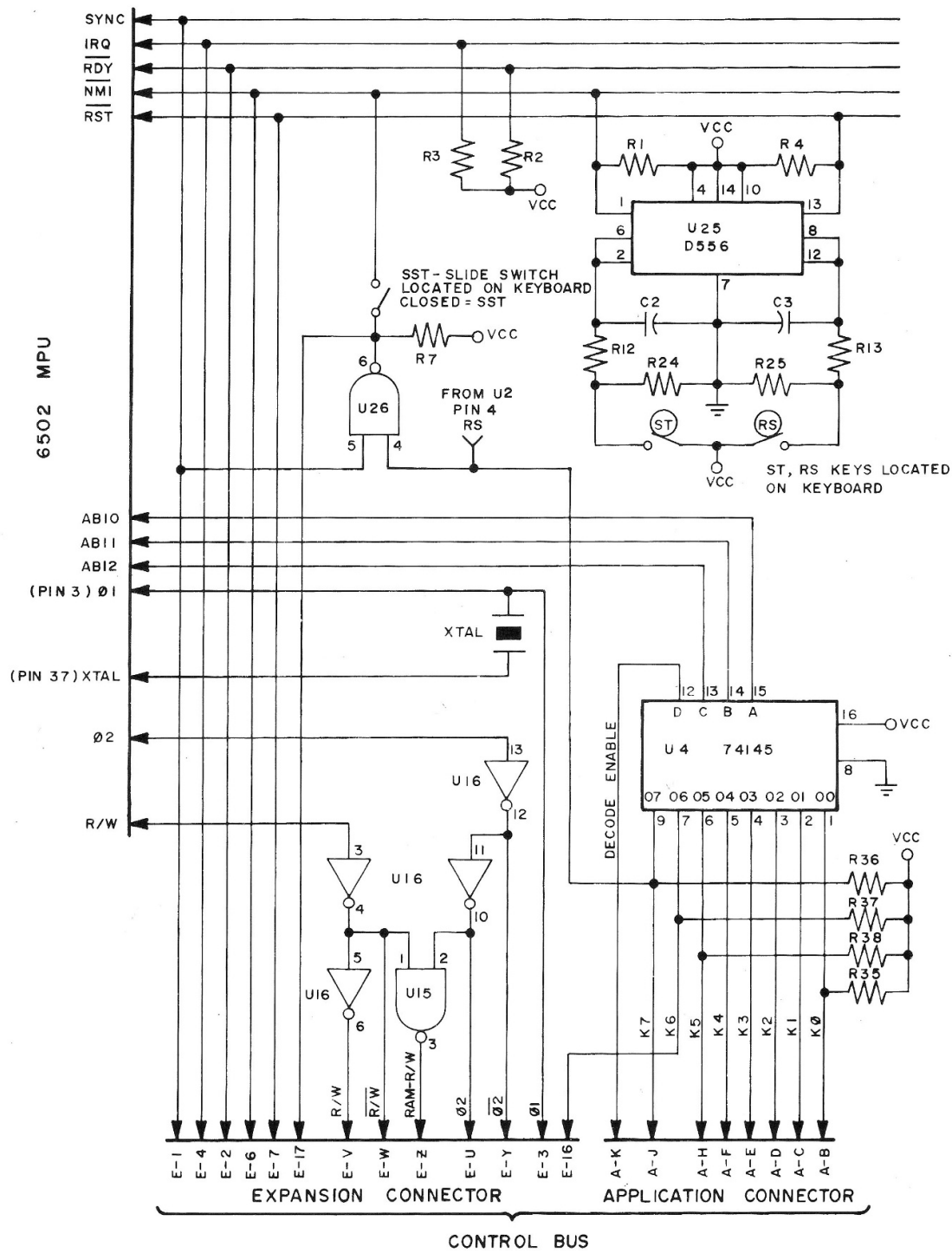
You may refer to the Hardware Manual included with your KIM-1 module for additional details on the operating characteristics of the 6502 and 6530 arrays as well as detailed information on system timing.



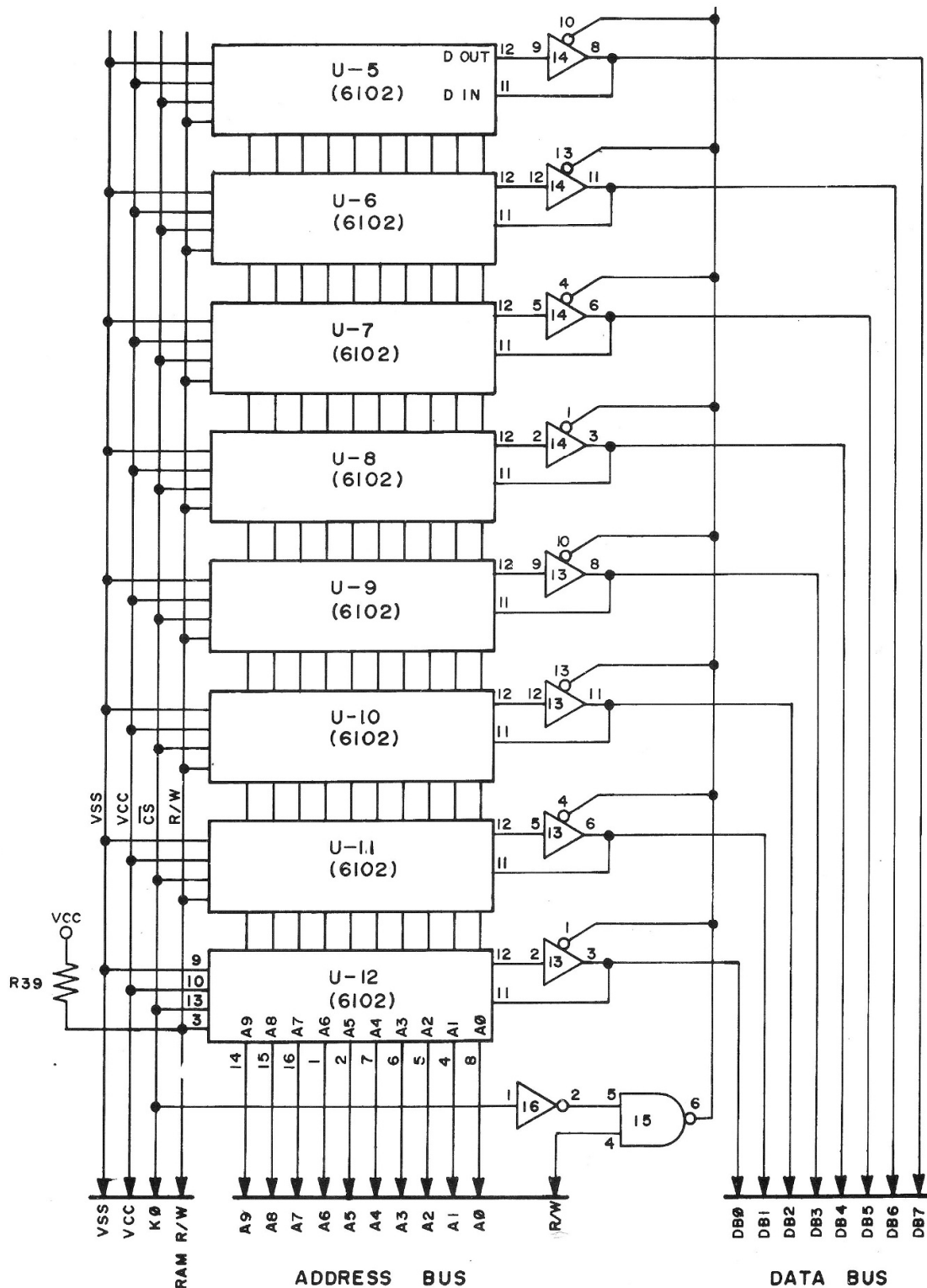
KIM-1 Block Diagram
FIGURE 3.1



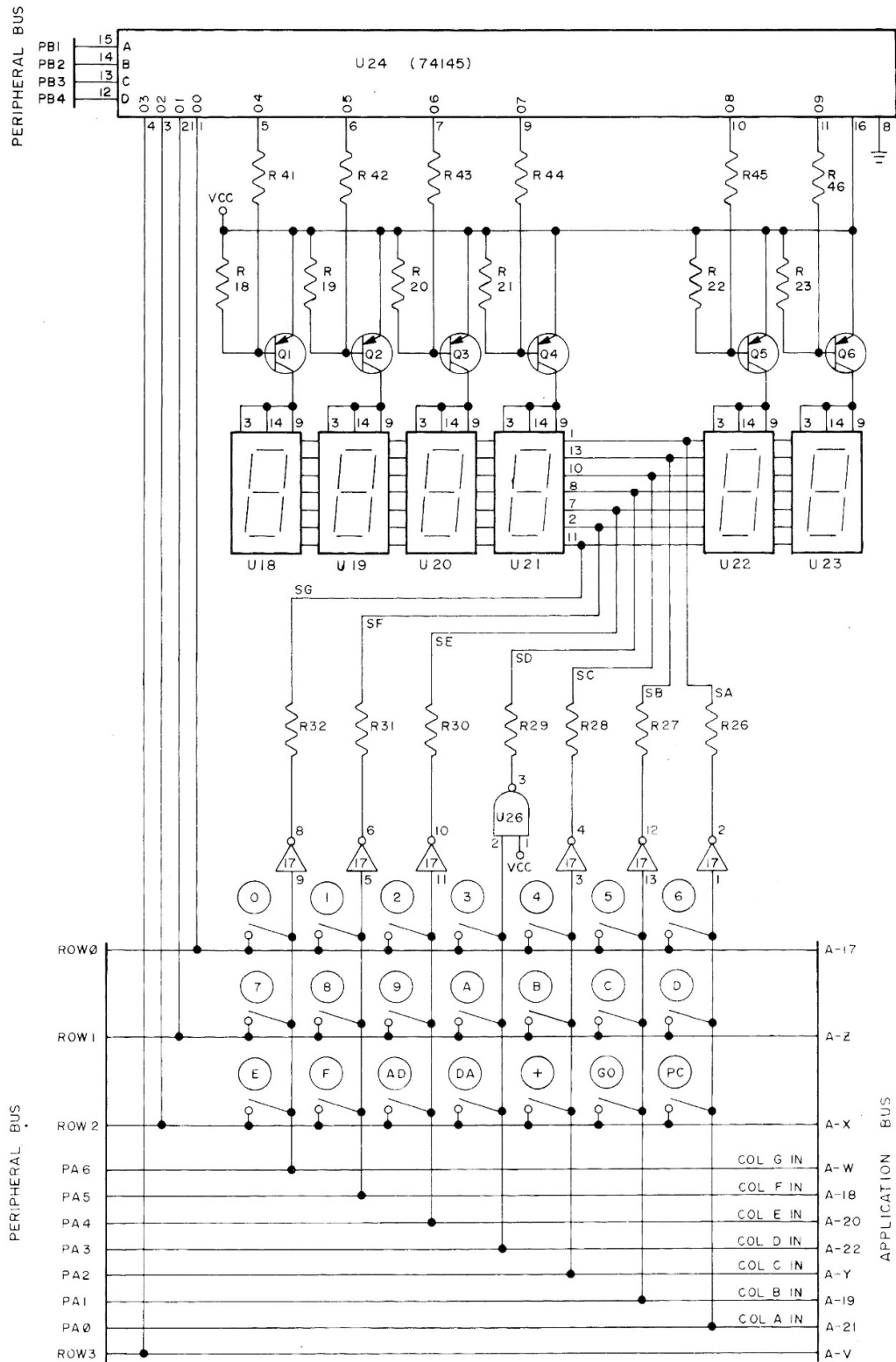
Detailed Block Diagram
FIGURE 3.2



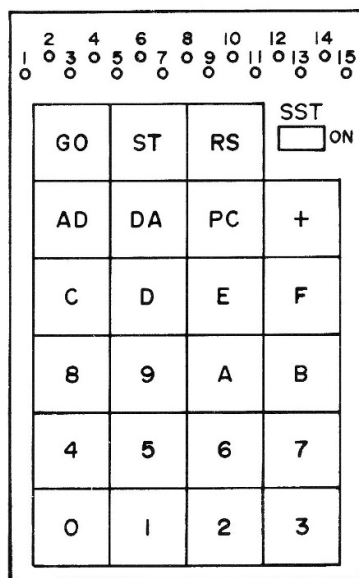
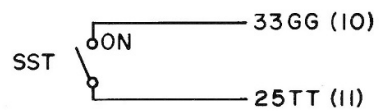
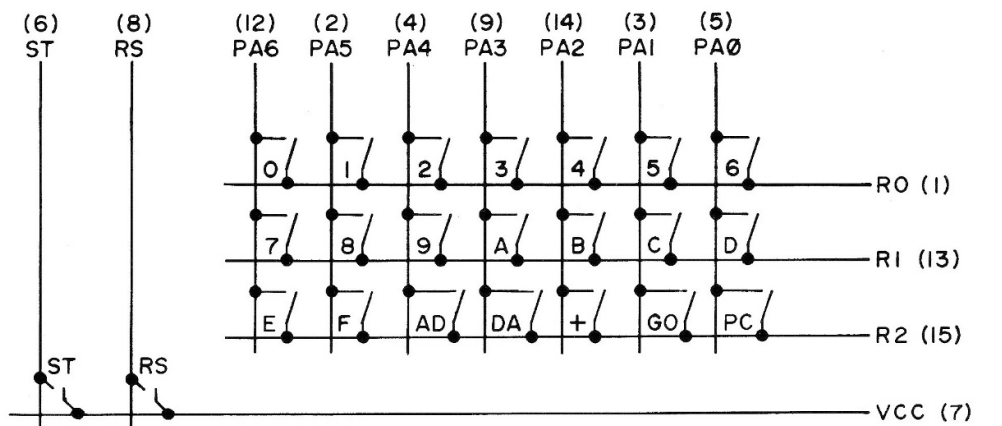
Control and Timing
FIGURE 3.3



