

# 8

## ***AN MCS6502-BASED MICROCOMPUTER—THE KIM-1***

### **INTRODUCTION**

This chapter is written for persons who need to become acquainted with the KIM-1<sup>®</sup> microcomputer. This chapter uses the programmed instruction approach. Beginning with a discussion of the actual power supply connections, the chapter then proceeds to discuss the function of each key on the KIM-1 keyboard. The function of each key is highlighted by tables showing keys pressed and the corresponding display. After the keyboard entry of a simple problem is discussed, the use of the Teletype<sup>®</sup> model ASR33 as an input/output terminal is presented. The use of the Teletype is discussed in detail with reproductions of actual printed outputs and punched paper tapes being provided. Next the use of an audio cassette tape recorder is presented. Examples illustrate the detailed step by step procedure for recording onto magnetic tape from the microcomputer, and vice versa. Simple programming examples are used in all cases so that the reader can focus his attention on the operation of a microcomputer at machine or assembly language levels.

This chapter provides detailed information on the interaction between a user and a microcomputer when assembly level (or machine) language is used. After studying this chapter the reader will be in a position to decide whether this detailed level of interaction is acceptable,

<sup>®</sup> KIM-1 is a registered trademark of the MOS Technology Corporation

or whether he should consider the use of a microcomputer equipped with a higher level language. A higher level language, such as BASIC, makes the computer much more transparent to the user. (The reader should be forewarned, however, that there are literally hundreds of different versions of BASIC, so purchase of a microcomputer with a resident BASIC does not solve all of the user's problems.)

## 8.1 WHAT IS A KIM-1?

A KIM-1 (Keyboard Input Monitor) is a microcomputer built around the 6502 microprocessor presently manufactured by MOS Technology and Rockwell. The KIM-1 presently (1977 - 1978) sells for a little over \$200 and is a complete microcomputer which includes the following:

- \* The microprocessor
- \* 2048 ROM bytes
- \* 1024 + 128 RAM bytes
- \* 30 input/output pins
- \* 2 timers
- \* Interface for audio cassette
- \* Interface for a Teletype
- \* 23 key keyboard
- \* Display consisting of 6 seven bar characters

The KIM-1 is an example of the so-called single-board microcomputers, because, except for the power supplies, it is fabricated on a single printed circuit board measuring approximately 20 by 28 cm, (8½ by 11 inches). Figure 8.1-1 is a photograph of the KIM-1 and the single 5 volt power supply which is required when using only the keyboard, a Teletype, or other compatible terminal. An additional 12 volt power supply is required if an audio cassette tape recorder is added. The KIM-1 circuit board is shown in Figure 8.1-2.

The KIM-1 is typical of the many microprocessor based microcomputers which are presently being sold. It is the authors' opinion that this chapter will be of more use to the reader if one microcomputer is presented in depth as opposed to discussing several microcomputers in more general terms. Somewhat like learning to drive an automobile, or to ride a bicycle, the carryover from the use of one microcomputer to another is large so the adjustment period in learning to use a different microcomputer is surprisingly short.



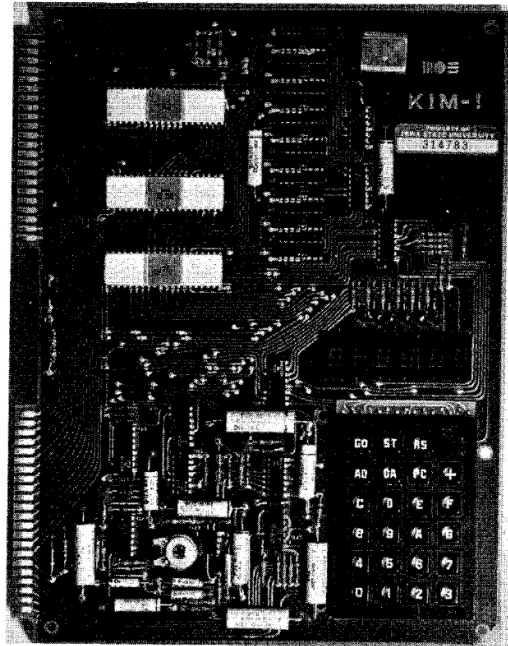
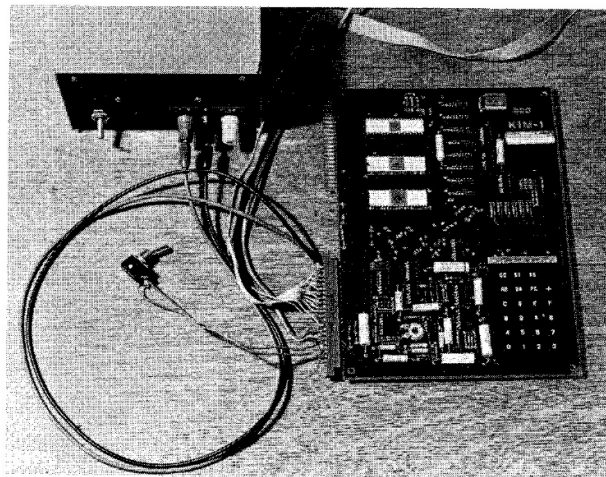


Figure 8.1-1(a) The KIM-1 microcomputer



(b) The KIM-1 microcomputer connected to a 5 volt power supply  
(Note Teletype mode switch.)

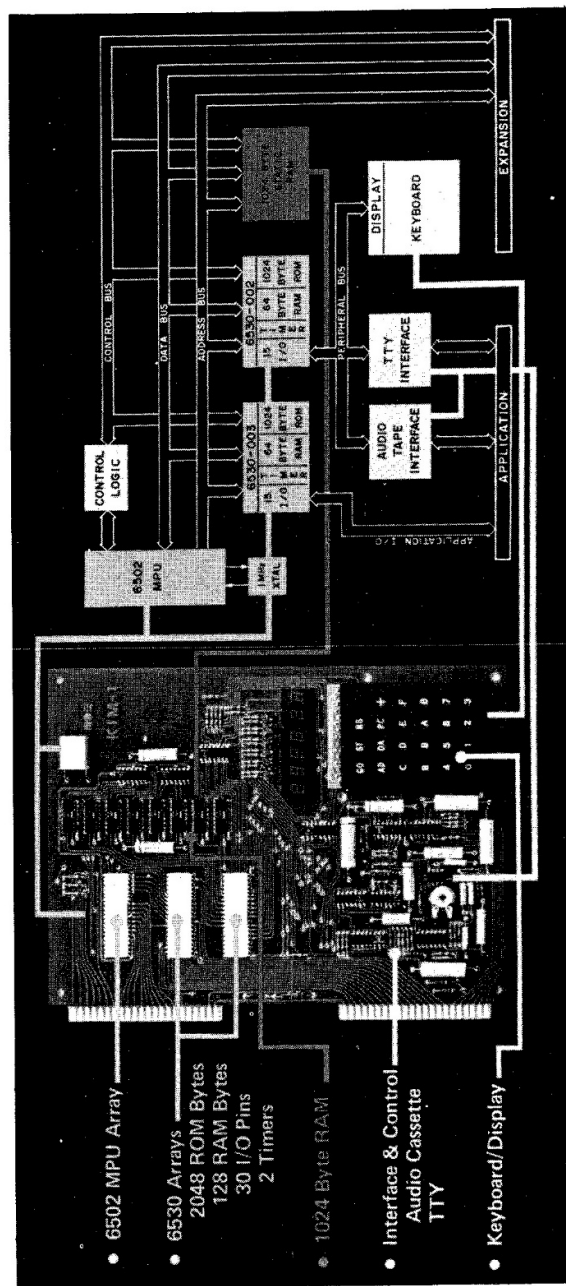


Figure 8.1-2 Identification of the major components on the KIM-1 microcomputer printed circuit board  
(Courtesy RCC News)

## 8.2 THE KIM-1 SYSTEM: (A Micro versus a Mini)

A microcomputer system, such as the KIM-1, differs from a minicomputer system in much more than just cost. With a minicomputer, such as the Digital Equipment Corporation PDP-11, the documentation provided by the manufacturer is extensive, diversified, and quite specialized. It is not unusual to find all of the following documents:

1. An Installer's Manual - provides unpacking information and power supply connections for the electrician.
2. An Operator's Manual - provides front panel information for the operator.
3. A Software Manual - provides information for the programmer.
4. A Hardware Manual - provides information about the theory of operation.
5. A Maintenance and Trouble Shooting Manual - provides the technician with schematic diagrams and periodic maintenance requirements, such as air filters which need to be changed or cleaned.

In the case of a microcomputer, the documentation is frequently very limited. (Fortunately, in the case of the KIM-1, rather extensive documentation is provided with three manuals; a Hardware Manual, a Software Manual, and a KIM-1 User's Manual.) Furthermore, in the case of the microcomputer, the installer, electrician, operator, programmer, and technician for maintenance and trouble shooting, are one and the same person.

## 8.3 AN EXAMPLE PROGRAM (EXAMPLE 8.1)

The purpose of this example program is to serve as a vehicle for illustrating the steps which are involved in going from the statement of a problem to a functioning computer program on the KIM-1 microcomputer. These steps will be the same, in principle, for any microprocessor; only the details will change from one particular microprocessor to another. These steps may be summarized as follows:

- \* statement of problem
- \* conversion to assembly language program
- \* conversion to machine language program
- \* entry of program into computer
- \* execution and debugging of program
- \* modification of program

The following additional steps are sometimes involved:

- \* entry and execution from a Teletype (includes use of paper tape)
- \* recording and playing back from an audio cassette tape recorder

**Example 8.1.** Add two positive numbers located initially in two memory locations and store the sum into a third memory location.

**Comments.** In order to keep this example as simple as possible:

1. We are considering only positive, one byte numbers.
2. We will not check for a carry out (overflow).

### 8.3.1 Example 8.1 – Assembly Language Program

The assembly language program for Example 8.1 is given in Table 8.3-1. The assembly language program is also frequently referred to as the “source code” and this designation is used in Table 8.3-1. The reader should note that the lefthand 4 columns are blank. At this time we could also insert the object code for all of the op codes, such as LDA, CLC, STA, etc. This has been done, and is shown in Table 8.3-2.

**TABLE 8.3-1 Source Code for Example 8.1.**

**Note that the object code columns are blank.**

OBJECT CODE				SOURCE CODE			
MEMORY	BYTE	BYTE	BYTE	LABEL	OP CODE	OPERAND	COMMENT
ADDRESS	1	2	3				
				NUM1			Designate one of the two numbers to be added as NUM 1.
				NUM2			Designate the second of the two numbers to be added as NUM2.
				SUM			Designate the sum of the two numbers to be SUM.
						⋮	
				SED			Set decimal Mode.
				LDA	NUM1		Load the accumulator with the numerical value of NUM1.
				CLC			Clear Carry.
				ADC	NUM2		Add with carry NUM1 + NUM2. (There is no ADD instruction.)
				STA	SUM		Store the sum of two numbers into memory location SUM.
				WAIT	JMP	WAIT	The 6502 microprocessor has no halt or stop command. This is a jump to itself.

**TABLE 8.3-2 Source Code but Only Partial Object Code for Example 8.1. The Memory Map is needed in order to complete this table.**

OBJECT CODE				SOURCE CODE		
MEMORY ADDRESS	BYTE 1	BYTE 2	BYTE 3	LABEL	OP CODE	OPERAND COMMENT
				NUM1		Designate one of the two numbers to be added as NUM1.
				NUM2		Designate the second of the two numbers to be added as NUM2.
				SUM		Designate the sum of the two numbers to be SUM.
					⋮	
F8				SED		Set Decimal Mode.
A5	00			LDA	NUM1	Load the accumulator with the numerical value of NUM1.
18				CLC		Clear carry.
65	01			ADC	NUM2	Add with carry NUM1 + NUM2. (There is no ADD instruction.)
85	02			STA	SUM	Store the sum of two numbers into memory location SUM.
4C				WAIT	JMP	WAIT The 6502 microprocessor has no halt or stop command. This is a jump to itself.

In order to complete Table 8.3-2, we need to decide where to store the program and where to store (or look for) data (the two numbers to be added and the sum). In other words, we need to know where RAM is located *and* if it is available to the user. Frequently, some bytes of RAM are reserved for use by the Monitor program in the microcomputer. As we shall see later, in the case of the KIM-1, 17 bytes (00EF to 00FF in hex) are reserved for the KIM-1 Monitor program. (The user may still decide to use these reserved memory locations, but with the caution that the Monitor program will store its own data in these memory locations.

We are now ready to proceed to the next section in order to learn more about the KIM-1 hardware. In particular, we are anxious to study the KIM-1 Memory Map, which will tell us how much memory we have, what type it is, and where it is located.

#### 8.4 KIM-1 MEMORY MAP AND TABLE

In order to be able to use any microcomputer (including the KIM-1) at the assembly language (or machine code) level, the user

must know the answers to the following questions:

- \* How much memory is there?
- \* What type is it; (ROM, RAM, PROM, etc.)
- \* Where is it in memory? (i.e., What is its address?)

The answers to these questions are provided in the memory map and the memory table. The memory map for the KIM-1 is shown in Figure 8.4-1. For convenience of notation, memory is divided into pages. Each page is  $256_{10}$  bytes long. For example, page 0 encompasses 0 through \$FF while page 4 encompasses \$400 through \$4FF. We will use decimal page numbers throughout this material. The reader will note that a user generated program may make use of the following areas of memory:

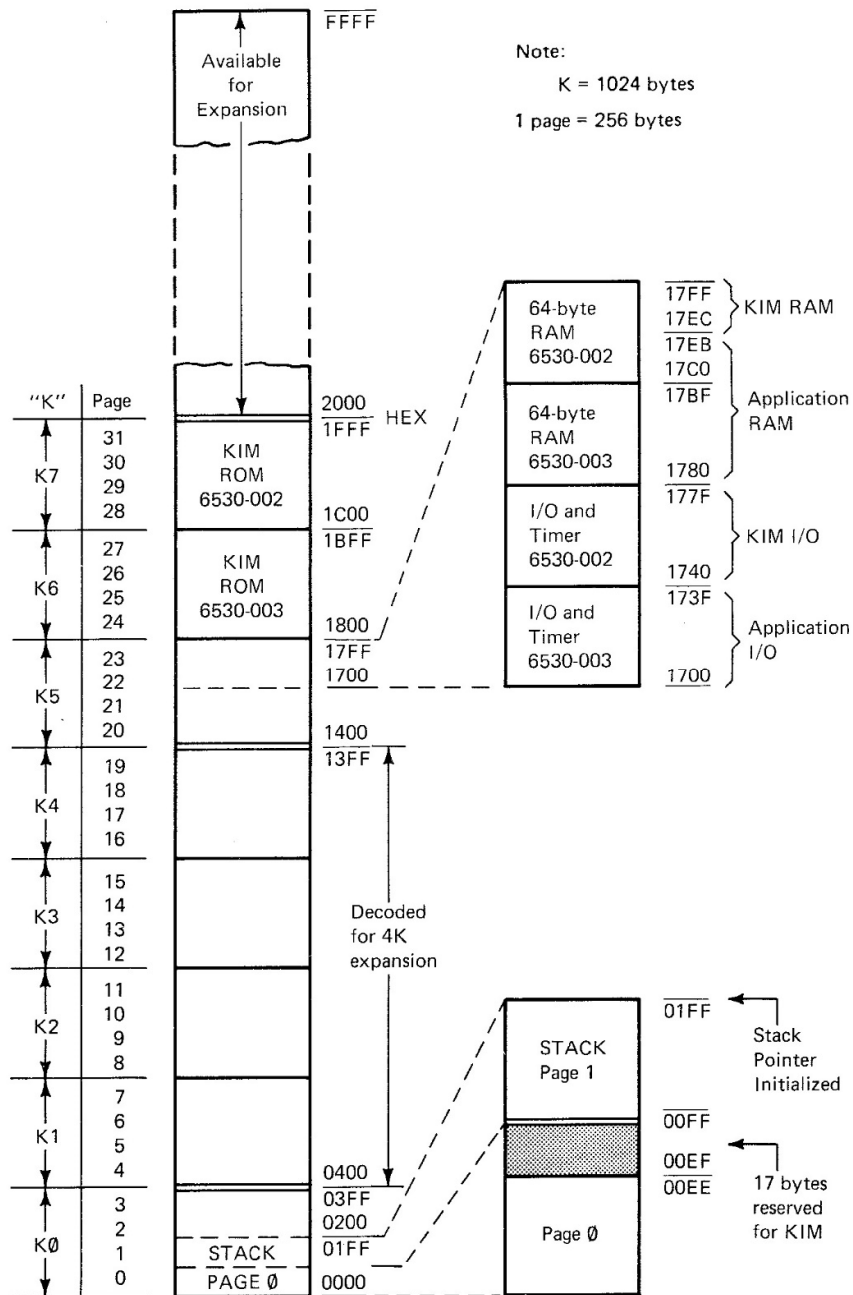
- \* all of page 0 except for 00EF to 00FF which are reserved for the KIM-1 Monitor program
- \* all of page 1 with the caution that the stack will use some of page 1
- \* all of page 2 and page 3
- \* in page 23 - all input/output locations from 1700 to 173F
  - all 64 bytes of RAM from 1780 to 17BF
  - an additional 44 bytes of RAM from 17C0 to 17EB

The memory table for the KIM-1 is given in Table 8.4-1. The reader will note that some of the information is the same as that provided in the memory map of Figure 8.4-1, but additional and more detailed address information is also provided. When using the KIM-1 the user will need to refer frequently to both of these memory information sources.

We will now use the information provided by the memory map and the memory table to complete the coding for Example 8.1 so that we may enter the program into the KIM-1 memory. This will be the subject for the next section.

## 8.5 MACHINE CODE FOR EXAMPLE 8.1

If we now refer back to Table 8.3-2, the reader will note that in order to complete that table, we need to decide where to store the program and where to store the data that will be used by the program in the RAM memory. It is usually a good programming technique to store the numerical data in page 0 because then we can use the so-called zero page instructions which require only two, rather than three, bytes



**Figure 8.4-1** Memory map for the KIM-1 microcomputer (MOS Technology)

**TABLE 8.4-1 Memory Table for the KIM-1 Microcomputer**  
(Courtesy MOS Technology, Inc.)

ADDRESS	AREA	LABEL	FUNCTION
00EF	Machine Register Storage Buffer	PCL	Program Counter - Low Order Byte
00F0		PCH	Program Counter - High Order Byte
00F1		P	Status Register
00F2		SP	Stack Pointer
00F3		A	Accumulator
00F4		Y	Y-Index Register
00F5		X	X-Index Register
1700	Application I/O	PAD	6530-003 A Data Register
1701		PADD	6530-003 A Data Direction Register
1702		PBD	6530-003 B Data Register
1703		PBDD	6530-003 B Data Direction Register
1704	Interval Timer		6530-003 Interval Timer
170F			(see Table 3.7.1-1, page 144)
17F5	Audio Tape Load & Dump	SAL	Starting Address - Low Order Byte
17F6		SAH	Starting Address - High Order Byte
17F7		EAL	Ending Address - Low Order Byte
17F8		EAH	Ending Address - High Order Byte
17F9		ID	File Identification Number
17FA	Interrupt Vectors	NMIL	NMI Vector - Low Order Byte
17FB		NMIH	NMI Vector - High Order Byte
17FC		RSTL	RST Vector - Low Order Byte
17FD		RSTH	RST Vector - High Order Byte
17FE		IRQL	IRQ Vector - Low Order Byte
17FF		IRQH	IRQ Vector - High Order Byte
1800	Audio Tape	DUMPT	Start Address - Audio Tape Dump
1873		LOADT	Start Address - Audio Tape Load
1C00	STOP Key + SST		Start Address for NMI using KIM "Save Machine" Routine (Load in 17FA & 17FB)
17F7	Paper Tape Dump (Q)	EAL	Ending Address - Low Order Byte
17F8		EAH	Ending Address - High Order Byte



of code, and hence less memory. Also, these two byte, zero page, instructions require less time to execute. We will not use this programming technique here because we require only three bytes of memory for the numerical data and the program itself is so short.

Using the memory map data given in Figure 8.4-1, we decide to:

- \* store NUM1, NUM2, and the SUM into memory locations 0000, 0001, and 0002 (base 16), respectively
- \* store the machine code into memory locations beginning with memory location 0010 (base 16)

As a result of making these memory allocation decisions, the machine code for Example 8.1 may now be completed and it is shown in Table 8.5-1.

We are now ready to load the machine code for Example 8.1 into the KIM-1 memory. This is the subject of the next section.

**TABLE 8.5-1 Assembly and Machine Code for Example 8.1**

OBJECT CODE				SOURCE CODE			
MEMORY ADDRESS	BYTE 1	BYTE 2	BYTE 3	LABEL	OP CODE	OPERAND	COMMENT
0000				NUM1			Designate one of the two numbers to be added as NUM1.
0001				NUM2			Designate the second of the two numbers to be added as NUM2.
0002				SUM			Designate the sum of the two numbers to be SUM.
					⋮		
0010	F8			SED			Set Decimal Mode.
0011	A5	00		LDA		NUM1	Load the accumulator with the numerical value of NUM1.
0013	18			CLC			Clear Carry.
0014	65	01		ADC		NUM2	Add with carry NUM1+NUM2. (There is no ADD instruction.)
0016	85	02		STA		SUM	Store the sum of two numbers into memory location SUM.
0018	4C	18	00	WAIT	JMP	WAIT	The 6502 microprocessor has no halt or stop command. This is a jump to itself.

Comment 1: We have arbitrarily decided to store NUM1, NUM2, and SUM into memory locations 0000, 0001, and 0002, respectively.