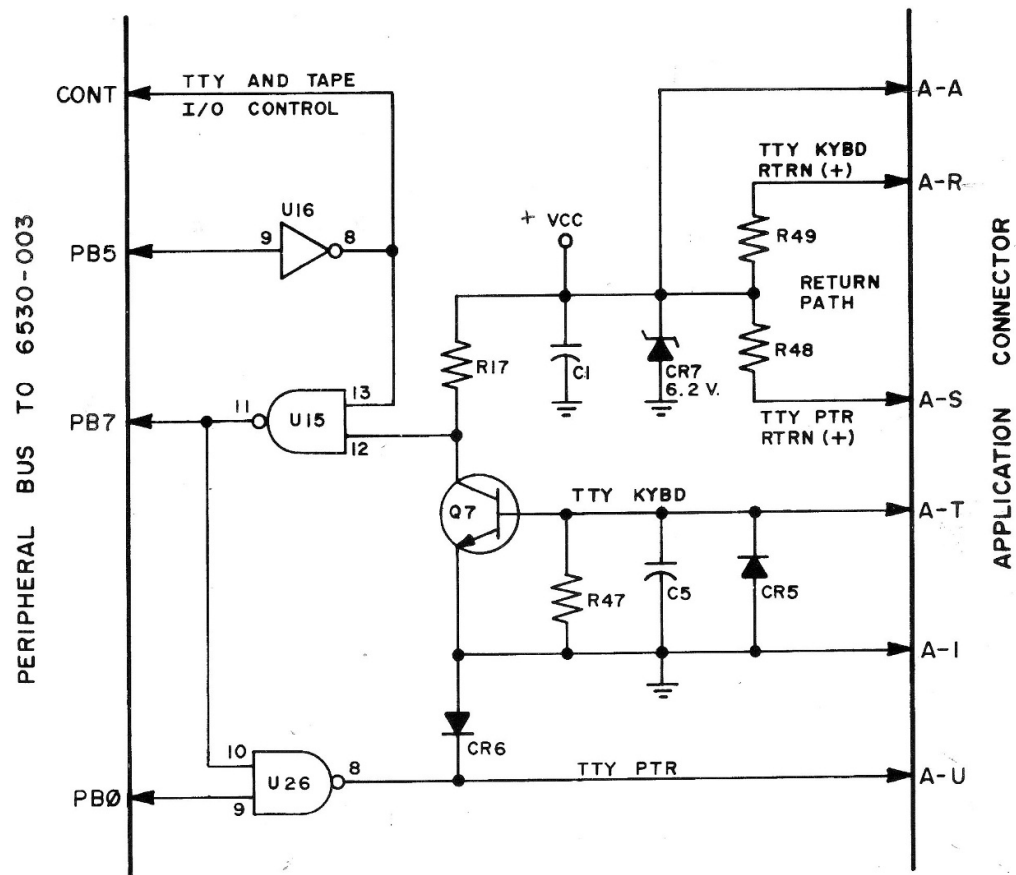
1Kx8 RAM Memory
FIGURE 3.4



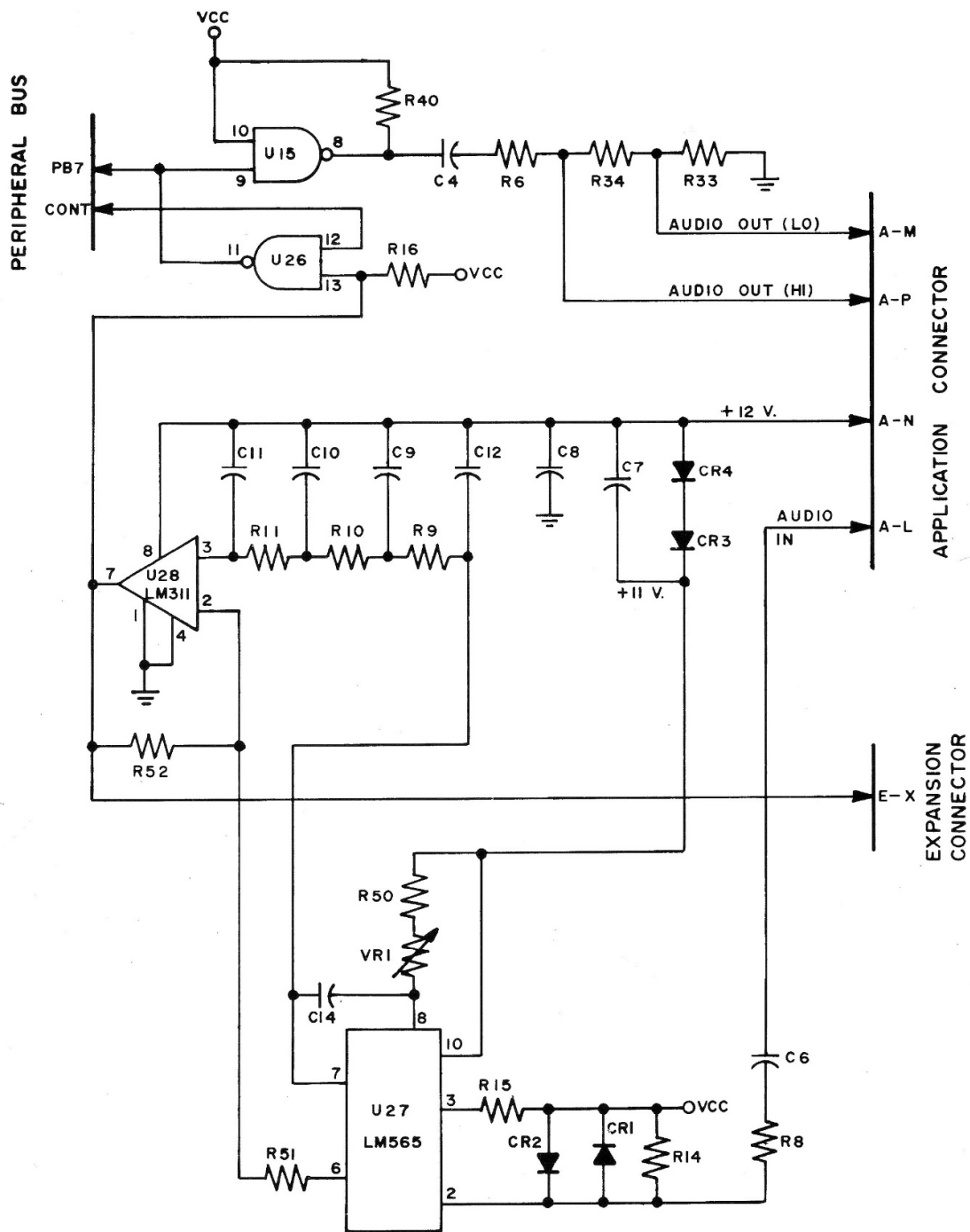
Keyboard and Display
FIGURE 3.5



Keyboard Detail
FIGURE 3.6



TTY Interface
FIGURE 3.7



Audio Tape Interface
FIGURE 3.8

22	KB Col D
21	KB Col A
20	KB Col E
19	KB Col B
18	KB Col F
17	KB Row 0
16	PB5
15	PB7
14	PA0
13	PB4
12	PB3
11	PB2
10	PB1
9	PB0
8	PA7
7	PA6
6	PA5
5	PA4
4	PA1
3	PA2
2	PA3
1	VSS GND

Z	KB Row 1
Y	KB Col C
X	KB Row 2
W	KB Col G
V	KB Row 3
U	TTY KYBD
T	TTY PTR
S	TTY PTR RTRN(+)
R	TTY KYBD RTRN(+)
P	AUDIO OUT HI
N	+12v
M	AUDIO OUT LO
L	AUDIO IN
K	DECODE ENAB
J	K7
H	K5
F	K4
E	K3
D	K2
C	K1
B	K0
A	VCC +5v

Application Connector
FIGURE 3.9

22	VSS GND
21	VCC +5
20	
19	
18	
17	SST OUT
16	K6
15	DB0
14	DB1
13	DB2
12	DB3
11	DB4
10	DB5
9	DB6
8	DB7
7	RST
6	NMI
5	RO
4	IRQ
3	01
2	RDY
1	SYNC

Z	RAM/R/W
Y	02
X	PLL TEST
W	R/W
V	R/W
U	02
T	AB15
S	AB14
R	AB13
P	AB12
N	AB11
M	AB10
L	AB9
K	AB8
J	AB7
H	AB6
F	AB5
E	AB4
D	AB3
C	AB2
B	AB1
A	AB0

Expansion Connector
FIGURE 3.10

3.2 KIM-1 MEMORY ALLOCATION

It has been stated that the 6502 microprocessor array included in the KIM-1 system is capable of addressing any of 65,536 memory locations. Obviously, we have not included that much memory in your KIM-1 system and this section is intended to detail for you exactly what memory locations are included in the system and where they are located (their exact addresses).

Each byte of memory in the system is understood to include 8 bits. Also, you should note that any addressable location in the system may be performing any one of four functions:

1. A ROM byte - read-only memory in which we have stored the operating program.
2. A RAM byte - read/write memory for storage of variable data.
3. An I/O location - these locations include both direction registers which define the I/O pins to be either input pins or output pins, and the actual data buffer locations containing the data to be transmitted on output pins or the data read from input pins. Any I/O location may be viewed as a read/write memory location with a specific address.
4. An Interval Timer location - a series of addresses are reserved for each interval timer in the system. Again, you may write to the timer to define its counting period or read from the timer to determine its exact state.

Figure 3-11 shows a block diagram detailing all memory blocks in the KIM-1 system. Figure 3-12 provides a memory map showing all addressable locations included in the system and their relationship to each other. Note also the areas in the memory map indicated as available for expansion. (Section 6 of the manual provides more detail on the subject of memory expansion). Finally, Figure 3-13 provides a complete listing of all important memory locations and will be referenced frequently by you when writing your application programs.

Referring to Figure 3-12, note that the memory map shows a block of 8192 address locations all existing in the lowest address space within the possible 65,536 address locations. This address space is further divided into eight blocks of 1024 locations each. Each 1024 block is further divided into four pages of 256 locations each. The "K" reference defines a specific block of 1024 locations and refers to the "K" number of the address decoder included within the system control logic. The "page" reference defines a specific group of 256 addresses. A total of 32 pages (0 to 31) are included in the 8192 address locations. The hex codes for certain addresses are shown at strategic locations in the memory map.

Beginning from the highest address location of the 8192, note that the first 1024 block (K7) is assigned to the ROM of the 6530-002 and the second 1024 block (K6) is assigned to the ROM of the 6530-003. The entire operating program of the KIM-1 system is included in these two blocks.

Next in order, a portion of the K5 block is dedicated to the RAM, I/O, and Timer locations of the two 6530 arrays. An expanded view of this address space is shown in Figure 3-12. Note that the RAM addresses for the 6530-002 (Hex 17EC to 17FF) are reserved for use by the operating program and should not appear in a user generated application program. The same is true for the I/O and Timer locations of the 6530-002 which also are reserved for use by the operating programs.

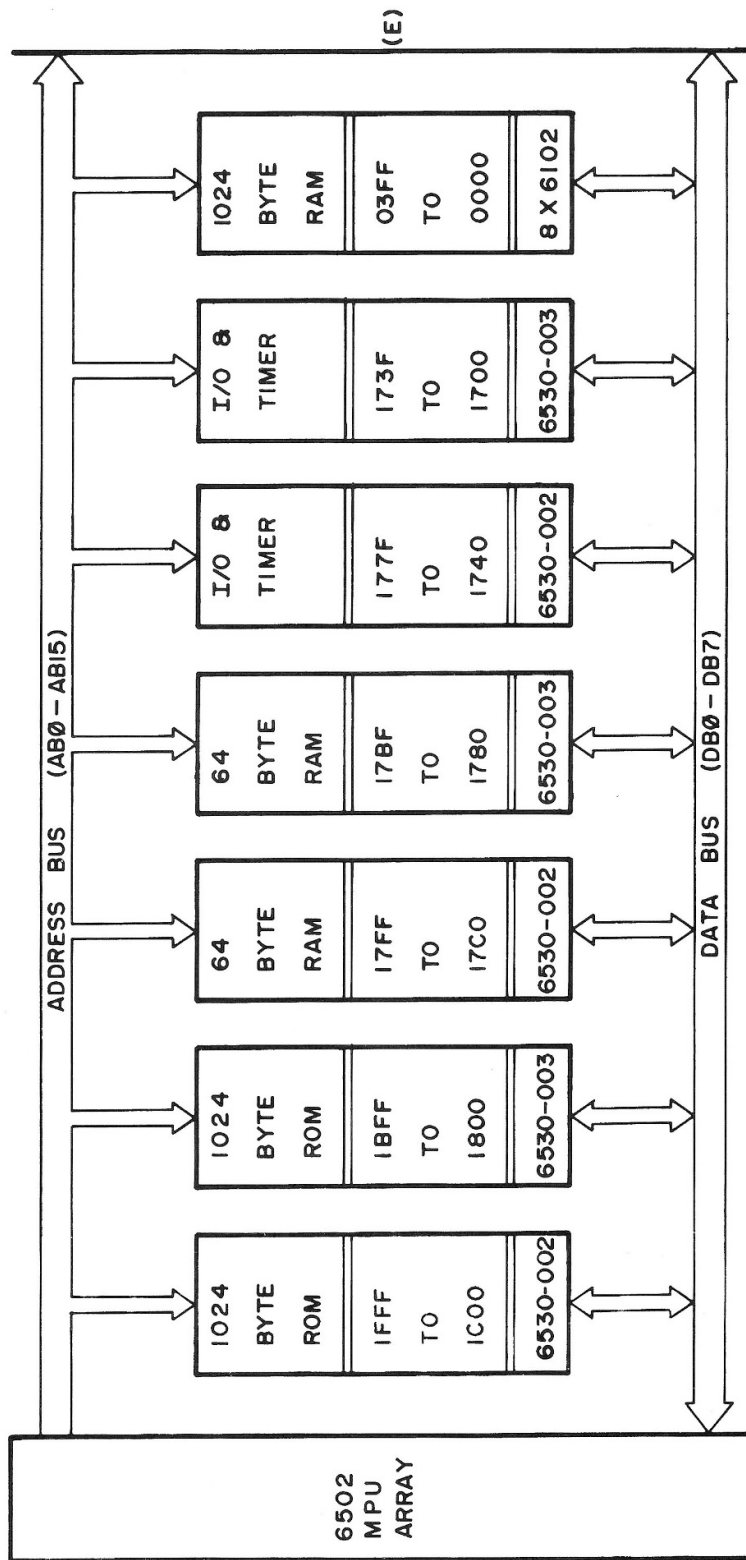
The next four blocks in order (K4, K3, K2, K1) are reserved for additional memory in an expanded system. In Section 6, the methods for adding memory will be discussed.

Finally, the lowest 1024 address locations (K0) are assigned to the static RAM included within the KIM-1 system. You should note that within this block, Page 0 and Page 1 have special significance. Page 1 is used as the system stack onto which return addresses and machine status words are pushed as the system responds to interrupts and subroutine commands. Page 0 has significance for certain of the special addressing modes available when programming for the 6502 microprocessor array.

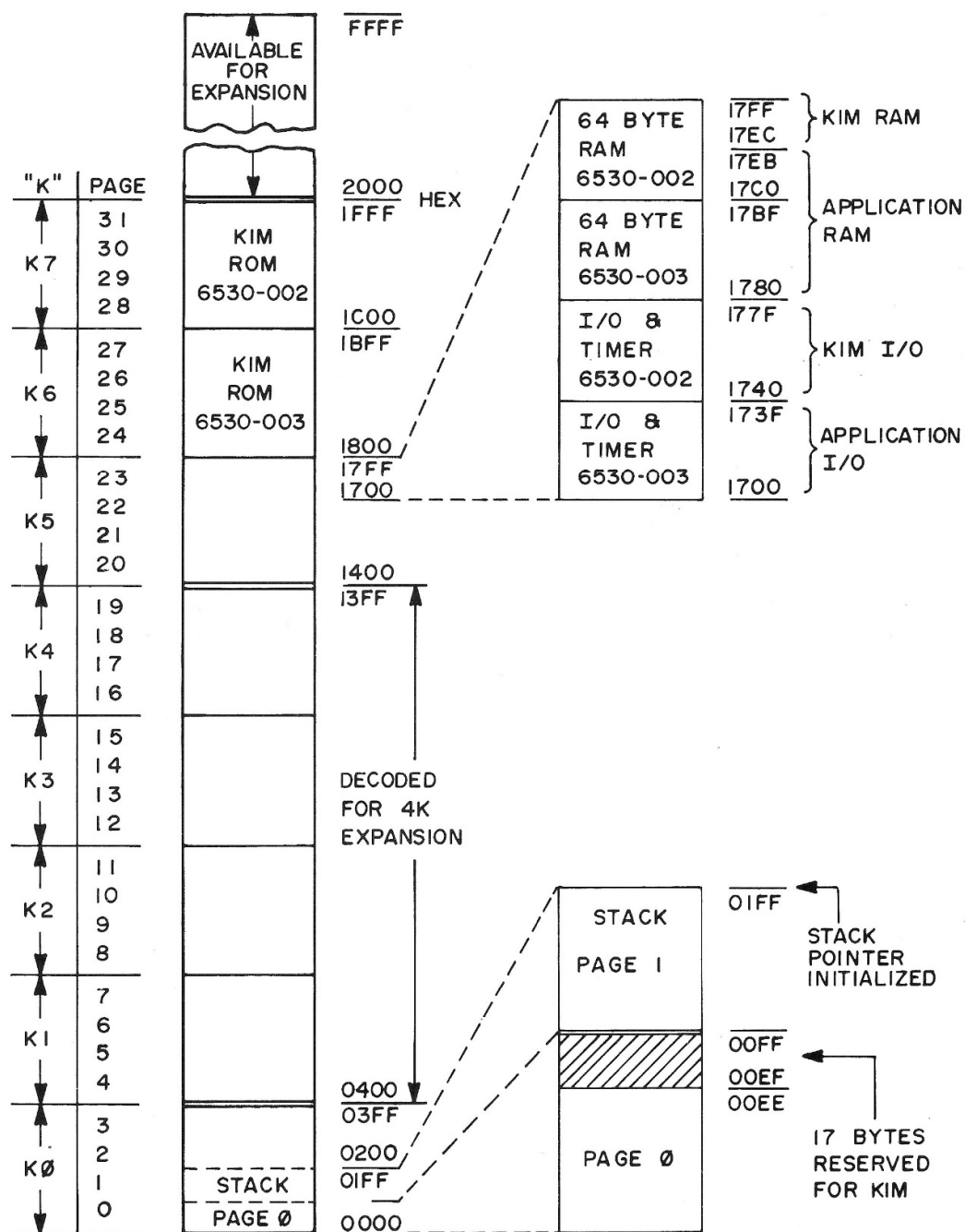
Figure 3-12 shows an expanded view of Page 0 and Page 1. Note that 17 addresses (00EF to 00FF) are reserved for use by the operating program and must never appear in the user generated application program. Also, note the comment that a maximum of eight locations may be required on the stack (Page 1) to service operating program interrupts.

In summary, the user generated application program may make use of the following areas of memory:

1. All of Page 0 except 00EF to 00FF
2. All of Page 1 (remember that the stack will extend an extra 8 bytes deep to accommodate the operating program).
3. All of Page 2 and Page 3.
4. In Page 23:
 - All I/O locations from 1700 to 173F
 - All 64 bytes of RAM from 1780 to 17BF
 - An additional 44 bytes of RAM from 17C0 to 17EB



Memory Block Diagram
FIGURE 3.11



Memory Map
FIGURE 3.12

ADDRESS	AREA	LABEL	FUNCTION
00EF	Machine Register Storage Buffer	PCL	Program Counter - Low Order Byte
00F0		PCH	Program Counter - High Order Byte
00F1		P	Status Register
00F2		SP	Stack Pointer
00F3		A	Accumulator
00F4		Y	Y-Index Register
00F5		X	X-Index Register
1700	Application I/O	PAD	6530-003 A Data Register
1701		PADD	6530-003 A Data Direction Register
1702		PBD	6530-003 B Data Register
1703		PBDD	6530-003 B Data Direction Register
1704	Interval Timer		6530-003 Interval Timer (See Section 1.6 of Hardware Manual)
170F			
17F5	Audio Tape Load & Dump	SAL	Starting Address - Low Order Byte
17F6		SAH	Starting Address - High Order Byte
17F7		EAL	Ending Address - Low Order Byte
17F8		EAH	Ending Address - High Order Byte
17F9		ID	File Identification Number
17FA	Interrupt Vectors	NMIL	NMI Vector - Low Order Byte
17FB		NMIH	NMI Vector - High Order Byte
17FC		RSTL	RST Vector - Low Order Byte
17FD		RSTH	RST Vector - High Order Byte
17FE		IRQL	IRQ Vector - Low Order Byte
17FF		IRQH	IRQ Vector - High Order Byte
1800	Audio Tape	DUMPT	Start Address - Audio Tape Dump
1873		LOADT	Start Address - Audio Tape Load
1C00	STOP Key + SST		Start Address for NMI using KIM "Save Machine" Routine (Load in 17FA & 17FB)
17F8	Paper Tape Dump (Q)	EAL	Ending Address - Low Order Byte
17F9		EAH	Ending Address - High Order Byte

Special Memory Addresses
FIGURE 3.13

3.3 KIM-1 OPERATING PROGRAMS

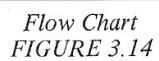
Figure 3-14 shows a simplified flow chart of the KIM-1 operating programs. This section provides a brief explanation of these programs to assist you in understanding the various operating modes of the system.

First, you should note that when power is first applied to your KIM-1 module and the RS (reset) key is depressed, control of the system automatically is assumed by the operating program. This is true, as well, for any succeeding depression of the reset key.

For each depression of the reset key, the system is initialized. At this time, stack pointer values are set, the I/O configuration is established, and essential status flags are conditioned. Next the program determines whether the system is to respond to TTY inputs or is to operate with the keyboard and display on the KIM-1 module.

If the TTY mode has been selected, the program halts and awaits a first key depression from the TTY (the RubOut Key). Upon receipt of this key depression, the program automatically performs a bit rate measurement and stores the correct value for use in receiving and decoding succeeding data transfers from the TTY. Note that this bit rate measurement is performed after each depression of the reset key.

The program will proceed immediately to a routine causing the prompting message ("KIM") to be typed on the TTY. Now, the program halts at the loop called "Get Character". As each key is depressed on the TTY, the coded data is accepted and analyzed in the routine called "Execute Key". The various keys depressed will cause the program to branch to the appropriate subroutines required to perform the desired operation. Upon completion of the individual key executions, the program returns to the "Get Key" loop and awaits the next key depression.



Exit from the TTY processing loop will occur in response to:

1. A depression of the reset key,
2. A depression of the G key which initiates execution of the application program, or
3. A change in the mode from TTY to Keyboard/Display.

If, after system reset and initialization, the Keyboard/Display mode (KB) is determined to be in effect, the program will proceed directly to display, and keyboard scan routines. The program will cause the display scan to occur continuously ("Display Cell") until one of the keys on the keyboard is depressed (AK?). Key validation is performed during an additional scan cycle. If the key is truly depressed (not noise), the program proceeds to the routine called "Get Key" in which the exact key depressed is defined. Next, the program moves to the "Execute Key" routine where branches to appropriate execution routines will be performed. Finally, after key execution, the program returns to the "Display Cell" routine and waits for the key to be released. When no key is depressed, the program returns to the normal "Display Cell" routine and awaits the next key depression.

In either the TTY or KB modes, the audio tape load or dump routines may be executed using appropriate commands from the selected keyboards. In either case, completion of the tape load or dump routine allows the program to return to the "Start" position which will, as usual, activate the KIM-1 display or cause the "KIM" prompting message on the TTY.

You should note the use of the Stop key to activate the non-maskable interrupt input ($\overline{\text{NMI}}$) of the 6502 microprocessor array. Depression of this key causes an unconditional termination of program execution, a saving of machine status registers on the stack, and a return to the control of the operating program.

A second interrupt input is available and referred to as $\overline{\text{IRQ}}$. This interrupt may be defined by the user and will cause the program to jump to any location defined by the user in his program.

CHAPTER 4

OPERATING THE KIM-1 SYSTEM

Now that you have a better idea of what is included in your KIM-1 system and how it operates, its time to provide you with detailed procedures for all of the operations you can perform with the system. We will separate our operating procedures into three areas giving specific direction for the use of the KIM-1 keyboard and display, the audio tape recorder, and the serial teleprinter (TTY).

4.1 USING THE KIM-1 KEYBOARD AND DISPLAY

A brief study of your keyboard shows a total of 23 keys and one slide switch. First, let's list the purpose of each key:

- 0 to F - Sixteen keys used to define the hex code of address or data
- AD - selects the address entry mode
- DA - selects the data entry mode
- + - increments the address by +1 but does not change the entry mode
- PC - recalls the address stored in the Program Counter locations (PCH, PCL) to the display
- RS - causes a total system reset and a return to the control of the operating program
- GO - causes program execution to begin starting at the address shown on the display
- ST - terminates the execution of a program and causes a return to the control of the operating program

You have seen in an earlier chapter that the six digit display includes a four digit display of an address (left four digits) and a two digit display of data (right two digits).

Using only the KIM-1 keyboard and display, you may perform any of the following operations:

1. Select an Address

Press **AD** followed by any four of the hex entry keys. The address selected will appear on the display. If an entry error is made, just continue to enter the correct hex keys until the desired address shows on the display. Regardless of what address is selected, the data field of the display will show the data stored at that address.

2. Modify Data

After selecting the proper address, press **DA** followed by two hex entry keys which correctly define the data to be stored at the selected address. The data entered will appear in the data field of the display to indicate that the desired code has already been entered.

Note that it is possible for you to select an address of a ROM memory cell or even the address of a memory cell that does not exist in your system. In these cases, you will not be able to change the data display since it is clearly not possible for the system to write data to a ROM cell or a non-existent memory location.

3. Increment the Address

By pressing the **+** key the address displayed is automatically increased by +1. Of course, the data stored at the new address will appear on the display. This operation is useful when a number of successive address locations must be read or modified. Note that the use of the **+** key will not change the entry mode. If you had previously pressed the **AD** key, you remain in the address entry mode and a previous depression of the **DA** means you remain in the data entry mode.

4. Recall Program Counter

Whenever the NMI interrupt pin of the 6502 microprocessor array is activated, the program execution in progress will halt and the internal registers of the 6502 are saved in special memory locations before the control of the system is returned to the operating program. In the KIM-1 system, the NMI interrupt may occur in response to a depression of the **ST** key (stop) or, when operating in the Single Step mode, after each program instruction is executed following the depression of the **GO** key.

The **PC** key allows you automatically to recall the value of the Program Counter at the time an interrupt occurred. You may have performed a variety of operations since the interrupt such as inspecting the contents of various machine registers stored at specific memory locations. However, when you press the **PC** key, the contents of the Program Counter at the time of the interrupt are recalled to the address field of the display. You now may continue program execution from that point by pressing the **GO** key.

5. Execute a Program

Select the starting address of the desired program. Now, press the **GO** key and program execution will commence starting with the address appearing on the display.

6. Terminate a Program

The **ST** key is provided to allow termination of program execution. As mentioned earlier, the **ST** key activates the NMI interrupt input of the 6502 microprocessor array.

Note: The **ST** key will operate correctly only if you store the correct interrupt vector at locations 17FA and 17FB. For most of your work with the KIM-1 system, you should store the address 1C00 in these locations as follows:

AD			
1	7	F	A
DA		0	0
+		1	C

Now, when the $\overline{\text{NMI}}$ interrupt occurs, the program will return to location 1C00 and will proceed to save all machine registers before returning control to the operating program.

You should remember to define the $\overline{\text{NMI}}$ vector each time the power to the system has been interrupted. A failure of the system to react to the $\boxed{\text{ST}}$ key means you have forgotten to define the $\overline{\text{NMI}}$ vector.

7. Single Step Program Execution

In the process of debugging a new program, you will find the single step execution mode helpful. To operate in this mode, move the SST slide switch to the ON position (to your right). Now, depress the $\boxed{\text{GO}}$ key for each desired execution of a program step. The display will show the address and data for the next instruction to be executed. Note that in the course of stepping through a program, certain addresses will appear to be skipped. A program instruction will occupy one, two, or three bytes of memory depending upon the type of instruction. In single instruction mode, all of the bytes involved in the execution of the instruction are accessed and the program will halt only on the first byte of each successive instruction.

Note: SST mode also makes use of the $\overline{\text{NMI}}$ interrupt of the 6502 microprocessor array. Again, the $\overline{\text{NMI}}$ vector must be defined as described in (6) above if the SST mode is to work correctly.

This covers all of the standard operations you may perform from the KIM-1 keyboard. Using combinations of the operations described, you may wish to perform certain specialized tasks as follows:

1. Define the $\overline{\text{IRQ}}$ Vector

You will recall that a separate interrupt input labelled $\overline{\text{IRQ}}$ is available as an input to the 6502 microprocessor array. If you wish to use this feature, you should enter the address to which the program will jump. The $\overline{\text{IRQ}}$ vector is stored in locations 17FE and 17FF.

2. Interrogate Machine Status

We have mentioned that after an $\overline{\text{NMI}}$ interrupt in response to the $\boxed{\text{ST}}$ key or during the SST mode, the contents of various machine registers are stored in specific memory locations. If you wish to inspect these locations, their addresses are:


```

00EF = PCL
00F0 = PCH
00F1 = Status Register (P)
00F2 = Stack Pointer (SP)
00F3 = Accumulator (A)
00F4 = X Index Register
00F5 = Y Index Register

```

4.2 USING THE AUDIO TAPE RECORDER

There are two basic operations possible when working with your audio tape system. You may transfer data from the KIM-1 memory and record it on tape. Or, you may read back a previously recorded tape, transferring the data on tape into the KIM-1 memory.

Recording on Audio Tape

The procedure for recording on audio tape requires that you perform the following steps:

1. Define an identification number (ID) for the data block you are about to record. This two digit number is loaded into address 17F9. Don't use ID = 00 or ID = FF.
2. Define the starting address of the data block to be transferred. This address is to be loaded into locations:


```

17F5 = Starting Address Low (SAL)
17F6 = Starting Address High (SAH)

```
3. Define the ending address as one greater than the last address in the data block to be recorded. The ending address is to be loaded into locations:


```

17F7 = End Address Low (EAL)
17F8 = End Address High (EAH)

```

As an example, assume you wish to record a data block from address 0200 up to and including address 03FF. (All of Pages 2 and 3). You wish to assign an ID number of 06 to this block. Using the KIM-1 keyboard, you should load the data shown into the addresses indicated so that:

```

17F5 = 00 (SAL)
17F6 = 02 (SAH)
17F7 = 00 (EAL)
17F8 = 04 (EAH) } = 03FF + 1
17F9 = 06 (ID)

```

Note that the ending address must be greater than the starting address for proper operation.