Comment 2: We have also arbitrarily decided to begin storing the object code into memory locations beginning with memory location 0010.

## 8.6 ENTERING EXAMPLE 8.1 CODE INTO KIM-1

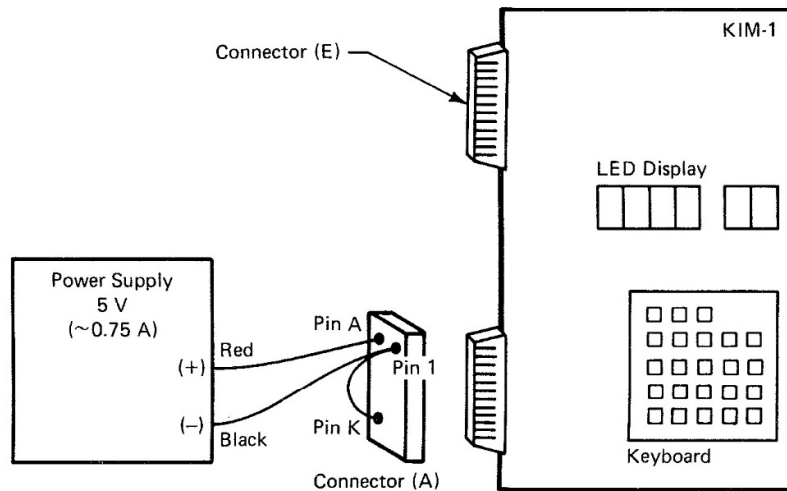
We are now ready to enter the machine code for Example 8.1 (as given in Table 8.5-1) into the KIM-1 memory locations we selected in Section 8.4.

The first step is to apply power to the KIM-1 microcomputer printed circuit board. In this case, we need 5 volts,  $\pm 5\%$ , with a current rating of at least 0.75 amperes. The circuit connections are shown in Figure 8.6-1. It is strongly recommended that these connections be checked very carefully each time the KIM-1 is used. Power supply polarity reversals can be catastrophic.

Table 8.6-1 presents a detailed key by key stroke entry of the machine code for Example 8.1 into the KIM-1 memory according to the decisions made and presented in Table 8.5-1. Table 8.6-1 also gives the information that the user will see on the display following the depression of the various keys. Comments are supplied, where necessary.

**TABLE 8.6-1 Key by Keystroke Entry of Example 8.1 Code per TABLE 8.5-1 into KIM-1 Memory**

Press Keys	See on Display	Comments
	blank	Apply 5 volts, $\pm 5\%$ , 0.75 amp. per Figure 8.6-1. There is no power ON-OFF switch on the KIM-1. The SST-ON switch is for selecting the single-step mode.
<b>RS</b>	xxxx xx	x implies an unpredictable display.
<b>AD</b>	no change	Puts KIM into address mode.
<b>0 0 1 0</b>	0010 xx	Contents of memory location 0010 are displayed.
<b>DA</b>	no change	Puts KIM into DATA mode.
<b>F 8</b>	0010 F8	Enters F8 code into memory location 0010.
<b>+</b>	0011 xx	Increments the memory address displayed by one. This is easier than pressing the four keys for location 0011, which would accomplish the same thing.
<b>A 5</b>	0011 A5	Enters A5 code into memory location 0011.
<b>+ 0 0</b>	0012 00	Enters 00 code into memory location 0012.
<b>+ 1 8</b>	0013 18	Similarly, enters 18 code into memory location 0013, etc.
<b>+ 6 5</b>	0014 65	
<b>+ 0 1</b>	0015 01	
<b>+ 8 5</b>	0016 85	
<b>+ 0 2</b>	0017 02	This completes the loading of the code for Example 8.1. It is suggested that the user go back to memory location 0010, and then, using the <b>+</b> key, check sequentially through memory location 001A for possible errors.
<b>+ 4 C</b>	0018 4C	
<b>+ 1 8</b>	0019 18	
<b>+ 0 0</b>	001A 00	



**Figure 8.6-1** When used without Teletype, CRT, or audio tape recorder, the KIM-1 requires only one power supply connected as shown

## 8.7 EXECUTION OF EXAMPLE 8.1 FROM THE KIM-1 KEYBOARD

Table 8.7-1 presents a key by key stroke sequence for entering the numerical values of NUM1 and NUM2 and then executing the program using the **GO** key. We have arbitrarily selected NUM1 = 43 and NUM2 = 25. Again, we have presented the displays which the user will see after depressing the various keys and appended comments, where necessary. The resulting SUM is displayed in the last step of the table and is seen to be equal to 68, which is the correct answer in the decimal mode.

It is suggested that the user modify the program by changing the SED (set decimal mode) op code to CLD (clear decimal mode) and try different numerical values for NUM1 and NUM2 in order to gain a greater understanding of what the decimal and binary modes in the KIM-1 really mean. Problems are provided at the end of this chapter to help guide the user in this endeavor.

### 8.7.1 Single Step Execution of Example 8.1 Program

Table 8.7-2 presents a key by key stroke sequence for entering the numerical values for NUM1 and NUM2 and then executing the program

**TABLE 8.7-1 Execution of the Example 8.1 Program from the KIM-1 Keyboard**  
**NUM1 + NUM2 = SUM where NUM1 = 43; NUM2 = 25**

Press Keys	See on Display	Comments
[RS]	xxxx xx	Reset. x implies display is not predictable.
[AD]	no change	Puts KIM-1 into ADDRESS input mode.
[0][0][0][0]	0000 xx	Present contents of memory location 0000 are displayed.
[DA]	no change	Puts KIM-1 into DATA input mode.
[4][3]	0000 43	Enter 43 as the numerical value of NUM1 into location 0000.
[+][2][5]	0001 25	Enter 25 as the numerical value of NUM2 into location 0001.
[+][0][0]	0002 00	Clears memory location 0002 where the SUM will be placed.
[AD][0][0][1][0]	0010 F8	Prepares to run program by loading the beginning address (0010) of our program.
[GO]	blank	Program executes in literally a few microseconds.
[RS]	0010 F8	Reset KIM-1.
[AD][0][0][0][0]	0000 43	NUM1 is still equal to 43.
[+]	0001 25	NUM2 is still equal to 25.
[+]	0002 68	SUM is 68 and is the correct answer in the decimal mode. Recall that we had previously cleared this memory location.

**TABLE 8.7-2 Single Step Execution of Example 8.1 Program from the KIM-1 Keyboard**  
**NUM1 + NUM2 = SUM where NUM1 = 43; NUM2 = 25**

Press Keys	See on Display	Comments
[RS]	xxxx xx	Reset.
[AD][1][7][F][A]	17FA xx	This sequence of 3 lines must be performed in order to use the [ST] key and/or the single step mode. To be more specific: memory location 1C00 is stored into interrupt vector locations 17FA and 17FB.
[DA][0][0]	17FA 00	
[+][1][C]	17FB 1C	
[AD][0][0][0][0]	0000 xx	Present contents of memory location 0000 are displayed.
[DA][4][3]	0000 43	Enter 43 for NUM1 into location 0001.
[+][2][5]	0001 25	Enter 25 for NUM2 into location 0002.
[+][0][0]	0002 00	Clear location 0002 where SUM will be placed.
[AD][0][0][1][0]	0010 F8	Loads beginning address (0010) of our program.
[GO]	0011 A5	Each time the [GO] key is depressed, a single line of code is executed and the memory address of the next line of code is displayed. A "line of code" is defined in TABLE 8.5-1. Note: The same numbers are missing here as in column 1 of TABLE 8.5-1.
[GO]	0013 18	
[GO]	0014 65	
[GO]	0016 85	
[GO]	0018 4C	
[GO]	no change	
[AD][0][0][0][0]	0000 43	NUM1 is still equal to 43.
[+]	0001 25	NUM2 is still equal to 25.
[+]	0002 68	SUM is 68 again. Recall that we had previously cleared this memory location to zero. Do not forget that the SST-ON switch is in the ON position.

for Example 8.1 in the *single step mode*. We have again selected NUM1 = 43 and NUM2 = 25. The usual display information and comments are given.

An important item for the reader to note in Table 8.7-2 is the need to properly initialize the interrupt vector locations 17FA and 17FB. This must always be done if the **ST** key and the single step switch mode are to function properly. For more information, please refer to the **NMI** section of the 6502 microprocessor discussion.

The KIM-1 user will find that the single step mode is used primarily during program debugging.

## 8.8 DECIMAL OR BINARY CODE

Any system which uses the 6502 microprocessor has the option of operating in either the decimal mode or the binary mode. In the case of the KIM-1, the designers choose not to include a **CLD** or **SED** command in the reset programs, so it is not possible to predict whether the KIM-1 will be in the decimal mode or in the binary mode when first energized. (These writers' experiences with about a dozen KIM-1s conform with this observation; about half come up on the binary mode and about half in the decimal mode; unfortunately, it is not always the same half.)

This uncertainty - decimal or binary mode - is responsible for most of the problems encountered by first-time KIM-1 users. These problems appear to the user to be one of the following two major types:

1. The numerical results to arithmetic operations do not appear to be correct.
2. Programs recorded on audio cassette tapes do not load.

Problems of the first type are solved if the user (programmer) develops the habit of preceding each arithmetic operation with a **CLD** or **SED** command. (The reader will recall that our solution to Example 8.1 in Table 8.2-1 followed this suggestion.)

Problems of the second type are solved if the user (operator) sets the KIM-1 into the binary or decimal mode immediately after powerup. This may be done by loading all 0s or all 1s into the memory location 00F1. When we press the Go button, the contents of this location will be loaded into the P register. Actually only the decimal bit in the status register needs to be set to a 1 for the decimal mode and set to 0 for the binary mode, but it is much easier for the user to set all 1s or all 0s. (There are 64 different ways to set one bit to a 1 or to a 0 in a byte.)

## 8.9 KIM-1 KEYBOARD KEY FUNCTIONS

The KIM-1 keyboard key functions are summarized in Figure 8.9-1.

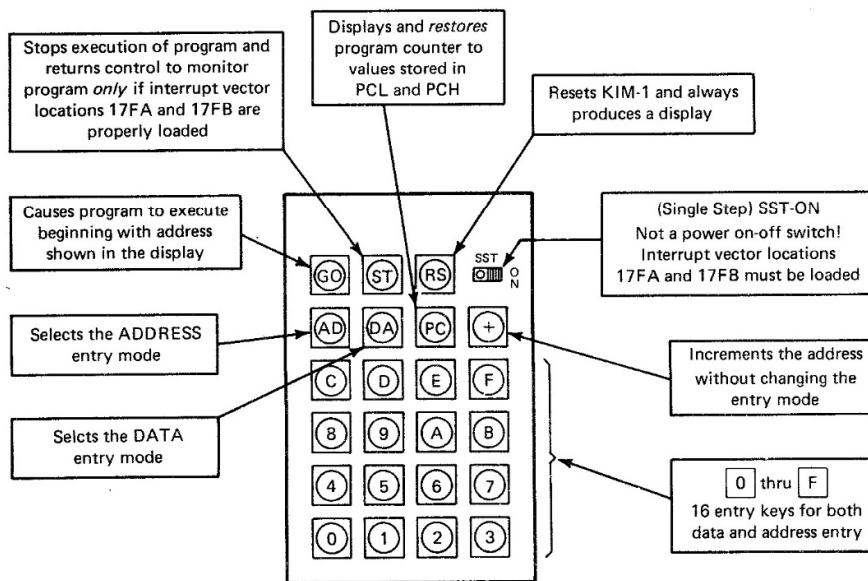


Figure 8.9-1 Summary of KIM-1 keyboard key functions and cautions

## 8.10 OPERATING THE KIM-1 VIA A TELETYPE

Before energizing a KIM-1 system that is using a serial teleprinter such as the TELETYPE<sup>®</sup> Model 33ASR (Automatic Send and Recive)\* the following items should be verified.