4. Assuming that you are using a new cassette on which no data has been stored previously, insert the cassette in the unit and rewind the tape to its start position.
5. Select the starting address of the tape record program. This address is 1800.
6. Select the Play/Record mode of the audio unit and allow several seconds for the tape to begin to move.
7. Press the **GO** key and the recording process will begin. The display will be blanked for a period and then will relight showing 0000 xx. This means that the data block selected has been recorded.
8. You may now stop the tape or allow some additional seconds of blank tape and then stop the unit.

Loading Data From Audio Tape

The procedure for loading data from an audio tape into the KIM-1 memory requires that you perform the following steps:

1. Define the ID number of the data block to be loaded from tape. The ID number is loaded into address 17F9.
2. Select the starting address of the Tape Load program. This address is 1873_{HEX}.
3. Press the **GO** key. The KIM-1 system is now waiting for the appearance of data from the tape unit.
4. Load the cassette and, presuming you do not know where on the tape the data block is recorded, rewind the tape to its starting position. Check the volume control setting.
5. Start the audio tape unit in its Play mode and observe that the tape begins to move.
6. Wait for the KIM-1 display to relight showing 0000 xx. This means the data block has been loaded successfully from the tape into the KIM-1 memory. If the display relights with FFFF xx, the correct data block has been found but there has been an error detected during the read operation. If the tape continues to run and the display never relights, the system has not been successful in finding the data block with the specific ID number you requested.

7. If in step (1), you had selected an ID = 00, the ID number recorded on the tape will be ignored and the system will read the first valid data block encountered on the tape. The data read from the tape will be loaded into memory address as specified on the tape.
8. If, in step (1), you had selected an ID = FF, the ID number recorded on the tape will be ignored and the system will read the first valid data block encountered on the tape. In addition, the data block will be loaded into successive memory locations beginning at the address specified in locations 17F5 and 17F6 (SAL, SAH) instead of the locations specified on the tape.

Special Operations with Audio Tape

The KIM-1 system causes data to be recorded on audio tape with a specific format as detailed in Appendix E. Each recorded data block is preceeded by a group of synchronizing characters together with an identification code to define the specific block. Data blocks may be of arbitrary length.

With a little care, there is no reason for you not to include a number of recorded data blocks on the same tape. If you are recording blocks in sequence and have not rewound the tape between blocks, you need only specify the parameters of each new block (ID, SAL, SAH, EAH, EAL) and proceed with recording the new block.

If the tape has been rewound, you will need to know the ID number of the last recorded data block. Rewind the tape to its starting point and set up the parameters required to read the last recorded data block. After reading this block, stop the tape and you may now proceed to add a new block or blocks to the tape.

If you wish, you may add voice messages between the recorded data blocks on the tape. The KIM-1 system will ignore these audio messages when the tape is read back. Of course, you will need to install an earphone or speaker in parallel with the KIM-1 audio tape data input pin in order to hear the voice messages.

We do not recommend that you attempt to record data blocks in areas of the tape which have been used previously for recorded data. Variations in tape speed and block lengths can result in overlapping of recorded data which may be read incorrectly by the KIM-1 system.

4.3 USING A SERIAL TELEPRINTER

The addition of a serial teleprinter (such as the Teletype Model 33ASR) to work with the KIM-1 system permits a variety of special operations to be performed. In all cases, you define desired operations by depressing the proper keys while simultaneously producing a hard-copy printed record of each operation. If your teleprinter is equipped with a paper tape reader/punch, you may generate or read paper tapes using the KIM-1 system. Using the serial teleprinter, you may perform the following operations:

Select an Address

Type four hex keys (0 to F) to define the desired address. Next, press the bar.

The printer will respond showing the address code selected followed by a two digit hex code for data stored at the selected address location:

```
Type:           1234 
Printer Responds: 1234 AF
```

showing that the data AF is stored at location 1234.

Modify Data

Select an address as in the previous section. Now type two hex characters to define the data to be stored at that address. Next type the \odot key to authorize the modification of data at the selected address:

```
Type:           1234 
Printer Responds: 1234 AF
Type:                                     6D  $\odot$ 
Printer Responds: 1235 B7
```

Note that the selected address (1234) has been modified and the system increments automatically to the next address (1235).

Note: Leading zero's need not be entered for either address or data fields: For example:

```
EF  selects address 00EF
E   selects address 000E
A   $\odot$  enters data 0A
    $\odot$  enters data 00 (etc.)
```

Step to Next Address

Type **CR** to step to the next address without modifying the current address:

See Printed:	1234	AF	
Type:			CR
Printer Responds:	1235	B7	
Type:			CR
Printer Responds:	1236	C8	(etc.)

Step to Preceding Address

Type **LF** to step back to the preceding address:

See Printed:	1234	AF	
Type:			LF
Printer Responds:	1233	9D	
Type:			LF
Printer Responds:	1232	8E	(etc.)

Abort Current Operation

Type **RUB OUT** to terminate the current operation. The prompting message will be printed ("KIM") indicating that a new operation may proceed:

Type:	1264	RUB OUT
Printer Responds:	KIM	
	xxxx xx	
Type:	1234	SPACE
Printer Responds:	1234	AF

In the example, the **RUB OUT** key is used to correct an erroneous address selection.

Note: The **RUB OUT** key must be depressed after each depression of the KIM-1 reset key in order to allow the operating program to define the serial bit rate for the teleprinter.

Load Paper Tape

Paper Tapes suitable for use with the KIM-1 system are generated using the format shown in Appendix F. To read such a tape into the KIM-1 system, proceed as follows:

1. Load the punched paper tape on to the tape mechanism
2. Type Ⓛ
3. Activate the paper tape reader

The paper tape will advance and data will be loaded into addresses as specified on the tape. A printed copy of the data read will be generated simultaneously with the reading of the paper tape.

Check-sums are generated during the reading of the paper tape and are compared to check-sums already contained on the tape. A check-sum error will cause an error message to appear in the printed copy.

Punch Paper Tape

The KIM-1 system can be used to punch paper tapes having the format described in Appendix F. The procedures for generating these tapes is as follows:

1. Define the starting address and ending address of the data block to be punched on the paper tape.
2. Load blank paper tape on the punch unit and activate the punch.

Type:	ⓕⓈ	SPACE
See Printed	00F5	xx
Type:	⓪⓪●	
See Printed:	00F6	xx
Type:	⓪②●	
See Printed:	00F7	xx
Type:	ⓕⓕ●	
See Printed:	00F8	
Type:	⓪③●	
See Printed:	00F9	xx

You have now loaded a starting address (0200) into address locations 00F5 (SAL) and 00F6 (SAH). The ending address (03FF) has been loaded into locations 00F7 (EAL) and 00F8 (EAH).

3. Now type ③

The paper tape will advance and punching of the data will proceed. Simultaneously, a printed record of the data will be typed.

List Program

A printed record of the contents of the KIM-1 memory may be typed. The procedure is the same as for punching paper tape except that the punch mechanism is not activated.

Execute Program

To initiate execution of a program using the TTY keyboard, the following procedures should be followed:

1. Enter the starting address of the program

2. Type ③

For example, to begin program execution from address location 0200:

Type: ②①① SPACE
See Printed: 0200 xx
Type: ③

Program execution begins from location 0200 and will continue until the ST or RS keys of the KIM-1 module are depressed. The single step feature may be employed while in the TTY mode.

CHAPTER 5

LET'S TRY A REAL APPLICATION

It is not practical in this manual to describe every possible application or programming technique. However, now that you have become familiar with the basic elements and operating procedures of the KIM-1 system, this section will show you how to apply what you have learned in a simple but realistic application example.

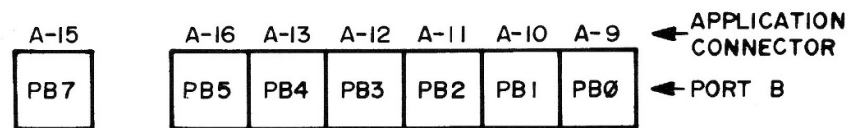
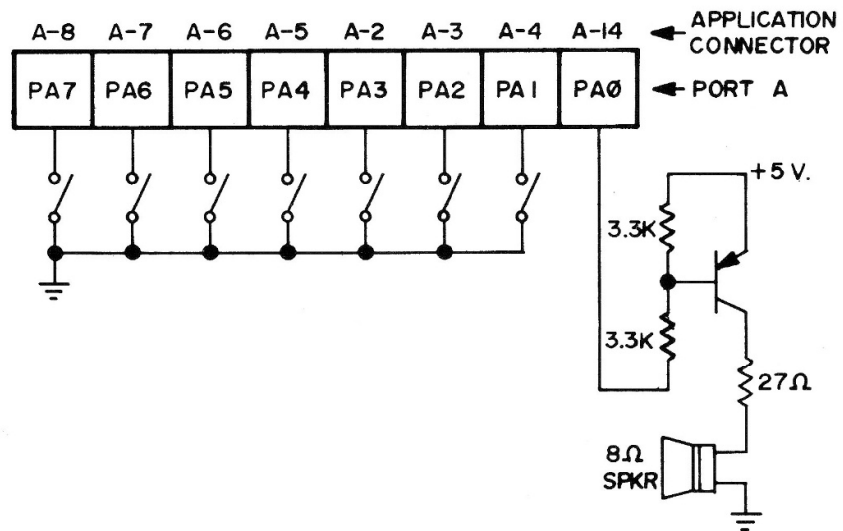
Our example will involve the generation of a variable frequency square wave which will be connected to a speaker to produce an audible tone. The frequency of the tone will be selected using a set of seven toggle switches. We will proceed through the example by defining the interface, writing and entering the program, and executing the program. Finally, we will study a series of program debugging techniques which will be useful to you for any new program you may write.

5.1 DEFINING THE INTERFACE

You will recall that a group of 15 I/O pins are brought to the Application connector from the 6530-003 array. The logic and circuit details concerning these I/O pins are described in Appendix H and in Section 1.6 of the Hardware Manual ("Peripheral Interface/Memory Device - - MCS 6530").

For our application example we will use eight of these I/O pins. One pin (PA0) will be used as an output line to supply a square wave to a driver circuit and speaker. The other seven I/O pins (PA1 to PA7) are defined as input points with a SPST toggle switch connected to each. Figure 5-1 shows the circuit configuration for this example. Note that the remaining seven I/O pins (the PB port) are not used for this problem.

For the switches connected to the input pins, we would like the sense of the switch to be defined as a logic "0" when open and a logic "1" when closed. By connecting the switches to ground, we are producing exactly the opposite sense and must remember to complement the switch states with software when we write our program. Also, we must define now that the switch at PA1 is to be the LSB (least significant bit) and the switch at PA7 is to be the MSB (most significant bit) of the seven bit binary word formed by all seven switches. In this way, the state of the switches can define a binary number from zero (all switches open) to 127_{DEC} (all switches closed).



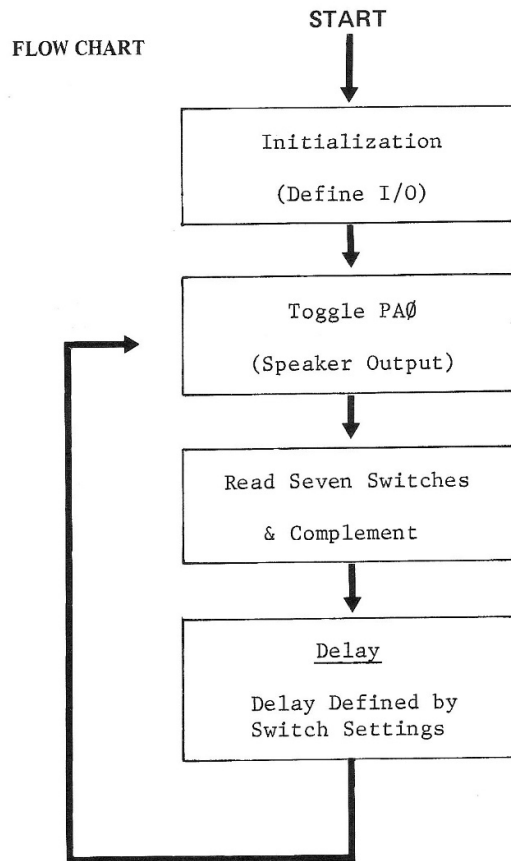
(THE B PORT IS NOT USED IN THIS EXAMPLE APPLICATION)

Speaker Application
FIGURE 5.1

5.2 WRITING THE PROGRAM

Having defined the interface for our application, we may proceed now to write our program. The effort proceeds in four stages:

1. Generate a flow chart
2. Generate assembly language code
3. Analyze the program
4. Generate machine language code



Briefly, our flow chart shows a first step of system initialization. During this step, we must define the I/O configuration of the system in that pin PA0 becomes the output to the speaker and that pins PA1 to PA7 become inputs from the seven switches.

After initialization, a loop is set up which begins by inverting the state of PA0 (Toggle PA0). Next, the state of the switches is read and the data is complemented to produce the correct "sense" from the switches. The value so read is used to define a delay before returning to the start of the loop and again toggling the state of PA0. A little thought will show that this loop will produce a square wave with a frequency determined by the setting of the seven switches.

Assembly Language Program

Our next task is to convert the simple flow chart into a program. The program is first written in "Assembly Language". You should refer to your Programming Manual to become familiar with all of the possible 6502 instructions (especially see Appendix B; Instruction Summary). Figure 5-2 shows the application example programmed in assembly language.

LABEL	OP CODE	OPERAND	MACHINE CYCLES	COMMENTS
INIT	LDA	#\$01	2	Define I/O 0=Input 1=Output
	STA	PADD	4	PADD = PORT A DATA DIRECTION REG.
START	INC	PAD	6	Toggle PA0, PA1-PA7 Inputs not affected
READ	LDA	PAD	4	READ switches into accumulator
	EOR	#\$FF	2	Complement switch value
	LSR	A	2	Shift Accumulator 1 bit to right
	TAX		2	Transfer final count into X-Index
DELAY	DEX		2	Delay by an amount specified
	BPL	DELAY	3,2	By the count in the X-Index
	BMI	START	3	Go To START
PADD	=\$1701			Define absolute address of Data Direction Reg. A
PAD	=\$1700			Define absolute address of Data Reg. A

Assembly Language Listing
FIGURE 5.2

You will note that each line of the program is broken into several fields:

- A label field permitting you to assign a "name" to a specific location in the program.
- An Operation Code field (Op Code) in which the exact instruction to be executed is defined.
- An Operand Field where the exact data required by the instruction is defined together with certain symbols defining addressing modes or data formats. Symbols encountered generally in MOS Technology, Inc. manuals are:

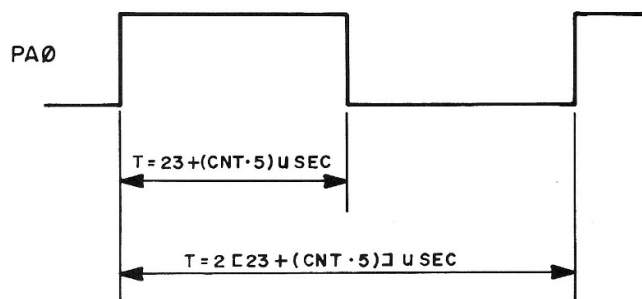
#	Immediate Addressing
\$	Hex Code
@	Octal Code
%	Binary Code
'	ASCII literal
=	Equates a label to a value

- A Machine Cycle field defining the total number of machine cycles required to execute an instruction. (This information is derived from Appendix B of the Programming Manual).
- A Comment Field where the programmer may define the intent of specific program steps.

Program Analysis

The inclusion of the "machine cycle" information of the program chart (Figure 5-2) allows us to analyze the exact timing relationships involved in our program example. Note that the KIM-1 system operates from a fixed frequency (1 MHz) oscillator with each machine cycle being 1 μ s. Therefore, an instruction like "INC PAD" which requires 6 machine cycles will be executed in a 6 μ s period.

By counting the total machine cycles occurring between each toggle of PA0, an equation for the square wave frequency can be developed. The actual frequency is determined by the position of the seven switches, the number of machine cycles between each toggle of PA0, and the basic clock rate (1 MHz) of the KIM-1 system. Figure 5-3 shows the waveform of the PA0 square wave and the derived equations for computing the exact frequency.



$$\text{FREQ} = \frac{1}{T} = \frac{10^6}{46 + 10 \cdot \text{CNT}} \text{ CPS}$$

NOTE: CNT EQUALS THE VALUE IN X-INDEX
WHICH WAS CALCULATED FROM THE
SEVEN SWITCHES $0 \leq \text{CNT} \leq 127$

Square Wave Output
FIGURE 5.3

Machine Language Coding

Our next problem is to convert our assembly language program into a program written in "machine language". The quickest and most foolproof method for accomplishing this conversion is by using the MOS Technology, Inc. Assembler (available for use on the time share services of United Computing Systems, Inc.). If you choose not to use this method, you will need to convert your source program to machine code using "paper-and-pencil" techniques.

You should proceed by constructing a table similar to that shown in Figure 5-4.

ADDRESS	INSTRUCTION			SOURCE CODE		
	BYTE 1	BYTE 2	BYTE 3	LABEL	OP CODE	OPERAND
0200	A9	01		INIT	LDA	#\$01
0202	8D	01	17		STA	PADD
0205	EE	00	17	START	INC	PAD
0208	AD	00	17	READ	LDA	PAD
020B	49	FF			EOR	#\$FF
020D	4A				LSR	A
020E	AA				TAX	
020F	CA			DELAY	DEX	
0210	10	FD			BPL	DELAY
0212	30	F1			BMI	START
0214						

Machine Language Code Table
FIGURE 5.4

The source code contained in your assembly language program (Figure 5-2) is entered into the table first. A column is provided to allow you to define the specific address at which an instruction is located. The Instruction column provides space for defining one, two, or three byte instructions. (Please refer to Appendix B of the Programming Manual or to your Programming Card for specific Op Codes).

As an example, the first source instruction is LDA #01 which, when translated, means load the accumulator with the byte stored in the next program location (hex 01). This is the "immediate" addressing mode defined by the "#" symbol. The Op Code for LDA# is A9. This value is entered in the first column under the heading, Instruction. The next column contains the hex 01 value defined by the source statement. The initial address for the program is inserted in the "Address" column as 0200 (an arbitrary selection). The total instruction LDA #01 now occupies address locations 0200 and 0201.

The next available address is 0202 which is inserted in the "Address" column for the next source instruction. In this manner, you will proceed through all of the source statements decoding each and entering one, two, or three bytes of machine code as required in the "Instruction" column. The "Address" column will contain the address of the first byte of machine code (the Op Code) for each source statement.

In cases where the operand of the source statement is a symbol, the address to which the symbol has been equated should be filled in as the proper machine code. For example, the source statement "INC PAD" requires the incrementing of data stored at a location "PAD" defined in our assembly programs to have the address: PAD = 1700. Therefore, the address 1700 is entered as the second and third bytes of the source statement "INC PAD". (See Figure 5-4). Note also that when entering an address, such as 1700, the low order byte (00) is entered first and immediately after the Op Code and the high order byte (17) is entered next as the third byte of the instruction.

When dealing with branch instructions (BPL, BMI, etc.), you will need to calculate the exact value of the offset which may be either positive (branch forward) or negative (branch backward). You should refer to Section 4.1.1 of the Programming Manual to explore "Basic Concept of Relative Branching." As an example, the source statement "BMI START" (See Figures 5-2 and 5-4) requires a branch backward by (-15) locations to the address labelled "START" (from address 0213 backward to 0205 inclusive).

(The 2's complement of the -15 displacement is F1_{HEX} which you should insert at location 0212). Had the branch been to a forward location the positive value of the offset would be inserted rather than the 2's complement value.

5.3 ENTERING THE PROGRAM

With the program now reduced to machine language code, you may enter the program address and data codes listed in Figure 5-4 following the procedures detailed in Section 2.4. The procedure for entering the program is as follows:

<u>Press Keys</u>	<u>See On Display</u>
AD 0 2 0 0	0200 xx
DA A 9	0200 A9
+ 0 1	0201 01
+ 8 D	0202 8D
+ 0 1	0203 01
+ 1 7	0204 17
+ E E	0205 EE
+ 0 0	0206 00
+ 1 7	0207 17
+ A D	0208 AD
+ 0 0	0209 00
+ 1 7	020A 17
+ 4 9	020B 49
+ F F	020C FF
+ 4 A	020D 4A
+ A A	020E AA
+ C A	020F CA
+ 1 0	0210 10
+ F D	0211 FD
+ 3 0	0212 30
+ F 1	0213 F1

Key Sequences: Enter Program
FIGURE 5.5

5.4 EXECUTING THE PROGRAM

With the program entered, you may proceed to program execution. First, if the NMI vector has not been defined previously, enter the vector as follows:

<u>Press Keys</u>	<u>See Displayed</u>
<input type="text" value="AD"/> <input type="text" value="1"/> <input type="text" value="7"/> <input type="text" value="F"/> <input type="text" value="A"/>	17FA xx
<input type="text" value="DA"/> <input type="text" value="0"/> <input type="text" value="0"/>	17FA 00
<input type="text" value="+"/> <input type="text" value="1"/> <input type="text" value="C"/>	17FB 1C

This procedure insures that the key will be effective in terminating the program. Now, select the starting address of your program (0200) and begin execution as follows:

<u>Press Keys</u>	<u>See Displayed</u>
<input type="text" value="AD"/> <input type="text" value="0"/> <input type="text" value="2"/> <input type="text" value="0"/> <input type="text" value="0"/>	0200 A9
<input type="text" value="GO"/>	(Dark)

The program will now execute. If your seven selector switches all are open, you will probably hear no sound from the speaker because the square wave frequency is too high. If all selector switches are closed, you will hear in the speaker the lowest frequency that can be generated with the program as currently written. You may experiment with other combinations of switch settings to hear a variety of tones from the speaker.

Depression of the key will cause the program execution to stop (the tone will terminate) and the KIM-1 display will relight. The display will show the address and data for the next instruction to be executed (probably 020F or 0210 since this is the delay loop where the program spends most of its running time).

5.5 PROGRAM DEBUGGING AND MODIFICATION

If your program did not execute correctly, you would follow a debugging procedure involving the following steps:

Step 1: List the Program

First make sure you have entered the program steps correctly. Select the starting address () and observe that the correct data (A9) is displayed. Now, using the key, step through the remaining program locations checking for the correct data stored in each location.

Step 2: Single Step the Program

Follow the procedures listed in Section 5-4 for program execution but before depressing the key, place the SST slide switch in the ON position. Now, press the key and the first instruction will be executed. The display will relight indicating that the operating program is again in control of the system. The address displayed will be the address of the first byte of the next instruction to be executed. You may press the key again to execute the next instruction or you may choose to investigate changes in the contents of machine registers stored in selected memory locations (See Figure 3-13). The procedure detailed in Figure 5-6 gives a good indication of the various operations you may wish to perform in the SST mode.

Step 3: Check the I/O Operations


If program entry has been verified and program execution in the SST mode appears to be normal, you may wish to verify the correct operation of your specific I/O configuration.

You should recall that writing to or reading from any I/O port is the same as reading from or writing to any other memory location in the system. Therefore, if you select the address of an I/O port, the KIM-1 display will show you the hex code for the data being read from that address and thus, directly indicate the state of each I/O pin in the port. For example, the

address of the I/O port used for your sample program is 1700. Press and the display will show the hex code corresponding to the settings of your selector switches. If you change the positions of your selector switches, you will see the hex code change in the data field of the display.

Now, leave the same address (1700) selected and press the key. If you press any of the hex keys to , you will write the data to the I/O port (1700). Since seven of the pins of this I/O port are defined as inputs, only one (PA0) will act as an output and will respond to the data entered by you from the keyboard. Try alternating rapidly between the and keys and you should hear clicking in the speaker indicating that you are successfully toggling the PA0 pin.

This concept of using the KIM-1 keyboard and display to exercise and verify the operation of I/O ports is a generally useful technique for debugging the hardware portions of most specific applications.

<u>Press Keys</u>	<u>See Displayed</u>	<u>Comments</u>
AD 0 2 0 0	0200 A9	Select first instruction address
SST  0	0200 A9	Set SST to ON; All selector switches open
GO	0202 8D	Accumulator now loaded with \$01
GO	0205 EE	PADD now loaded
GO	0208 AD	PA0 now toggled
GO	020B 49	Switch values (PA1-PA7) now loaded
GO	020D 4A	Accumulator now complemented
GO	020E AA	Accumulator now right shifted 1 Bit
AD 0 0 F 3	00F3 xx	Display Accumulator
+	00F4 xx	Display Y - INDEX
+	00F5 00	Display X - INDEX
PC	020E AA	Restore PC (TAX will execute next)
GO	020F CA	Accumulator now loaded in X-INDEX
AD 0 0 F 3	00F3 00	Display Accumulator
+	00F4 xx	Display Y-INDEX
+	00F5 00	Display X-INDEX (A=0→X)
PC	020F CA	Restore PC
GO	0210 10	DEX now completed
AD 0 0 F 5	00F5 FF	Display X-INDEX (X<0)
PC	0210 10	Restore PC
GO	0212 30	No branch (Result of DEX <u>not</u> positive)
GO	0205 EE	Branch (Result of DEX <u>is</u> negative).

SST Mode: Sample Operation
FIGURE 5.6

CHAPTER 6

EXPANDING YOUR SYSTEM

In earlier sections you have learned that the MCS 6502 Microprocessor Array is capable of directly addressing up to 65,536 locations (bytes) of memory. (Usually abbreviated to 65K where "K" for the remainder of this section is to mean 1024 memory locations). In this section, we will discuss first the techniques for adding memory or I/O locations to the system and next, the proper handling of interrupt vectors in an expanded system.

6.1 MEMORY AND I/O EXPANSION

In the KIM-1 system, the management of input/output data is handled exactly the same as transfers to or from any other memory location in the system. There are no instructions dealing specifically with input/output transfers. Instead, transfer of data is accomplished by reading from or writing to registers connected to the data bus and to I/O pins in specific I/O interface devices (such as the 6530 array). These registers have a specific address in the system just as does any other memory location. Therefore, when we speak of expanding the memory of the KIM-1 system, we are defining the methods for expanding both the real memory (RAM, ROM, PROM, etc.) as well as the I/O ports since they are both treated exactly alike as far as address assignments are concerned.

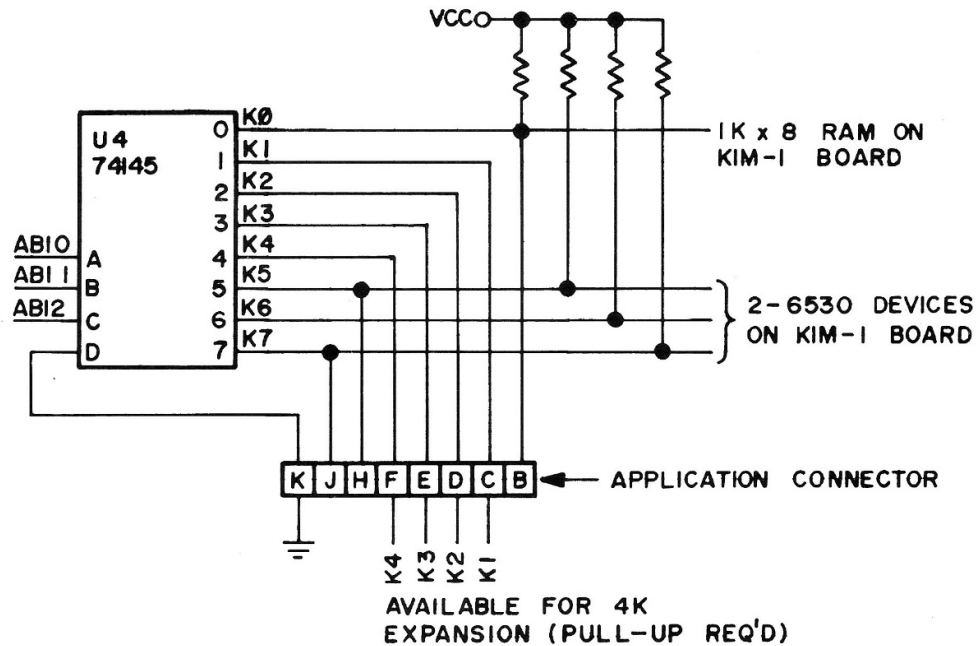
The first and most easily implemented memory expansion is the addition of up to 4K of memory space. You will recall that the lowest 8K memory locations are defined by an address decoder included on the KIM-1 module, (Device U4 on the schematic). The eight outputs of this decoder (K0 to K7) each define a 1K block of addresses in the lowest 8K of the memory map. Three of the outputs (K5, K6, K7) are used to select ROM, RAM, I/O and Timer locations on the two 6530 arrays while a fourth (K0) is used to select the 1024 locations of the static RAM memory. The remaining four outputs (K1, K2, K3, K4) are not used on the KIM-1 module but instead, are brought out to the Expansion connector for use as chip selects for memory or I/O additions.

Figure 6-1 shows the proper method for deriving the four chip select signals for the additional 4K of memory. Note that one of input pins of the decoder (D) was brought out to the Application Connector. It was this pin which we asked you to connect to ground in Chapter 2 of this manual. As long as this point remains connected to ground, the decoder will always select the lowest 8K addresses of the memory field regardless of the state of AB13, AB14, and AB15.

If you wish to expand the memory and I/O address space beyond the lower 8K addresses, you must arrange to de-select the lower 8K memory block while selecting some other 8K block. One suggested method for expanding beyond the lower 8K space is shown in Figure 6-2.

Note that the three high order address bits (AB13, AB14, AB15) are connected to a decoder. The eight outputs of the decoder act to divide the total 65K memory space into eight blocks of 8K each (8K0, 8K1, etc.). Now, the 8K0 output may be returned as the fourth input (D) to the decoder (U4) on the KIM-1 module causing the proper selection and de-selection of this block within the total address space. The remaining seven outputs (8K1 to 8K7) may be used to select and de-select the additional decoders shown in Figure 6-2. You need add only as many decoders (one for each 8K block of memory) as you need for your desired memory expansion.

A word of caution is in order when you decide to add memory to your system. You have noticed the inclusion of the line receivers for the AB10, AB11, and AB12 signals, (See Figure 6-2). These devices are included because of loading limitations placed on the address bus lines of the 6502 array (Each such line is capable of driving one standard TTL load and 130pf of capacity. See Appendix G).