1. The Model 33 is wired for full duplex 20 mA operation.
2. The 5 volt power supply is connected to the KIM-1 application connector as shown in Figure 8.6-1. Be sure a jumper is connected between pins 1 and K.

\* Models 33KSR or RO can also be used.

3. Pins R, U, S, and T of the KIM application connector are connected to the Teletype as follows:

R	X-7 or P2-8
U	X-6 or P2-7
S	X-3 or P2-5
T	X-4 or P2-6

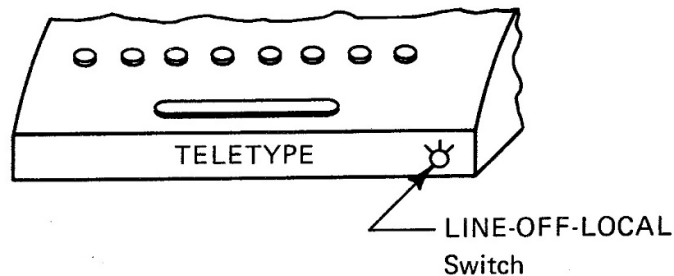
X denotes Terminal Block X and P2 denotes Plug 2. (On the 33ASR Teletype terminal block X can be located by noting the ac line cord connected to X1 and X2. It is the second receptacle from the left on the top row as viewed from the rear of the machine.

4. The toggle switch TTY-KB (keyboard) is in the TTY (TELETYPE) position.

### 8.10.1 The KIM-1 Prompting Message

If you have checked all of the electrical connections specified in Sec. 8.1, you may now proceed to do the following:

1. Energize the 5 volt power supply.
2. Plug the TELETYPE into a 115 volts, 60 Hz. power outlet and turn the switch marked LINE-OFF-LOCAL to LINE. This switch is located at the right side of the small vertical front panel strip which bears the letters T E L E T Y P E as shown in Figure 8.10.1.





**Figure 8.10-1** Closeup view of LINE-OFF-LOCAL switch on model 33ASR TELETYPE

The following steps are most important because the KIM-1 system adjusts to the bit rate of the serial teleprinter and requires this sequence of key strokes to establish the proper bit rate:

Press KIM key	Press TTY key	See printed
		KIM xxxx xx

where xxxx xx denotes an arbitrary, unpredictable set of six (6) printed characters. If everything is working properly, you should immediately observe the message KIM xxxx xx being typed in the two line format shown above. This is a prompting message telling you that the TTY is on-line and that the KIM-1 system is ready to accept commands from the TTY keyboard.

If the "KIM" prompting message is not typed, repeat the two key sequence given above; that is, press the  key on the KIM-1 keyboard and the  key on the TTY. If the "KIM" message is still not typed, recheck all connections and try again. If the problem still persists, obtain assistance.

### 8.10.2 Possible Input and Output Modes

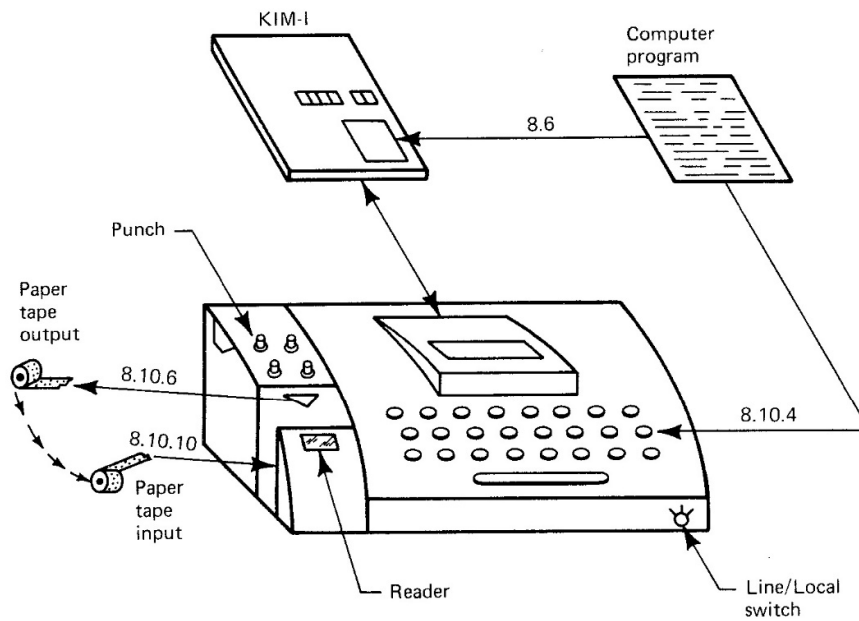
Immediately after the "KIM" prompting message is typed on the TTY, the KIM-1 system is ready to accept commands from the TTY keyboard. As soon as you decide what type of input and/or output you desire, you may refer to Figure 8.10.2 and determine what sections describe the procedure you must follow.

### 8.10.3 Another Simple Example (Example 8.2)

The purpose of this simple example is to explain and to illustrate the different possible input and output modes using the TELETYPE Model 33ASR. It also illustrates another way of terminating a program. Rather than just jumping to itself, the program jumps into the Monitor Program and displays the numerical value of the SUM and its memory location.

**Example 8.2.** Add two positive integers located in two memory locations and store the sum into a third memory location. Output and display the sum and its memory location. It is not necessary to check for overflow. The solution to Example 8.2 is presented in Table 8.10-1.





**Figure 8.10-2** Pictorial view of possible input/output modes using a TELETYPE 33ASR with the KIM-1

**TABLE 8.10-1** Source code and machine code for adding two numbers, storing, and displaying the sum. Refer to Example 8.2 for more details.

Memory Map	Byte 1	Byte 2	Byte 3	SOURCE CODE		
				LABEL	OP CODE	OPERAND
00				NUM1		
01				NUM2		
02				SUM		
03	18			PROGRAM	CLC	
04	F8				SED	
05	A5	00			LDA	NUM1
07	65	01			ADC	NUM2
09	85	02			STA	SUM
0B	A9	02			LDA	#02
0D	85	FA			STA	\$FA
0F	A9	00			LDA	#\$00
11	85	FB			STA	\$FB
13	4C	4F	1C		JMP	1C4F

### 8.10.4 Using the TELETYPE Keyboard to Enter a Program

While using the TELETYPE keyboard to enter a program, you may perform the following operations:

#### To Select an Address

1. Type four hex keys (0 to F) to define the address.
2. Press the **SPACE** bar.

The TTY will respond by typing the address code selected followed by a two hex digit code corresponding to the data stored at the selected address location.

*Example:*

Type: 1802 **SPACE**  
TTY printer responds: 1802 8D

showing that data 8D is stored at memory location 1802, which is in ROM, so its contents cannot be changed.

*Time saving hint:* Leading zeros need not be entered.

*Examples:*

EF **SPACE** selects address 00EF  
A **SPACE** selects address 000A  
**SPACE** selects address 0000

#### To Modify an Address

1. Select an address as above. Don't forget to press the **SPACE** bar.
2. Type two hex characters to define the data to be stored at this address.
3. Press the TTY period key, hereafter denoted by  $\odot$ .


*Example:*


Type: 0234 **SPACE**  
TTY printer responds: 0234 xx  
Type: 6D  $\odot$   
TTY printer responds: 0235 xx

*Note:* The selected address 0234 has been incremented automatically by one to the next address 0235 and the contents xx of memory location 0235 have been displayed.

*Time saving hint:* Leading zeros need not be entered.


*Examples:*

A  enters data 0A

 enters data 00

(This is a frequent source of errors for even experienced programmers who clear memory locations inadvertently by using this short-cut.)

#### To Step To Next Address (Without modifying the contents of the current address)

1. Press key 

*Example:*

See printed: 1234 xx

Type:



Printer responds: 1235 xx


Type:



Printer responds: 1236 xx

etc.

#### To Step to Preceding Address (Without modifying the contents of the current address)

1. Press key 

*Example:*

See printed: 1234 xx

Type:



Printer responds: 1233 xx


Type:





Printer responds: 1232 xx


etc.



**To Abort Current Operation**

1. Press key 




*Example:*

See printed:	1264	
Type:		
Printer responds:	KIM	
	xxxx xx	
Type:	1234	
Printer responds:	1234 xx	



In this example, the  key was used to correct an erroneous address selection.



*Note:* The  key must be depressed after each depression of the KIM-1  key in order to allow the operating program to define the serial bit rate for the particular Teletype being used.

**8.10.5 Using the TELETYPE Keyboard to Execute a Program****To Execute a Program (From the TTY keyboard)**

1. Enter the starting address of the program followed by .
2. Type  .[This is analogous to the  key on the KIM-1 keyboard.]

*Example 8.2* may be executed as follows:

Type:	0003 
See printed:	0003 18
Type:	 (on TELETYPE keyboard!)

*Comments:* In this example, program execution begins from location 0003 and will continue until memory location 0015 is reached. In other cases, program execution continues until the  or  keys on the KIM-1 keyboard are depressed.

Actual TTY output for Sec. 8.10.3 Example 8.2 stored in the KIM-1 with 8 stored in memory location 0000 and 9 stored in memory location 0001. As you may recall, the sum = 17 and is stored in memory location 0002. The Teletype output is shown in Figure 8.10.3

			COMMENTS
		0003	Program begins in memory location 0003.
0003	18	G	Note that G is typed on the same line as the previous response.
KIM			
0002	17		Sum = 17 and is stored in memory location 0002.

**Figure 8.10-3** Actual TELETYPE output for Example 8.2  
 NUM1 + NUM2 = SUM for NUM1 = 8 and NUM2 = 9  
 SUM = 17

#### To Execute a Program (Single step from TTY keyboard)

1. Enter the starting address of the program followed by **SPACE**.
2. Slide switch on KIM-1 keyboard marked SST-ON to ON.
3. Type key **G** once for each step of program you wish to have executed.

*Example 8.2* may be executed from the TELETYPE keyboard in the *single-step* mode as follows:

Type:	0003 <b>SPACE</b>
See printed:	0003 18
Type:	<b>G</b> and continue depressing the <b>G</b> key until output does not change. See Figure 8.10.4 for the actual output as well as comments.