4K Expansion
FIGURE 6.1

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Before deciding how to expand your system, we recommend a careful study of all of the loading limitations of the KIM-1 signals since almost certainly you will require additional buffering circuits if correct operation is to be achieved.

6.2 INTERRUPT VECTOR MANAGEMENT

We have referred several times in earlier sections to the interrupt features of the 6502 Microprocessor Array. We suggest now a careful reading of Section 9 of the Programming Manual for the subject "Reset and Interrupt Considerations".

In summary, there are three possible types of interrupt: Reset, NMI, and IRQ. Each will occur in response to an activation of one of the three pins of the 6502 array ($\overline{\text{RST}}$, $\overline{\text{NMI}}$, $\overline{\text{IRQ}}$). In response to these inputs, the 6502 array will fetch the data stored at a specific pair of addresses and load the data fetched into the program counter. The addresses are hardware determined and not under the control of the programmer. The specific addresses for each type of interrupt are:

FFFA, FFFB - $\overline{\text{NMI}}$ Vector

FFFC, FFFD - $\overline{\text{RST}}$ Vector

FFFE, FFFF - $\overline{\text{IRQ}}$ Vector

You will note that these addresses define the highest six locations in the 65K memory map.

In the KIM-1 system, three address bits (AB13, AB14, AB15) are not decoded at all. Therefore, when the 6502 array generates a fetch from FFFC and FFFD in response to a $\overline{\text{RST}}$ input, these addresses will be read as 1FFC and 1FFD and the reset vector will be fetched from these locations. You now see that all interrupt vectors will be fetched from the top 6 locations of the lowest 8K block of memory which is the only memory block decoded for the unexpanded KIM-1 system.

It is typical in any system to store the interrupt vectors in ROM so that they are immediately available after power-on. However, it is desirable that for the $\overline{\text{NMI}}$ and $\overline{\text{IRQ}}$ interrupts, the programmer be allowed to define as a variable the exact vector to which these interrupts will direct the system. Accordingly, the $\overline{\text{NMI}}$ and $\overline{\text{IRQ}}$ vector locations contain an indirect jump instruction referencing a RAM location into which the programmer will store the specific vector for the two types of interrupt. In the KIM-1 system, locations 17FA and 17FB contain the actual $\overline{\text{NMI}}$ vector and 17FE with 17FF contain the actual $\overline{\text{IRQ}}$ vector. The $\overline{\text{RST}}$ vector is not handled in this manner and always directs the system to the first step of the power-on initialization routine.

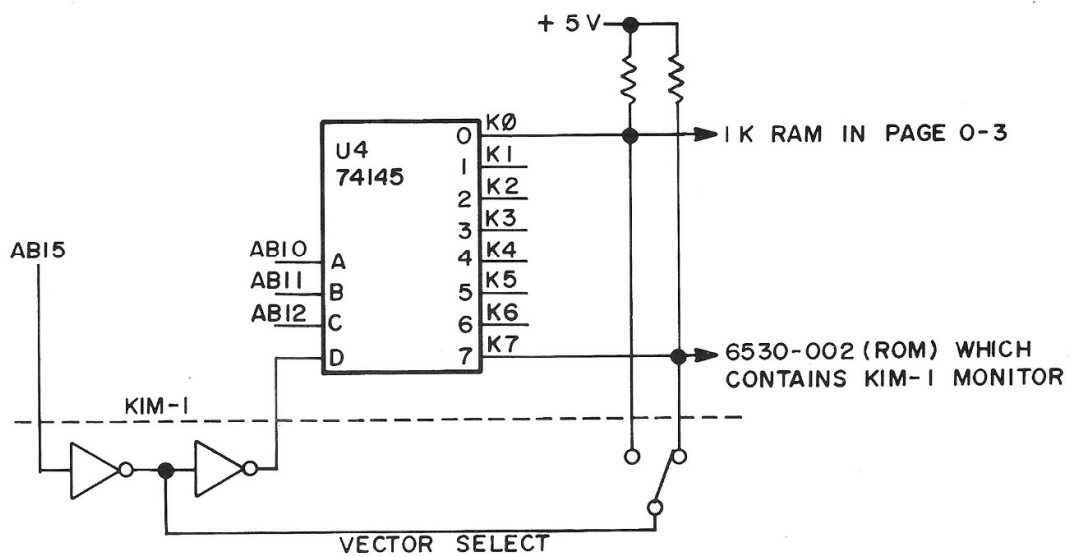
But what happens if we expand our memory above the lowest 8K block included in the KIM-1 system? Recall that we now must use AB13, AB14, and AB15 to decode the additional address locations of the memory. By so doing, the interrupt vector locations are no longer located in the K7 memory block since the decoder (U4) is de-selected in response to the addresses generated by the 6502 array in fetching the interrupt vectors (FFFA for example). We would have the same problem even in an unexpanded system if we wished to use a $\overline{\text{RST}}$ vector and initialization routine different than what the KIM-1 system provides and if the $\overline{\text{RST}}$ vector was to be located in a 1K block lower than K7 (K0 for instance).

The solution to this dilemma is to generate logically a special signal for interrupt select. Referring to Figure 6-2, a special signal called "Vector Select" is created to define the highest 1K memory block (K65). The fetch of any interrupt vector will cause this signal to go low "Select". Assuming that the K65 state is not used to select RAM, this signal may be "wire-or'd" with any one of the other "K" signals (K0 to K64) to define exactly which 1K block is to contain the interrupt vectors.

As an example, assume that you have connected the K65 "Vector Select" line to the K0 line. When a $\overline{\text{RST}}$ occurs, the 6502 array generates a fetch from locations FFFC and FFFD. These addresses cause K65 to be selected which, in turn, accesses the K0 field of the memory and causes the actual fetch of the $\overline{\text{RST}}$ vector from locations 03FC and 03FD. (Had you chosen to connect K65 to K7, the fetch of the reset vectors would occur from locations 1FFC and 1FFD).

In this way, the highest six addresses of any 1K block of memory may be used to supply the interrupt vectors for the system. If desired, a switch could be installed to allow you to select different areas of memory as the source locations for the interrupt vectors. (By the way, we selected the 75145 type decoders in Figure 6-2 specifically to allow the "wire-or" of K65 with any other K. This is possible because the 75145 decoder is provided with open-collector outputs which allows "wire-or" of several states using an external load resistor.)

An even simpler arrangement using the "Vector Select" approach is shown in Figure 6-3. Here, the KIM-1 system is assumed to have only the lower 8K of memory in place. The address decoder (U4) is de-selected using the AB15 signal which becomes "true" whenever an interrupt vector fetch is initiated by the system. The same signal (AB15) is inverted and "wire-or'd" through a switch to the K0 or the K7 chip select lines. Now, depending upon the position of the switch, interrupt vectors will be fetched from the top 6 addresses of either block K0 or K7. K0 in the KIM-1 system is the RAM and K7 is the ROM in the 6530-002 array (the operating program). In this way, you may have two different sets of interrupt vectors in your system and may select which set is to be used with a simple switch.



Vector Selection
FIGURE 6.3

CHAPTER 7

WARRANTY AND SERVICE

Should you experience difficulty with your KIM-1 module and be unable to diagnose or correct the problem, you may return the unit to MOS Technology, Inc. for repair.

7.1 IN-WARRANTY SERVICE

All KIM series Microcomputer Modules are warranted by MOS Technology, Inc. against defects in workmanship and materials for a period of ninety (90) days from date of delivery. During the warranty period, MOS Technology, Inc. will repair or, at its option, replace at no charge components that prove to be defective provided that the module is returned, shipping prepaid, to:

KIM Customer Service Department
MOS Technology, Inc.
950 Rittenhouse Road
Norristown, Pennsylvania 19401

This warranty does not apply if the module has been damaged by accident or misuse, or as a result of repairs or modifications made by other than authorized personnel at the above captioned service facility.

No other warranty is expressed or implied. MOS Technology, Inc. is not liable for consequential damages.

7.2 OUT-OF-WARRANTY SERVICE

Beyond the ninety (90) day warranty period, KIM modules will be repaired for a reasonable service fee. All service work performed by MOS Technology, Inc. beyond the warranty period is warranted for an additional ninety (90) day period after shipment of the repaired module.

7.3 POLICY ON CHANGES

All KIM series modules are sold on the basis of descriptive specifications in effect at the time of sale. MOS Technology, Inc. shall have no obligation to modify or update products once sold. MOS Technology, Inc. reserves the right to make periodic changes or improvements to any KIM series module.

7.4 SHIPPING INSTRUCTIONS

It is the customer's responsibility to return the KIM series module with shipping charges prepaid to the above captioned service facility.

For in-warranty service, the KIM module will be returned to the customer, shipping prepaid, by the fastest economical carrier.

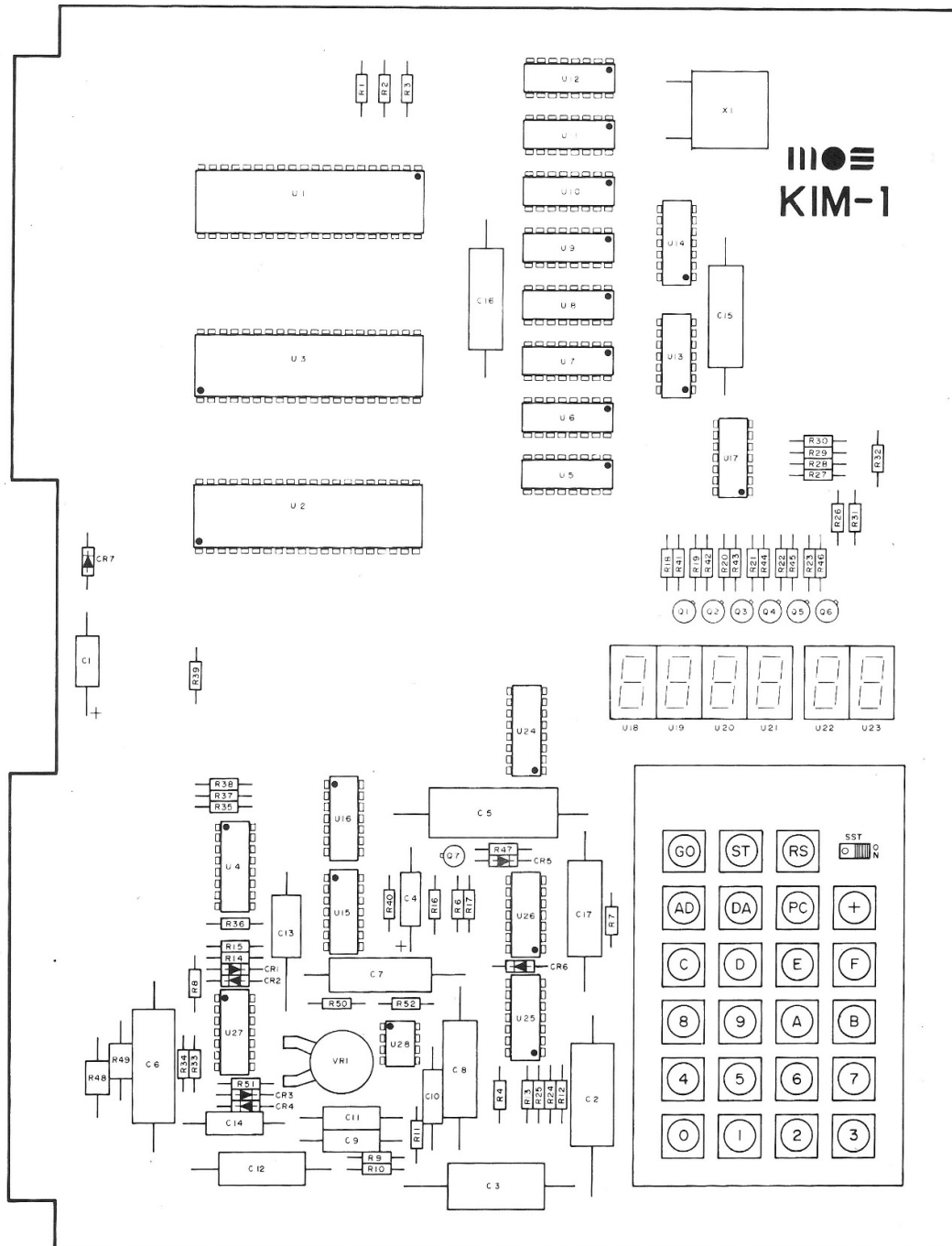
For out-of-warranty service, the customer will pay for shipping charges both ways. The repaired KIM module will be returned to the customer C.O.D. unless the repairs and shipping charges are prepaid by the customer.

Please be certain that your KIM module is safely packaged when returning it to the above captioned service facility.

APPENDIX A

ITEM	PART	QTY.	DESCRIPTION
1.	U1	1	6502 Microprocessor
2.	U2	1	6530 ROM RAM I/O Chip-02
3.	U3	1	6530 ROM RAM I/O Chip-03
4.	U5 through U12	8	6102 RAM 500ns Acc, 0ns
5.	U18 through U23	6	7 SEG .3" Red Display
6.	U25	1	556 Timer IC
7.	U27	1	565 Phase Lock Loop
8.	U28	1	311 Comparator
9.	U4, U24	2	74145 BCD Decoder IC
10.	U13 & U14	2	74125 TRI STATE Buffer
11.	U15	1	7400 Quad Nand IC
12.	U16	1	7404 Hex Inverter IC
13.	U17	1	7406 Hex Inv. O/C IC
14.	U26	1	7438 Quad Nand O/C IC
15.	CR1, CR2, CR3 & CR4	4	20 MA. 50v Diode - IN914
16.	CR5, CR6	2	1A 50v Diode - IN4001
17.	CR7	1	6.2v 1/2w Z. Diode - IN4735
18.	Q7	1	NPN Transistor B>20, VCE>12 - 2N5371
19.	Q1 through Q6	6	PNP Transistor B>20, VCE>6 - 2N5375
20.	R24 & R25	2	47KΩ ±10% 1/4w Resistor
21.	R1, R2, R3, R4, R6	5	3.3KΩ ±10% 1/4w Resistor
22.	R34 & R50	2	2.2KΩ ±10% 1/4w Resistor
23.	R12-R17, R41-R46	12	1.0KΩ ±10% 1/4w Resistor
24.	R35 through R40	6	560Ω ±10% 1/4w Resistor
25.	R18-R23, R47	7	220Ω ±10% 1/4w Resistor
26.	R33	1	47Ω ±10% 1/4w Resistor
27.	R52	1	5 Meg. ±10% 1/4w Resistor
28.	R51	1	30KΩ ±5% 1/4w Resistor
29.	R7, R8, R9, R10 & R11	5	10KΩ ±5% 1/4w Resistor
30.	R48, R49	2	150Ω ±5% 1/2w
31.	R26 through R32	7	82Ω ±5% 1/4w
32.	VR1	1	5KΩ Potentiometer
33.	C2, C3, C6	3	.22±10% uf.>12 wv. cap.
34.	C1, C4	2	1uf+80-10%>12WV Cap
35.	C5	1	.33 uf±10%>12WV Cap
36.	C7, C8, C15, C16, C17	5	.1uf+80-10%>12WV Cap
37.	C9, C10, C11	3	.0068uf±10%>12WV
38.	C12	1	.047uf±10%>12WV
39.	C13	1	.022uf±10%>12WV
40.	C14	1	.001uf±10%>12WV
41.		1	44 Pin Edge Conn.
42.	XL	1	1 MHz XTAL
43.		1	PCB.
44.		1	24 Key KBD
45.		6	Rubber Pads
46.		1	Shipping Bag (Static Free)
47.		1	Shipping Box
48.		1	Hardware Manual
49.		1	Software Manual
50.		1	KIMManual
51.		1	Warranty Card
52.		1	Wall Chart
53.		2	#2 x 1/4 SS Screws (Keyboard)
54.		1	Program Card

APPENDIX B
KIM-1 PARTS LAYOUT



APPENDIX C

IN CASE OF TROUBLE

SYMPTOM: Display Not Lit

1. Test +5 volt power supply. Using a VOM check for +5 volts between Pin E-21 and E-22. Also check for +5 volts between Pin A-A and Pin A-1. KIM-1 power supply should be set at $+5v \pm 5\%$.
2. Test KB/TTY option wiring (Figure 2-4). Pin A-21 should not be connected to Pin A-V.
3. Make sure decoder is enabled. See Figure 2-2 and insure that Pin A-K is connected to ground.
4. Depress the reset key and check all other keys to insure that no key is stuck.
5. Place a VOM between Pin E-21 (+5v) and Pin E-7 (Reset). Alternately depress and release the reset key checking to see if the voltage swings from ($>4v$) to ($<1v$).
6. Test Pin E-V (\emptyset_2) with an oscilloscope and insure 1 MHz operation.