### 8.10.6 Using the TELETYPE Keyboard to Punch a Paper Tape

#### To Punch Paper Tape

1. Load and thread blank paper tape into the punch unit.
2. Decide on the starting address and ending address of the data block to be punched on the paper tape.

A paper tape for *Example 8.2* may be punched as follows:



Type: 17F7 SPACE  
 See printed: 17F7 xx  
 Type: 15 Q  
 See printed: 17F8 xx  
 Type: 00 Q  
 See printed: 17F9 xx      Comment: Ignore this memory location  
                                                                                          because our next step is to  
                                                                                          specify the starting memory  
                                                                                          location.

Type: 0000 SPACE  
 See printed: 0000 xx      Activate punch by depressing but-  
                                                                                          ton on punch marked ON.

Press key: Q      Comment: Watch and listen to the tape  
                                                                                          being punched!

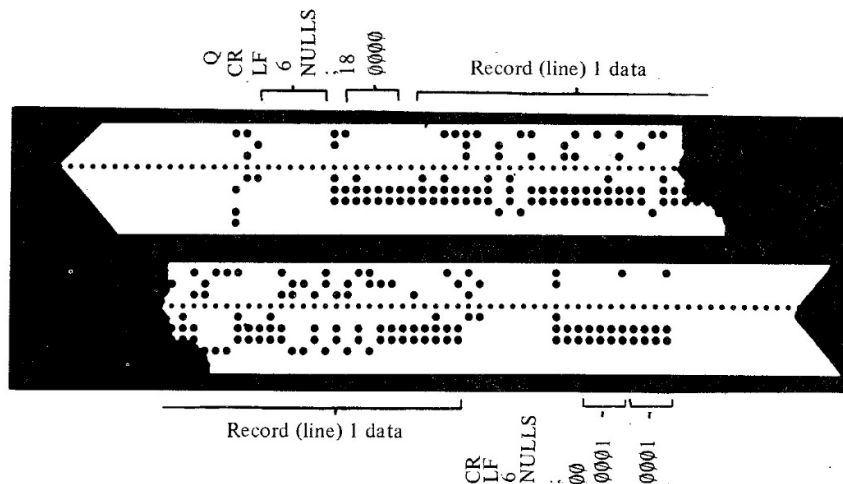
Please refer to Figure 8.10.5 for the actual printed output corresponding to Example 8.2. Figure 8.10.6 presents the actual punched paper tape for Example 8.2.

Q

```

;180000000091718F8A50065018502A90285FAA90085FB4C4F1C22040812
;0000010001
  
```

**Figure 8.10-5** Actual printed output for Example 8.2

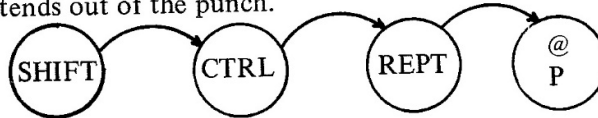


**Figure 8.10-6** Actual punched paper tape for Example 8.2

*A Handy Hint:*

It is convenient to have a leader on each paper tape which has the minimum number of holes punched so that you can write your name, date, and other information. You can do this in three steps:

1. Press keys *in this order and hold down* until about 2 inches of tape extends out of the punch.



2. Tear off and discard this paper tape. The “V” cutting bar will form the tip of an arrow so that you will know in what direction the tape is to move. The sprocket holes, which are off-center, will prevent you from placing the tape in up-side-down.
3. Repeat Step 1 until you have the desired amount of leader at the front of your tape.

### 8.10.7 Using the TELETYPE Keyboard to List a Program

A printed record of the contents of the KIM-1 memory may be typed on the TELETYPE. The procedure is the same as for punching a paper tape *except* that the punch mechanism is not activated. This is assured if you press the OFF button located on top of the punch mechanism.

### 8.10.8 Using the TELETYPE Keyboard to List KIM-1 Memory

There is no difference between using the TELETYPE keyboard to list the contents of the KIM-1 memory and to list a program as described in Sec. 8.10.7.

### 8.10.9 KIM-1 Paper Tape Format

The paper tape DUMP and LOAD routines in the KIM-1 Monitor store and retrieve data in a format designed to reduce errors. Each byte of data is converted to two half bytes (nibbles). Each half byte (in hex) is translated to its ASCII equivalent and written out onto paper tape in ASCII form.

Each line of the output begins with a “,” character (ASCII 3B) to indicate the start of a valid record. The next byte (18 in hex) is the number of data bytes in that line. Then the starting address lower order

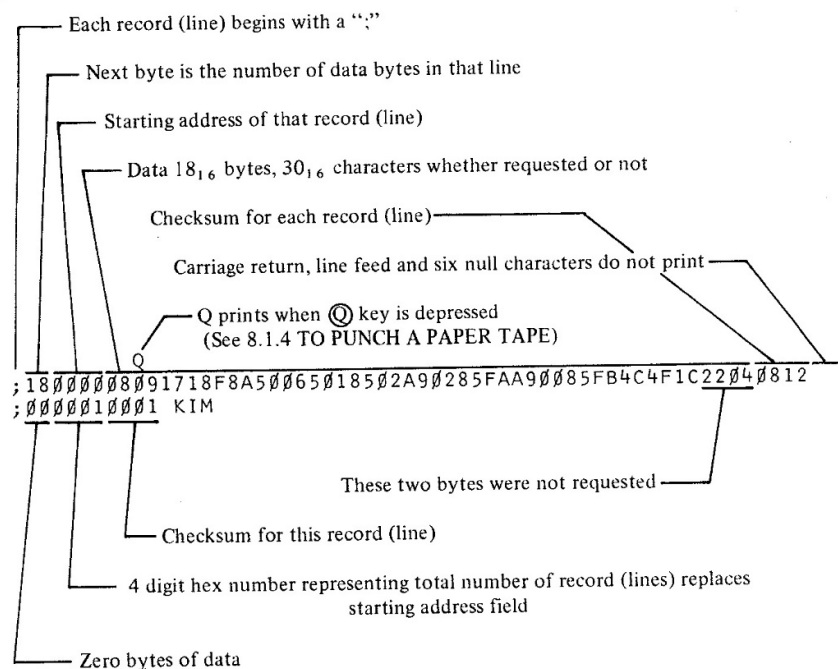


byte (1 byte, 2 characters), starting address higher order byte (1 byte, 2 characters), and then then data ( $18_{16}$  bytes,  $30_{16}$  characters) follow. Each record (line) is terminated by its checksum (2 bytes, 4 characters), a carriage return (ASCII 0D), line feed (ASCII 0A), and six “Null” characters (ASCII 00).

The last record (line) has zero bytes of data (indicated by ;00). The starting address field is replaced by a four digit hex number representing the total number of data records (lines), followed by the checksum. A “XOFF” character ends the transmission.

During a load, all incoming data is ignored until a “;” character is encountered. The receipt of nonASCII data or a mismatch between the calculated checksum and the checksum read from the tape will cause an error condition and the KIM-1 Monitor Program will cause the message “KIM ERROR” to be printed. The checksum is calculated by adding all data in each record (line) except the “;” character.

Figure 8.10.7 illustrates and interprets how the above description of the paper tape format applies to the Example 8.2 TELETYPE printed output.



**Figure 8.10-7** Illustration and interpretation of the KIM-1 paper tape format as it applies to Example 8.2 output

### 8.10.10 Reading a Paper Tape into the KIM-1

Paper tapes read into the KIM-1 system must have the format specified in Section 8.10.9. Paper tapes generated by the KIM-1 have, of course, this format. To read a paper tape with the proper format, proceed as follows:

1. Set the Paper Tape Reader switch to FREE.
2. Insert the tape into the tape reader mechanism being careful to align the sprocket holes. Move the tape slightly back and forth to make sure it is free.
3. Press **[RS]** on the KIM-1 keyboard and **(RUB OUT)** key on TELETYPE.
4. Type **(L)** on TELETYPE (mnemonic for Load).
5. Move the tape reader switch to START and remove your hand. Switch will automatically return to its center position, and the reader will begin to read the tape.

Caution: Keep fingers away from the paper tape reading mechanism!

Figure 8.10.8 illustrates the Teletype output corresponding to Example 8.2.

```

      LQ
;18000008091718F8A50065018502A90285FAA90085FB4C4F1C22040812
;0000010001 KIM

KIM
0001 09

```

**Figure 8.10-8** Actual TELETYPE printed output when paper tape for Example 8.2 is loaded into the KIM-1

## 8.11 ADDING AN AUDIO TAPE RECORDER TO THE KIM-1

### 8.11.1 Why Add an Audio Tape Recorder?

There are at least four reasons to add an audio tape recorder to the KIM-1 microcomputer:

1. Programs and data which the user has stored in the KIM-1 RAM may be saved before the power is turned off.
2. An audio tape recorder is much less expensive (as low as \$20) compared to a paper tape punch and reader.

3. Computer programs and data which are stored on magnetic tape may be readily duplicated and exchanged with other KIM-1 users.
4. Magnetic audio tape, as compared to paper tape, permits a storage density, a smaller more convenient package to handle, and is reusable.

Although both reel-to-reel and cassette tape recorders may be used, the cassette type is used in almost all KIM-1 applications because of its small size, portability, and cost. The discussion that follows will apply to both reel-to-reel and cassette types. (Other types of tape recorders, such as miniaturized recorders without external jacks, are usually intended primarily for dictating applications and require internal modifications. Recorders requiring internal modifications will not be discussed here.)

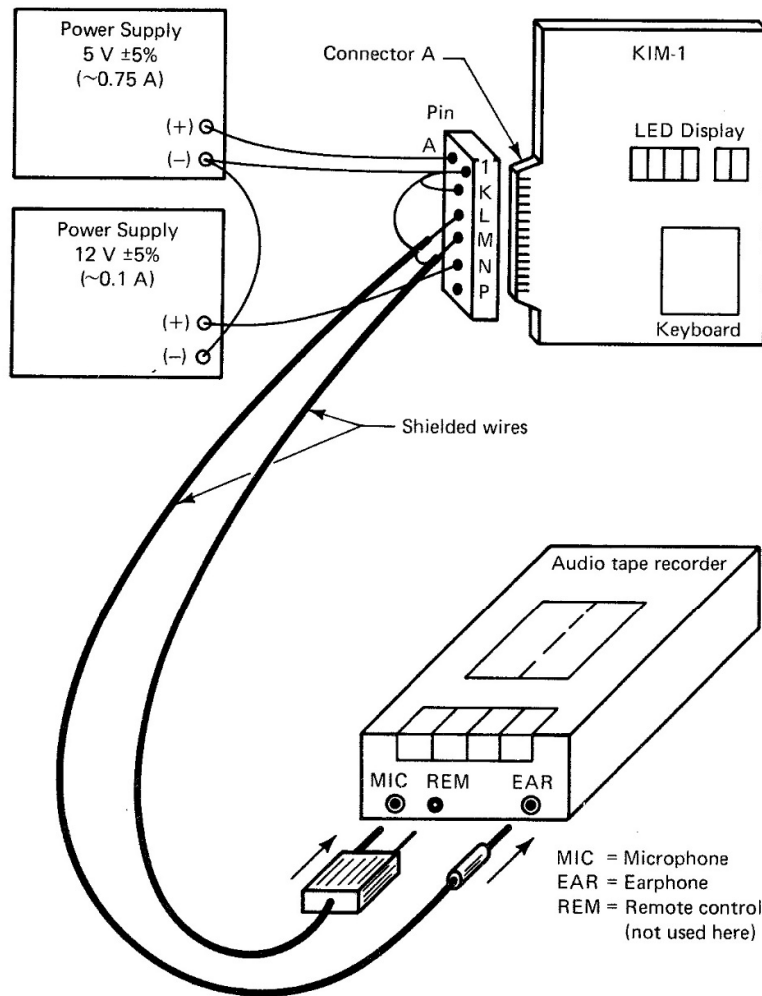
### 8.11.2 The Audio Tape Recorder Connection

The audio tape recorder unit to be used should possess the following features:

1. A jack, usually labeled microphone or input, to permit recording the electrical signal provided by the KIM-1 circuit board onto the magnetic tape.
2. A jack, usually labeled earphone or external speaker, to supply the electrical signal of the appropriate level to the KIM-1 circuit board for loading into the microcomputer memory.
3. The standard control switches for PLAY, RECORD, REWIND, and STOP. Other features, such as FAST FORWARD or FAST REVERSE, are convenient for locating programs when more than one program is stored on a tape side.

With the power off to all of the components, check that all of the cables and jumpers are present as shown in Figure 8.11.1. All wires, except for the shielded wires, should be as short as possible and kept away from other wires which might introduce electrical noise.

The Audio Data Out (LO) at pin M has a level of approximately 15 mv peak. The Audio Data Out (HI) at pin P has a level of approximately 1 volt peak. Pin P is used with some of the more expensive tape recorders which have an input (usually labeled "LINE") that accepts higher voltage input signals. The lower level Audio Data Output at pin M is the one which is used with most small inexpensive audio cassette tape recorders.



**Figure 8.11-1** With an audio tape recorder, the KIM-1 system requires one 5 volt power supply, one 12 volt power supply, one jumper, and two shielded wires connected as shown above

### 8.11.3 How to Record On an Audio Tape

Before we list the steps to be followed for recording a program and/or data from memory on an audio magnetic tape, let us document the assumptions we are making:

- a) There is a program and/or data stored in memory with a known starting and a known ending address.
- b) The program and data have been thoroughly tested and the program produces known results given specific data. This prevents confusing program errors with recording/playback errors.
- c) The special addresses for the Audio Tape Load and Audio Tape Dump subroutines are known or obtained from Figure 8.11.2.

The procedure for recording on an audio magnetic tape from the KIM-1 memory is:

(If not already on, turn on the 5 volt and the 12 volt power supplies and check power to the tape recorder, if it is not battery powered)\*\*

Step 1.\* Verify that the 6502 microprocessor is in the binary mode. This may be done by inspecting the contents of memory location \$00F1 and recalling that the decimal mode bit (1 = True) is bit three (b3 of b0 - b7) of the processor status word (which is stored in memory location \$00F1). Since bit 3 is zero for so many different hexadecimal numbers, it is frequently easier just to *clear* the decimal mode and place into the binary mode by performing the following steps:

Press Keys	Display	Comments
<b>AD</b> <b>0</b> <b>0</b> <b>F</b> <b>1</b>	00F1 xx	Selects the address mode and displays the contents of memory location \$00F1.
<b>DA</b> <b>0</b> <b>0</b>	00F1 00	Stores 00 into memory location \$00F1.

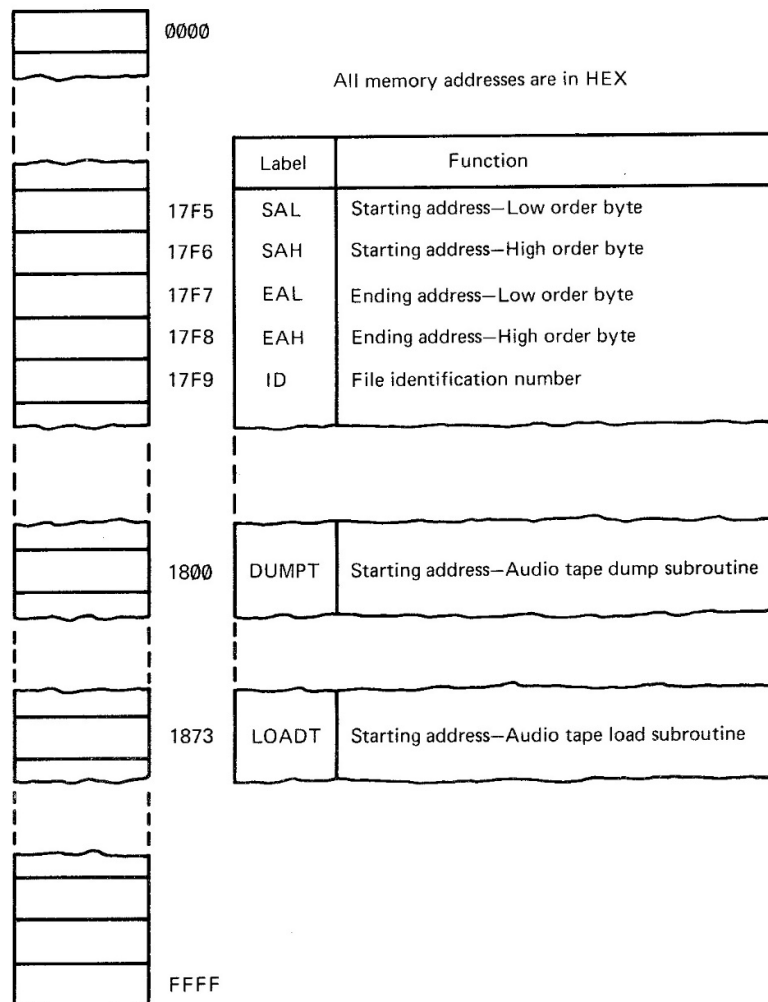
Step 2. Store the starting address in memory locations 17F5 and 17F6. Be sure to place the low order byte into 17F5 and the high order byte into 17F6.

Step 3. Store the ending address into memory locations 17F7 and 17F8. Be sure to place the low order byte into 17F7 and the high order byte into 17F8.

*The ending address is defined to be one greater than the last actual memory location of our program.*

\* If it were financially acceptable, we could add CLD instructions to the Audio Tape subroutines, manufacture new ROMs, and remove the need for Step 1.

\*\* Some users have reported erratic operation when using the internal batteries in some of the tape recorders and they suggest using 115 volt 60 Hz adaptors instead.



**Figure 8.11-2** Special memory addresses for Audio Tape Load (1873) and Audio Tape Dump (1800), starting address, ending address, and ID

Step 4. Pick a two digit hexadecimal number as the file identification number (ID) and store into memory location 17F9.

Do *not* choose ID = 00 or ID = FF.

For an explanation, please refer to Section 8.11.6, Hint #4.

Step 5. Refer to Figure 8.11.2 and note that the starting address for the Audio Tape Dump subroutine is 1800. Load 1800 into the field of the KIM-1. Do not press the **GO** key because the tape is not moving at this time.

Step 6. Select the RECORD mode of the tape recorder and wait a few seconds for the tape to start moving and to attain a constant speed.

Press **GO** key.

*Comment: The display will become dark for a time and then the display will light and show*

*0000 xx                      where xx is arbitrary*

*The amount of time during which the display is dark varies with the length of the program. In our example, which involves only about a dozen memory locations, the display is dark for about 16 seconds.*

The recording is complete as soon as the display relights.

Step 7. STOP the tape recorder.

Step 8. (optional) REWIND the tape cassette to its starting position.

Step 9. (optional) Listen to the tape. With experience, you will be able to identify the beginning 100 sync pulses, the program in the middle, and the terminating characters at the end of each record by their distinctive sounds.

#### 8.11.4 How to Read an Audio Tape into the KIM-1

Before we list the steps to be followed for reading from an audio tape into the KIM-1 microcomputer, let us list the assumptions we are making:

- a) The file identification number (ID) of the record we wish to transfer into the KIM-1 memory is known.
- b) The starting address of the Audio Tape Load (LOADT) subroutine is known to be \$1873.

The procedure for transferring information from an audio magnetic tape into the KIM-1 memory is:

(If not already on, turn on all power supplies.)

Step 1.\* Verify that the 6502 microprocessor is in the binary mode and not in the decimal mode. Please refer to Section 8.11.3, Step 1, for details on how to do this.

Step 2. Place the cassette into the tape recorder and REWIND the tape, if necessary, so that the tape will start to move and will attain a constant speed about 3 seconds before reaching the spot where the record with the desired ID exists on the tape.

Step 3. Store the known file identification number (ID) into memory location \$17F9.

Step 4. Load \$1873 (the starting address of the Audio Tape Load subroutine) into the address field of the display.

Step 5. Press the GO key.

*Comment: The display will become dark and the KIM-1 will search for a tape input with the specified ID number. The tape is not moving so the search is not successful.*

Step 6. Set the volume control on the tape recorder to approximately its midpoint position.

Step 7. Place the tape recorder into the PLAY mode. The tape should begin to move forward.

*Comment: As soon as the data record with the specified ID number is located and completely read in, the display will relight and show*

0000 xx                      where xx is arbitrary

Step 8. STOP the tape recorder after the display relights.

In reading an audio tape into a KIM-1, three types of problems may be encountered.

Trouble #1. In step 7, the display relights and shows FFFF xx. This means that the record with the specified ID was located but the check sum test failed as a result of either a record or a playback error.

FIX: Repeat the entire recording and playback procedures checking each step carefully. In a short program, step through all of the

\* If it were financially acceptable, we could add CLD instructions to the Audio Tape subroutines, manufacture new ROMs, and remove the need for Step 1.



memory locations involved and compare with your written program. If the problem still persists, refer to the KIM USERS MANUAL, Appendix C, "In Case of Trouble."

Trouble #2. In step 7, the tape continues to run to the end and the display remains dark. This usually means that the record with the specified ID number was not "found."

FIX: a) Verify that the record was actually on this cassette and on the correct side. (Cassettes have side A and side B recordings.) Repeat the entire recording and playback procedure checking each step carefully.

b) Actually listen to the recording through the speaker or an earphone. With some practice, you will be able to identify the beginning 100 sync pulses, the program in the middle, and the terminating characters at the end of each record by their distinctive sounds.

c) Verify again that the 6502 microprocessor is in the binary mode and not in the decimal mode.

Trouble #3. If you use the tape recorder to erase a tape, you may record a "noise" which seems to be related to a subharmonic of the clock frequency.