SYMPTOM: Cannot Dump to Audio Tape
Cannot Load From Audio Tape

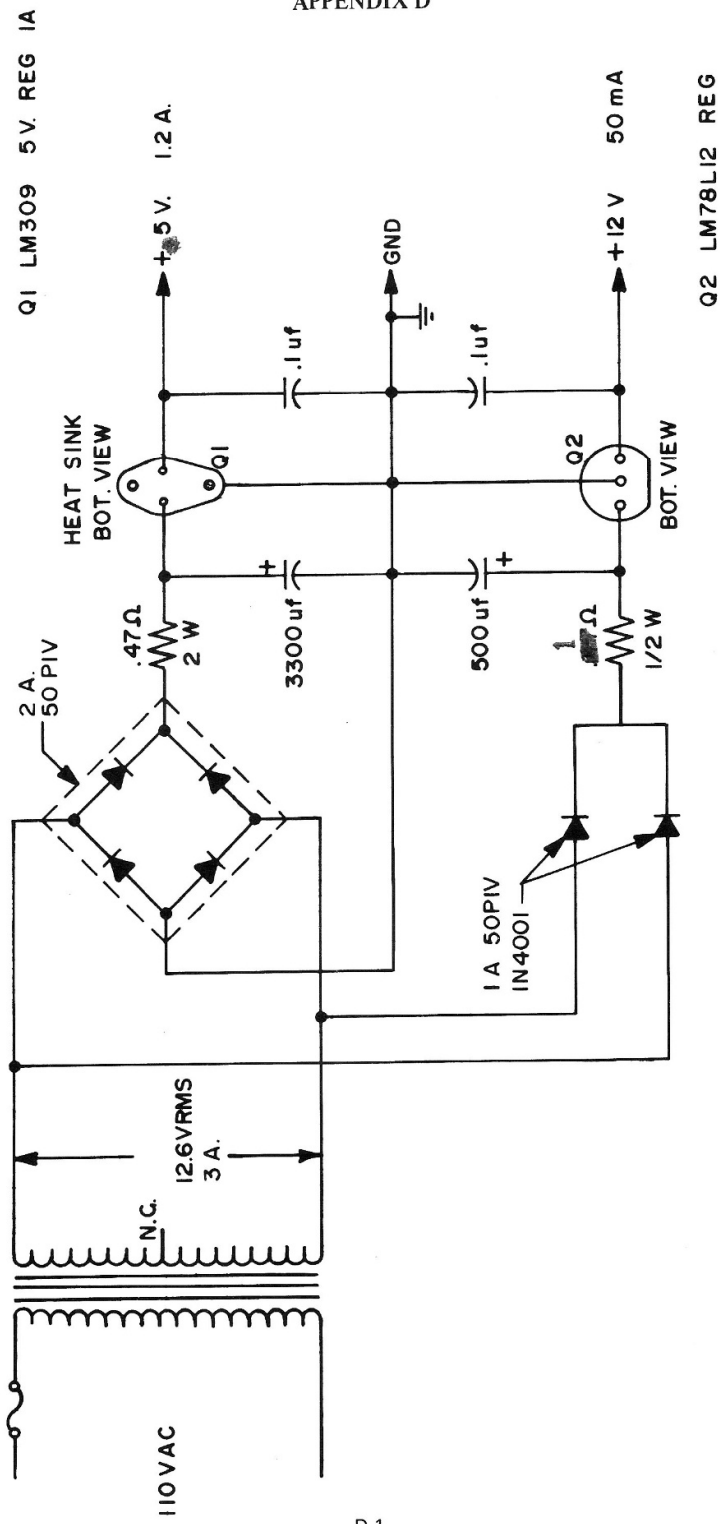
1. Test +12 volt power supply. Using a VOM check for +12 volts between Pin A-N (+12v) and Pin A-1 (GND). Set power supply to $+12v \pm 5\%$. (See Figure 2-2).
2. Check volume control on the tape recorder (Set at half way point).

3. Make sure that you are using the proper tape output pin.
See Figure 2-3.
4. Check the tape interface circuit by disconnecting the tape recorder and shorting Pin A-B (Audio Out High) to Pin A-L (Audio In). Set up KIM-1 monitor to dump a section of memory. Using an oscilloscope observe data at Pin E-X (PLL TEST). See Appendix E for correct data format and calibration procedure.
5. Record voice on a section of tape and play it back to insure that the tape recorder is working. Connect another tape recorder to the system or try another cassette.

SYMPTOM: TTY Interface Problems

1. Make sure that Pin A-21 is connected to Pin A-V (Figure 2-4) to allow TTY operation.
2. Compare the connections on Figure 2-4 with interface schematics in your TTY manual (or any other serial teleprinter).
3. Depress the reset key on the KIM-1 keyboard followed by a rub out character from the TTY.

APPENDIX D



Suggested Power Supply

APPENDIX E

AUDIO TAPE FORMAT

Data is stored out onto your audio cassette recorder in a specific format designed to insure an error free recovery. In the unlikely event that a playback error does occur, several "ERROR DETECTION" methods are incorporated to warn you of this condition.

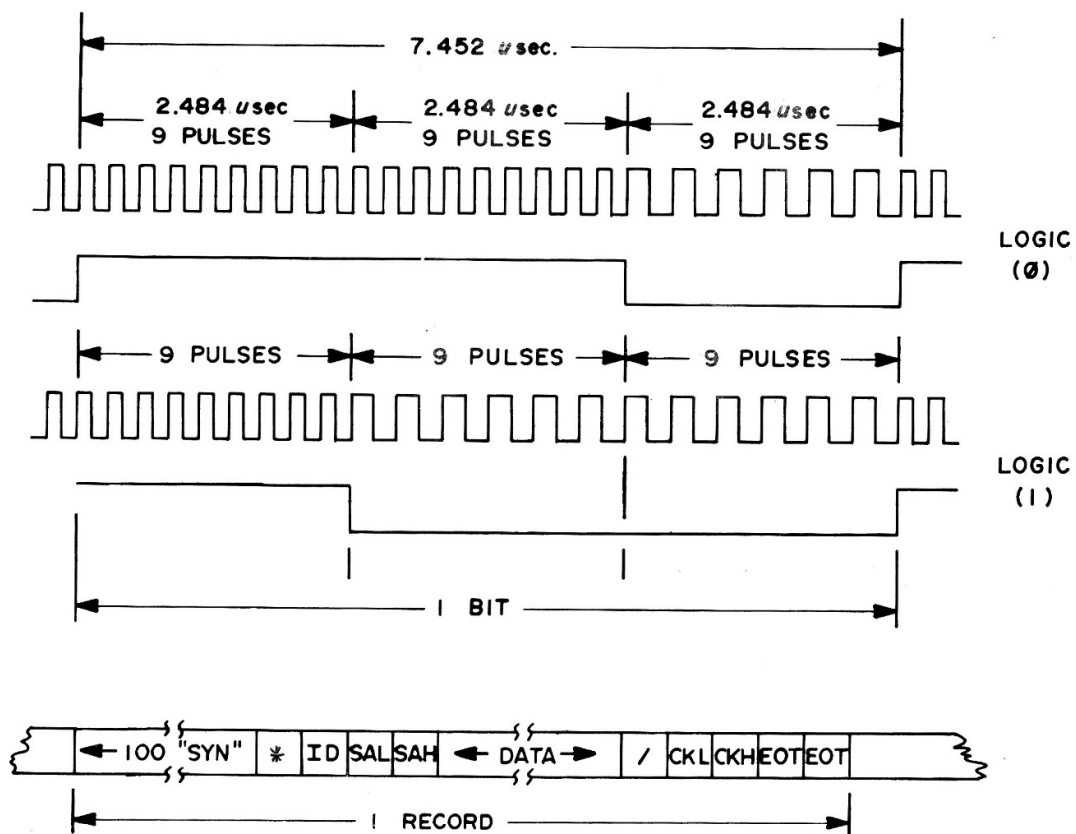
Data is transmitted to the tape recorder in the form of serial "ASCII" encoded characters (seven data bits plus Parity bit). Data retrieved from the memory is converted into this form by separating each byte into two half bytes. The half bytes are then converted into their ASCII equivalents.

Each record transmitted begins with a leader of one hundred "SYN" characters (ASCII 16) followed by a * character (ASCII 2A). During playback, this pattern allows your micro-computer to detect the start of a valid data record and synchronize to the serial data stream. Following the *, the record identification number (ID), and starting address low (SAL) and the starting address high (SAH) are transmitted. The data specified by the starting (SAL, SAH) and ending limits (EAL, EAH) is transmitted next followed by a "/" character (ASCII 2F) to indicate the end of the data portion of the record. Following the "/" two "CHECK-SUM" bytes are transmitted for comparison with a calculated check-sum number during playback to further insure that a proper data retrieval has taken place. Two "EOT" characters (ASCII 04) mark the end of record transmission.

Each transmitted bit begins with a 3700 hertz tone and ends with a 2400 hertz tone. "Ones" have the high to low frequency transition at one-third of the bit period. "Zeros" have the transition at two-thirds of the period. During playback the 565 phase locked loop locks to, and tracks these two frequencies producing (through the 310 comparator) a logic "1" pulse of one-third the bit period for a "One". A pulse two thirds the bit period is likewise produced for a "Zero". Your microcomputer uses a software controlled algorithm for converting this signal into eight bit data words.

The frequency shift keyed phase lock loop method of data recovery is relatively insensitive to amplitude and phase variations. The "FREE RUNNING" frequency of the phase lock loop has been adjusted at the factory to a frequency half way between the two data frequencies (called the Center Frequency). This adjustment is accomplished by strapping Pin A-P (Audio Out High) to Pin A-L (Audio In). A program starting at address 1A6B_{HEX} provides the center frequency reference that allows the loop to be adjusted by potentiometer VR1. Pin E-X (PLL TEST) is monitored with a voltmeter while the pot is rotated until the voltmeter reading is at the transition point between a logical "1" (+5v) and "0" (GND).

THIS ADJUSTMENT HAS BEEN FACTORY PRESET AND SHOULD ONLY REQUIRE ADJUSTMENT DUE TO COMPONENT REPLACEMENT!



Audio Tape Format
FIGURE E-1

APPENDIX F

PAPER TAPE FORMAT

The paper tape LOAD and DUMP routines store and retrieve data in a specific format designed to insure error free recovery. Each byte of data to be stored is converted to two half bytes. The half bytes (whose possible values are 0 to F_{HEX}) are translated into their ASCII equivalents and written out onto paper tape in this form.

Each record outputted begins with a ";" character (ASCII 3B) to mark the start of a valid record. The next byte transmitted (18_{HEX}) is the number of data bytes contained in the record. The record's starting address low (1 byte, 2 characters), starting address Hi (1 byte, 2 characters), and data (18 bytes, 36 characters) follow. Each record is terminated by the record's check-sum (2 bytes, 4 characters), a carriage return (ASCII 0D), line feed (ASCII 0A), and six "NULL" characters (ASCII 00).

The last record transmitted has zero data bytes (indicated by ;00). The starting address field is replaced by a four digit Hex number representing the total number of data records contained in the transmission, followed by the records usual check-sum digits. A "XOFF" character ends the transmission.

;180000FFEEDDCCBBAA0099887766554433221122334455667788990AFC
;0000010001

During a "LOAD" all incoming data is ignored until a ";" character is received. The receipt of non ASCII data or a mismatch between a records calculated check-sum and the check-sum read from tape will cause an error condition to be recognized by KIM. The check-sum is calculated by adding all data in the record except the ";" character.

The paper tape format described is compatible with all other MOS Technology, Inc. software support programs.

APPENDIX G

6502 CHARACTERISTICS

Clocks (\emptyset_1 , \emptyset_2)

The MCS 6502 is supplied with an internal clock generator. The frequency of this clock is crystal controlled.

Address Bus (A_0 - A_{15})

These outputs are TTL compatible, capable of driving one standard TTL load and 130pf.

Data Bus (D_0 - D_7)

Eight pins are used for the data bus. This is a bi-directional bus, transferring data to and from the device and peripherals. The outputs are tri-state buffers capable of driving one standard TTL load and 130pf.

Ready (RDY)

This input signal allows the user to single cycle the microprocessor on all cycles except write cycles. A negative transition to the low state during or coincident with phase one (\emptyset_1) will halt the microprocessor with the output address lines reflecting the current address being fetched. This condition will remain through a subsequent phase two (\emptyset_2) in which the Ready signal is high. This feature allows microprocessor interfacing with low speed PROMS as well as fast (max. 2 cycle) Direct Memory Access (DMA). If Ready is low during a write cycle, it is ignored until the following read operation.

Interrupt Request ($\overline{\text{IRQ}}$)

This TTL level input requests that an interrupt sequence begin within the microprocessor. The microprocessor will complete the current instruction being executed before recognizing the request. At that time, the interrupt mask bit in the Status Code Register will be examined. If the interrupt mask flag is not set, the microprocessor will begin an interrupt sequence. The Program Counter and Processor Status Register are stored in the stack. The microprocessor will then set the interrupt mask flag high so that no further interrupts may occur. At the end of this cycle, the program counter low will be loaded from address FFFE, and program counter high from location FFFF, therefore transferring program control to the memory vector located at these addresses. The RDY signal must be in the high state (for control to the memory vector) located at these addresses. The RDY signal must be in the high state for any interrupt to be recognized. A 3K Ω external register should be used for proper wire-OR operation.