FIX: Disconnect the tape recorder from the KIM-1 board while using the tape recorder to erase the tape.

### 8.11.5 An Example of Record and Playback

The purpose of this example is to illustrate the procedure for recording programs and/or data on audio magnetic tape from the KIM-1 system, and vice versa.

**Example 8.11.1.** Record and play back into the KIM-1 memory a program which will add two decimal numbers located in two memory locations and will store the sum in a third memory location.

**Solution:** The source program code and machine code for this example are given in Table 8.11-1. Because the 6502 microprocessor has no stop, wait, or halt instructions, the last instruction in the program is an absolute jump to itself. Pressing the RESET key will cause the KIM-1 system to exit from this infinite loop. In subsequent examples, we will jump to a display subroutine.

**TABLE 8.11-1 Source Program Code and Machine Code  
for Adding Two Numbers and Storing the Sum**

Memory Map	Byte			SOURCE CODE			Comments
	1	2	3	LABEL	OP CODE	OPERAND	
00				NUM1			First number in location 00.
01				NUM2			Second number in 01.
02				SUM			Sum in location 02.
03	F8			PROGRAM	SED		Set decimal mode.
04	A5	00			LDA	NUM1	Load NUM1 into accumulator.
06	18				CLC		Clear carry.
07	65	01			ADC	NUM2	Add NUM2 to NUM1 with carry.
09	85	02			STA	SUM	Store SUM into location 02.
0B	4C	0B	00	HALT	JMP	HALT	Jump to itself over and over again.
0E							

Using the KIM-1 keyboard, load the program given in Table 8.11-1 into memory locations \$03 through \$0D. If you do not know how to do this, or have forgotten how, please refer to Section 8.6.

Now thoroughly test and verify that this program is correctly loaded and that it performs as expected. Be sure to load two numbers, run the program, and inspect the contents of memory location \$02 to see if the correct sum is present there.

**To record a program from the KIM-1 on audio tape.** Please refer to Section 8.11.3 and note that we have satisfied all of the assumptions stated there and we are ready to proceed with performing Steps 1 through 8. The results are summarized in Table 8.11-2.

**To playback the audio tape into the KIM-1.** The assumptions which we have made require us to know that:

- a) The desired tape record ID is \$11;
- b) The starting address of the Audio Tape Load subroutine is \$1873.

Recall that in this example, we have arbitrarily chosen \$11 as the file ID. However, the starting address \$1873 is fixed for the KIM-1 Monitor Program and will always have this value.

Please refer to Section 8.11.4 and note that we have satisfied all of the assumptions stated there and we are now ready to proceed with performing Steps 1 through 6. The results are given in Table 8.11.3.

**TABLE 8.11-2 Solution to Example 8.1 for the case of recording on an audio tape from the KIM-1 Microcomputer. The step numbers correspond with those given in Section 8.11-3**

<i>Press Keys</i>	<i>Display</i>	<i>Comments</i>
<b>[RS]</b>	xxxx xx	Step 1:
<b>[AD] [0] [0] [F] [1]</b>	00F1 xx	Places in binary mode by clearing decimal mode.
<b>[DA] [0] [0]</b>	00F1 00	
<b>[AD] [1] [7] [F] [5]</b>	17F5 xx	Step 2:
<b>[DA] [0] [3]</b>	17F5 03	Stores starting address 0003 into memory locations 17F5 (low byte) and 17F6 (high byte).
<b>[+] [0] [0]</b>	17F6 00	
<b>[+] [0] [E]</b>	17F7 0E	Step 3:
<b>[+] [0] [0]</b>	17F8 00	Stores ending address 000E into memory locations 17F7 (low byte) and 17F8 (high byte).
<b>[+] [1] [1]</b>	17F9 11	
<b>[+] [1] [1]</b>	17F9 11	Step 4: Loads ID = 11 into location 17F9.
<b>[AD] [1] [8] [0] [0]</b>	1800 A9	Step 5: Loads the starting address of the subroutine Audio Tape Dump.
On Tape Recorder <b>[RECORD]</b>		
wait about 3 seconds <b>[GO]</b>	dark (about 16 sec.) 0000 xx	Step 6: Program is being recorded. Recording is complete.
On Tape Recorder <b>[STOP]</b>		
	0000 xx	Step 7: Tape recorder is stopped.
	0000 xx	Step 8: (optional) Rewind tape.

**TABLE 8.11-3 Solution to Example 8.11.1 for case of reading from an audio tape into the KIM-1 microcomputer memory. The step numbers correspond with those given in Section 8.11.4**

<i>Press Keys</i>	<i>Display</i>	<i>Comments</i>
<b>RS</b>	xxxx xx	Step 1:
<b>AD</b> <b>0</b> <b>0</b> <b>F</b> <b>1</b>	00F1 xx	Places KIM-1 into binary mode
<b>DA</b> <b>0</b> <b>0</b>	00F1 00	by clearing the decimal mode.
	00F1 00	Step 2: Place cassette into tape recorder and rewind tape, if necessary.
<b>AD</b> <b>1</b> <b>7</b> <b>F</b> <b>9</b>	17F9 xx	Step 3:
<b>DA</b> <b>1</b> <b>1</b>	17F9 11	Load ID = 11 into memory location \$17F9.
<b>AD</b> <b>1</b> <b>8</b> <b>7</b> <b>3</b>	1873 A9	Step 4: Loads the starting address of the subroutine Audio Tape Load.
<b>GO</b>	dark	Step 5: KIM-1 is searching for ID = 11.
On tape recorder	dark	Step 6: Set volume control to midpoint.
<b>PLAY</b>	dark	Step 7: Places recorder into PLAY mode and tape begins to move.
	(about 16 sec.) 0000 xx	Playback is complete.
<b>STOP</b>		Step 8: Stops tape recorder.

### 8.11.6 Some Helpful Tape Hints

The following hints are offered to help you make the most effective use of your cassette tapes without encountering some of the more commonly associated problems:

**Hint #1.** Voice messages may be added between the record data blocks on the tape. The KIM-1 system will ignore these audio messages when the tape is read back provided the voice messages occur only between record blocks. However, you will have to install an earphone or speaker in parallel with the KIM-1 audio tape data pin in order to hear the voice messages.

**Hint #2.** If more than one data block is to be recorded per cassette side, proceed as follows:

*Case 1.* Data blocks are to be recorded in sequence without re-winding between blocks.

You need only to specify the parameters of each new block (ID, SAL, SAH, EAH, EAL) and proceed with recording each block following the standard steps listed in Section 8.11.3. It is best to use unique ID numbers for each block.

*Case 2.* Tape has been rewound or the location of the end of the last data block has been lost.

You must know the ID number of the last recorded data block. Rewind the tape to well before the anticipated starting point and set up the parameters to read the last data block following the steps given in Section 8.11.3, Step 4.

If the data transfer is successful (as indicated by the display showing 0000 xx), you may proceed to load the next data block. Refer to Section 8.11.3.

**Hint #3.** Avoid placing a data block between two existing data blocks or in an area of the tape which has previously been used for recorded data. You may recall that data blocks may be of arbitrary length and hence variations in block length and tape speed may result in overlapping of recorded data blocks.

**Hint #4.** Recall that when *recording* on a magnetic tape, you cannot select ID = 00 or ID = FF. However, when *reading* from a magnetic tape into the KIM-1 memory, we can use these ID numbers as follows:

*Case 1.* If we select ID = 00, the ID number recorded on the tape will be ignored and the system will read the first valid data block encountered on the tape. The data read from the tape will be loaded into memory addresses as specified on the tape.

*Comment:* This is useful if you forget your ID number or if the ID number was not provided to you. Note that you still need the original starting and ending addresses in order to locate and execute the program.

*Case 2.* If we select ID = FF, the ID number recorded on the tape will be ignored and the system will read the first valid data block encountered on the tape. In addition, the data block will be loaded into successive memory locations beginning at the address specified in locations 17F5 and 17F6 (starting address low, SAL, starting address high, SAH) instead of the locations specified on the tape.

*Comment:* This is useful if you forget, or do not know, the ID number and the starting address. Note that you still should know the maximum length of the program so that you don't "write over" and destroy other data stored in memory.

### 8.12.1 Manipulating the KIM-1 Display (Example 8.12.1)

The KIM-1 display is controlled by the 6530-002 I/O ports. A diagram of the KIM-1 display and 6530 RAM ROM I/O interface is shown in Figure 8.12.1. This diagram is simplified to the point where it does not show any of the drivers, inverters, or pin connections. The primary purpose of Figure 8.12.1 is to help to explain how the KIM-1 six character, 7 segment LED display may be used to display many of the alphabetical characters. In addition, it will be found that it is possible to blink and to shift the characters which are displayed.

The A-side data determines which of the seven segments are illuminated for the character selected by the B-side data. Segments A through G are selected by bits 0 through 6, respectively.

**Example 8.12.1.** Determine A-side and B-side data which will cause

E	r	r	o	r
---	---	---	---	---

to be displayed in the first four character positions.