Non-Maskable Interrupt ($\overline{\text{NMI}}$)

A negative going edge on this input requests that a non-maskable interrupt sequence be generated within the microprocessor.

$\overline{\text{NMI}}$ is an unconditional interrupt. Following completion of the current instruction, the sequence of operations defined for $\overline{\text{IRQ}}$ will be performed, regardless of the state of the interrupt mask flag. The vector address loaded into the program counter, low and high, are locations FFFA and FFFB respectively. The instructions loaded at these locations causes the microprocessor to branch to a non-maskable interrupt routine in memory.

$\overline{\text{NMI}}$ also requires an external 3K Ω resistor to Vcc for proper wire-OR operations.

Inputs $\overline{\text{IRQ}}$ and $\overline{\text{NMI}}$ are hardware interrupts lines that are sampled during ϕ_2 (phase 2) and will begin the appropriate interrupt routine on the ϕ_1 (phase 1) following the completion of the current instruction.

Set Overflow Flag (S.O.)

This TTL level input signal allows external control of the overflow bit in the Status Code Register.

SYNC

This output line is provided to identify those cycles in which the microprocessor is doing an Op Code fetch. The SYNC line goes high during ϕ_1 of an Op Code fetch and stays high for the remainder of that cycle. If the RDY line is pulled low during the ϕ_1 clock pulse in which SYNC went high, the processor will stop in its current state and will remain in the state until the RDY line goes high. In this manner, the SYNC signal can be used to control RDY to cause single instruction execution.

RESET

This input is used to reset or start the microprocessor from a power down condition. During the time that this line is held low, writing to or from the microprocessor is inhibited. When a positive edge is detected on the input, the microprocessor will immediately begin the reset sequence.

After a system initialization time of six clock cycles, the mask interrupt flag will be set and the microprocessor will load the program counter from the memory vector locations FFFC and FFFD. This is the start location for program control.

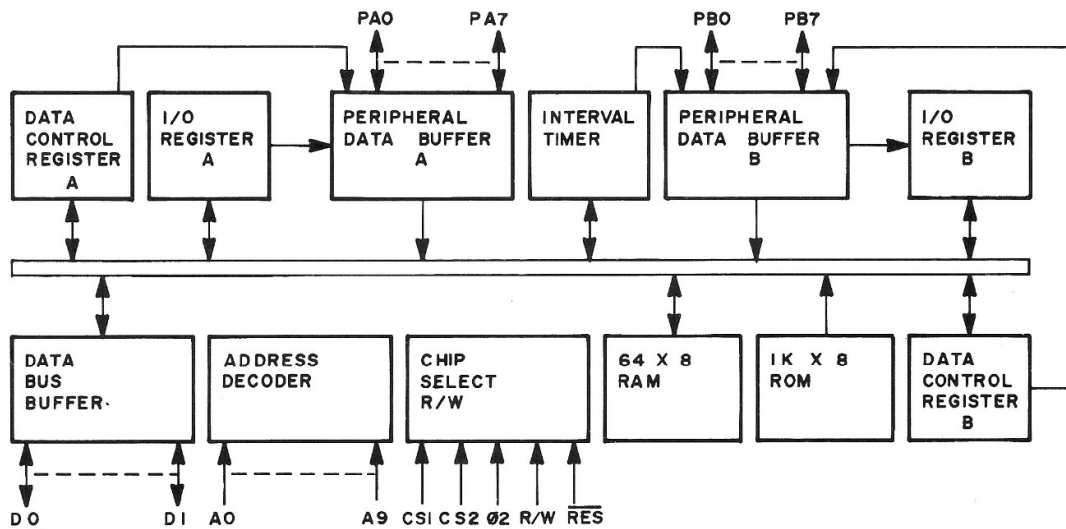
After Vcc reaches 4.75 volts in a power up routine, reset must be held low for at least two clock cycles.

When the reset signal goes high following these two clock cycles, the microprocessor will proceed with the normal reset procedure detailed above.

APPENDIX H

6530 CHARACTERISTICS

The MCS 6530 is designed to operate in conjunction with the MCS 650X Microprocessor Family. It is comprised of a mask programmable 1024 x 8 ROM, a 64 x 8 static RAM, two software controlled 8 bit bi-directional data ports allowing direct interfacing between the microprocessor unit and peripheral devices, and a software programmable interval timer with interrupt, capable of timing in various intervals from 1 to 262,144 clock periods.



MCS 6530 Block Diagram
FIGURE H.1

Reset ($\overline{\text{RES}}$)

During system initialization a Logic "0" on the $\overline{\text{RES}}$ input will cause a zeroing of all four I/O registers. This in turn will cause all I/O buses to act as inputs thus protecting external components from possible damage and erroneous data while the system is being configured under software control. The Data Bus Buffers are put into an OFF-STATE during Reset. Interrupt capability is disabled with the $\overline{\text{RES}}$ signal. The $\overline{\text{RES}}$ signal must be held low for at least one clock period when reset is required.

Input Clock

The input clock is a system Phase Two clock which can be either a low level clock ($V_{\text{IL}} < 0.4$, $V_{\text{IH}} > 2.4$) or high level clock ($V_{\text{IL}} < 0.2$, $V_{\text{IH}} = V_{\text{CC}}^{+.3}_{-.2}$).

Read/Write (R/W)

The R/W signal is supplied by the microprocessor array and is used to control the transfer of data to and from the microprocessor array and the MCS 6530. A high on the R/W pin allows the processor to read (with proper addressing) the data supplied by the MCS 6530. A low on the R/W pin allows a write (with proper addressing) to the MCS 6530.

Interrupt Request ($\overline{\text{IRQ}}$)

The $\overline{\text{IRQ}}$ pin is an interrupt pin from the interval timer. This same pin, if not used as an interrupt, can be used as a peripheral I/O pin (PB7). When used as an interrupt, the pin should be set up as an input by the data direction register. The pin will be normally high with a low indicating an interrupt from the MCS 6530.

Data Bus (D0-D7)

The MCS 6530 has eight bi-directional data pins (D0-D7). These pins connect to the system's data lines to allow transfer of data to and from the microprocessor array. The output buffers remain in the off state except when a Read operation occurs.

Peripheral Data Ports

The MCS 6530-002, MCS 6530-003 both have 15 pins available for peripheral I/O operations. Each pin is individually software programmable to act as either an input or an output. The 15 pins are divided into 2 8-bit ports, PA0-PA7 and PB0-PB7. PB6 was used as a chip select and is not available to the user. The pins are set up as an input by writing a "0" into the corresponding bit of the data direction register. A "1" into the data direction register will cause its corresponding bit to be an output. When in the input mode, the peripheral output buffers are in the "1" state and a pull-up device acts as less than one TTL load to the peripheral data lines. On a Read operation, the microprocessor unit reads the peripheral pin. When the peripheral device gets information from the MCS 6530 it receives data stored in the data register. The microprocessor will read correct information if the peripheral lines are greater than 2.0 volts for a "1" and less than 0.8 volts for a "0" as the peripheral pins are all TTL compatible. Pins PA0 and PB0 are also capable of sourcing 3 ma at 1.5v, thus making them capable of Darlington drive. Pin PB7 has no internal pull-up (to allow collector-oring with other devices).

Address Lines (A0-A9)

There are 10 address pins. In addition to these 10, there is the ROM SELECT pin. The above pins, A0-A9 and ROM SELECT, are always used as addressing pins. There are 2 additional pins which are mask programmable and can be used either individually or together as CHIP SELECTS. They are pins PB5 and PB6. When used as peripheral data pins they cannot be used as chip selects. PB5 was used as a data pin while PB6 was used as a chip select and is not available to the user.

A block diagram of the internal architecture is shown in Figure H-1. The MCS 6530 is divided into four basic sections, RAM, ROM, I/O and TIMER. The RAM and ROM interface directly with the microprocessor through the system data bus and address lines. The I/O section consists of 2 8-bit halves. Each half contains a Data Direction Register (DDR) and an I/O Register.

ROM 1K Byte (8K Bits)

The 8K ROM is in a 1024 x 8 configuration. Address lines A0-A9, as well as RS0 are needed to address the entire ROM. With the addition of CS1 and CS2, seven MCS 6530's may be addressed, giving 7168 x 8 bits of contiguous ROM.

RAM 64 Bytes (512 Bits)

A 64 x 8 static RAM is contained on the MCS 6530. It is addressed by A0-A5 (Byte Select), RS0, A6, A7, A8, A9 and CS1.

Internal Peripheral Registers

There are four internal registers, two data direction registers and two peripheral I/O data registers. The two data direction registers (A side and B side) control the direction of the data into and out of the peripheral pins. A "1" written into the Data Direction Register sets up the corresponding peripheral buffer pin as an output. Therefore, anything then written into the I/O Register will appear on that corresponding peripheral pin. A "0" written into the DDR inhibits the output buffer from transmitting data to or from the I/O Register. For example, a "1" loaded into data direction register A, position 3, sets up peripheral pin PA3 as an output. If a "0" had been loaded, PA3 would be configured as an input and remain in the high state. The two data I/O registers are used to latch data from the Data Bus during a Write operation until the peripheral device can read the data supplied by the microprocessor array.

During a read operation the microprocessor is not reading the I/O Registers but in fact is reading the peripheral data pins. For the peripheral data pins which are programmed as outputs the microprocessor will read the corresponding data bits of the I/O Register. The only way the I/O Register data can be changed is by a microprocessor Write operation. The I/O Register is not affected by a Read of the data on the peripheral pins.

Interval Timer

The Timer section of the MCS 6530 contains three basic parts: preliminary divide down register, programmable 8-bit register and interrupt logic. These are illustrated in Figure 4.

The interval timer can be programmed to count up to 256 time intervals. Each time interval can be either 1T, 8T, 64T or 1024T increments, where T is the system clock period. When a full count is reached, an interrupt flag is set to a logic "1". After the interrupt flag is set the internal clock begins counting down to a maximum of -255T. Thus, after the interrupt flag is set, a Read of the timer will tell how long since the flag was set up to a maximum of 255T.

The 8 bit system Data Bus is used to transfer data to and from the Interval Timer. If a count of 52 time intervals were to be counted, the pattern 0 0 1 1 0 1 0 0 would be put on the Data Bus and written into the Interval Time register.

At the same time that data is being written to the Interval Timer, the counting intervals of 1, 8, 64, 1024T are decoded from address lines A₀ and A₁. During a Read or Write operation address line A₃ controls the interrupt capability of PB₇, i.e., A₃ = 1 enables IRQ on PB₇, A₃ = 0 disables IRQ on PB₇. When PB₇ is to be used as an interrupt flag with the interval timer it should be programmed as an input. If PB₇ is enabled by A₃ and an interrupt occurs PB₇ will go low. When the timer is read prior to the interrupt flag being set, the number of time intervals remaining will be read, i.e., 51, 50, 49, etc.

When the timer has counted down to 0 0 0 0 0 0 0 0 on the next count time an interrupt will occur and the counter will read 1 1 1 1 1 1 1 1. After interrupt, the timer register decrements at a divide by "1" rate of the system clock. If after interrupt, the timer is read and a value of 1 1 1 0 0 1 0 0 is read, the time since interrupt is 28T. The value read is in two's complement.

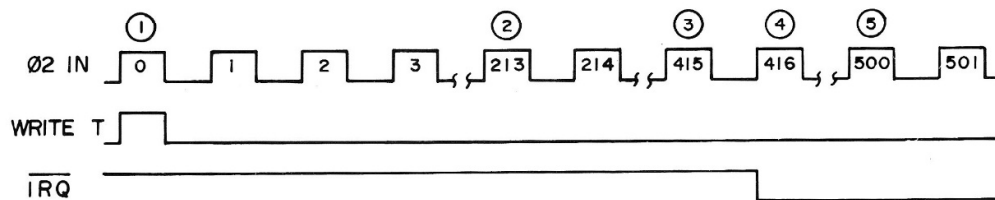
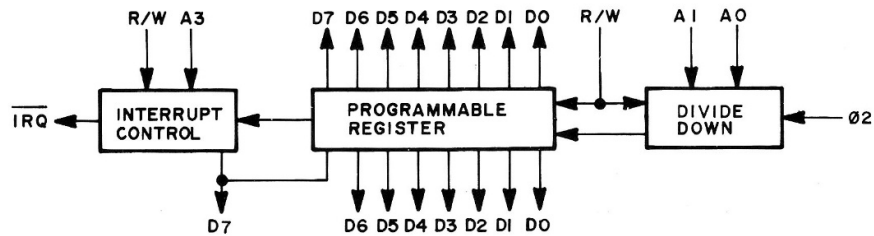
Value read = 1 1 1 0 0 1 0 0

Complement = 0 0 0 1 1 0 1 1

ADD 1 = 0 0 0 1 1 1 0 0 = 28.

Thus, to arrive at the total elapsed time, merely do a two's complement add to the original time written into the timer. Again, assume time written as 0 0 1 1 0 1 0 0 (=52). With a divide by 8, total time to interrupt is $(52 \times 8) + 1 = 417T$. Total elapsed time would be $416T + 28T = 444T$, assuming the value read after interrupt was 1 1 1 0 0 1 0 0.

After interrupt, whenever the timer is written or read the interrupt is reset. However, the reading of the timer at the same time the interrupt occurs will not reset the interrupt flag.



1. Data written into interval timer is $00110100 = 52_{10}$

2. Data in Interval timer is $00011001 = 25_{10}$

$$52 - \frac{213-1}{8} = 52-26-1 = 25$$

3. Data in Interval timer is $00000000 = 0_{10}$

$$52 - \frac{415-1}{8} = 52-51-1 = 0$$

4. Interrupt has occurred at ϕ_2 pulse #416

Data in Interval timer = 11111111

5. Data in Interval timer is 10101100

two's complement is $01010100 = 84_{10}$

$$84 + (52 \times 8) = 500_{10}$$

When reading the timer after an interrupt, A3 should be low so as to disable the $\overline{\text{IRQ}}$ pin. This is done so as to avoid future interrupts until after another Write timer operation.

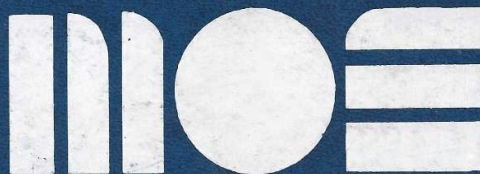
Basic Elements of Interval Timer
FIGURE H.2

HEADQUARTERS -
MOS TECHNOLOGY, INC. 950 Rittenhouse Road
Norristown, Pa. 19401, (215) 666-7950, TWX: 510/660/4033

EASTERN REGION -
Mr. William Whitehead
MOS TECHNOLOGY, INC., Suite 312,
410 Jericho Turnpike, Jericho, N.Y. 11753
(516) 822-4240

WESTERN REGION -
MOS TECHNOLOGY, INC. 2172 Dupont Drive,
Patio Bldg., Suite 221, Newport Beach, CA. 92660
(714) 833-1600

Mr. Petr Sehnal, Regional Applications Mgr.
MOS TECHNOLOGY, INC., 26921 Grasmere Place,
Hayward, CA. 94542
(415) 881-8080



MOSTECHNOLOGY INC