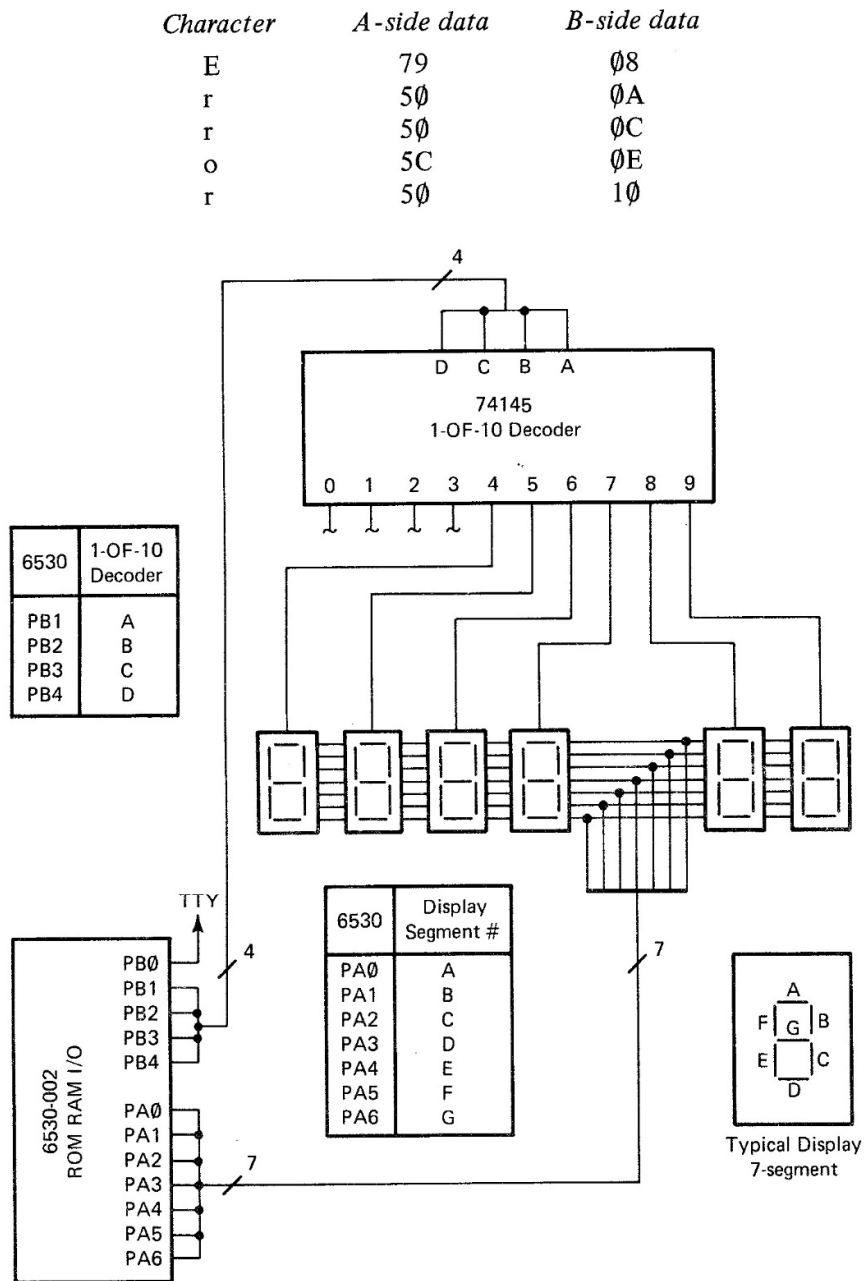


Figure 8.12-1 Simplified diagram of KIM-1 display and control by the 6530-002

**Comment 1.** A-side data are determined as follows:

	<i>Segments On</i>							<i>Hex Code</i>
	<i>G</i>	<i>F</i>	<i>E</i>	<i>D</i>	<i>C</i>	<i>B</i>	<i>A</i>	
Character E	1	1	1	1	0	0	1	79
Character r	1	0	1	0	0	0	0	50
Character o	1	0	1	1	1	0	0	5C

**Comment 2.** Computation of the B-side data may appear to be confusing at first. Because PB1, PB2, PB3, PB4 (instead of the more typical PB0, PB1, PB2, PB3) are sent to the 1-OF-10 decoder, the decoder supplies an output which is only one-half of the actual input. In other words, we have to supply an input which is twice as large so that character positions 04, 05, 06, 07, 08, 09 actually require as B-side data inputs: 08, 0A, 0C, 10, and 12.

The specific memory addresses associated with the 6530-002 I/O ports are as follows:

1740	SAD	A-side Data
1741	PADD	A-side Data Direction
1742	SBD	B-side Data
1743	PBDD	B-side Data Direction

A clearer perspective of the relative memory locations may be obtained by referring to Figure 8.12-2 which is Figure 8.4-1 with the above memory locations added. It may be seen that some pins on both data direction registers must be set for output. (Remember 1 = output and 0 = input.) Just exactly which pins ports A and B are set for input and which pins are set for output is given in Figure 8.12-3.

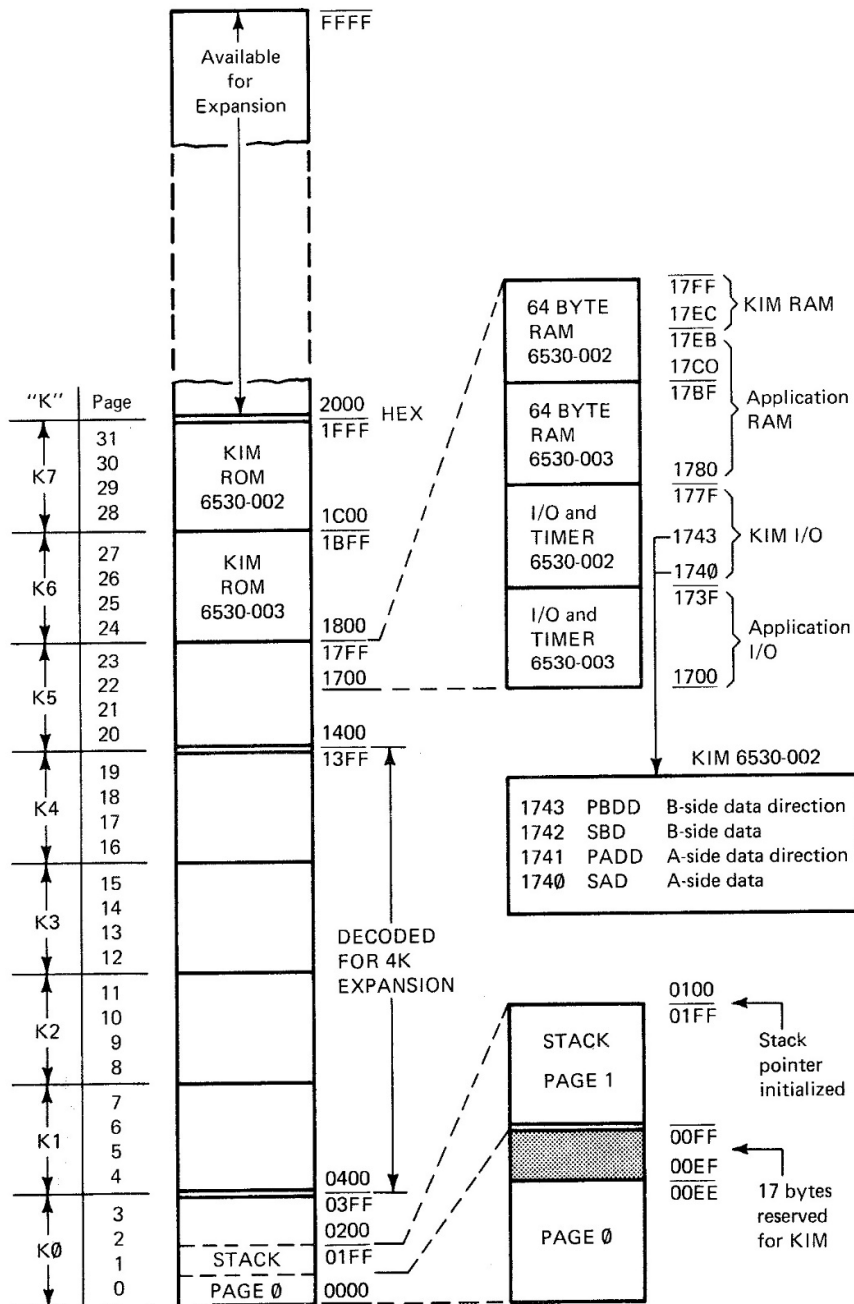
A program which will load the A and B data direction registers with 7F and 1E respectively, is as follows:

<i>Op Codes</i>	<i>Source Codes</i>
A9 7F	LDA #7F
8D 41 17	STA PADD
A9 1E	LDA #1E
8D 43 17	STA PBDD

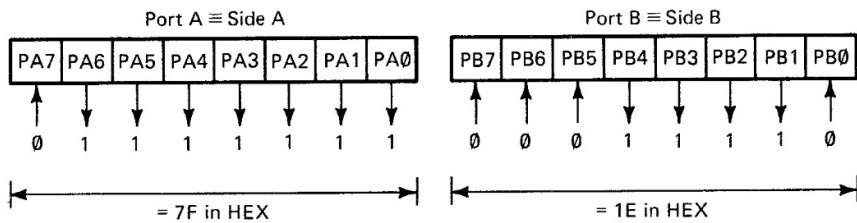
**Example 8.12.2.** Write a program which will display the letters "EE" in the two left-handmost characters of the KIM-1 display. The remaining four characters are to be dark.

The solution to Example 8.12.2 is shown in Table 8.12-1.





**Figure 8.12-2** MEMORY MAP with 6530-002 ports A and B added  
(Courtesy MOS Technology)



**Figure 8.12-3** Input/Output pin configuration for ports A and B in Example 8.12.1  
1 = output; 0 = input

**TABLE 8.12-1** Source and Op Codes which display the letters "EE" in the KIM-1 display

Memory Map	Byte 1	Byte 2	Byte 3	LABEL	OP CODE	SOURCE CODE OPERAND	Remarks
0030	A9	7F			LDA	#7F	Set output pins on
32	8D	41	17		STA	PADD	Port A.
35	A9	1E			LDA	#1E	Set output pins on
37	8D	43	17		STA	PBDD	Port B.
3A	A9	79			LDA	#79	Load the letter E.
3C	8D	40	17		STA	SAD	
3F	A9	08		AGAIN	LDA	#08	Display the letter E
41	8D	42	17		STA	SBD	in the first position.
44	A9	0A			LDA	#0A	Display the letter E
46	8D	42	17		STA	SBD	in the second position.
49	4C	3F	00		JMP	AGAIN	Jump back and display EE again.

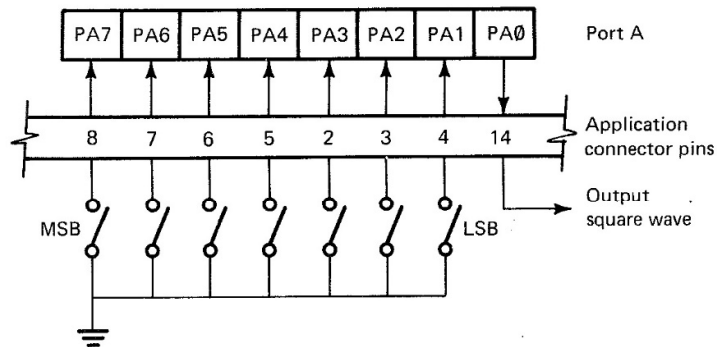
### 8.12.2 Generating a Variable Frequency Square Wave (Example 8.12.3)

The purpose of this example is to illustrate one of the many different ways of defining the Input/Output interface ports of the KIM-1 microcomputer for both input and output. In order to be more specific and practical as well, we will use Example 8.12.3, which may be stated as follows:

**Example 8.12.3\*.** Write a computer program for the KIM-1 which will generate a variable frequency square wave with amplitudes of 0 and +5 volts.

As shown in Figure 8.12-4, the eight pins of Port A are programmed such that one pin (PA0) is the output and the other seven pins (PA1 to and including PA7) are connected to seven switches which serve as the input control to determine the frequency of the square wave.

It is further desired that the switch connected to PA1 is to be the least significant bit (LSB) and the switch connected to PA7 is to be the most significant bit (MSB). We would also like each switch to act as a 1 when closed, and as a 0 when open. If this is done, the status of the switches define binary numbers from zero (all switches open) to  $127_{10}$  (with all switches closed).



**Figure 8.12-4** Input/Output connections for the variable frequency square wave generator (Example 8.12.3)

**A Solution to Example 8.12.3.** One of the many possible solutions to Example 8.12.3 is given in Figure 8.12-5, which presents in a modified form, both the assembly language program as well as the machine code. It is said to be “modified” form because of the additional column labeled MACHINE CYCLES. This column has been added because of the importance of determining exactly how much time (number of machine cycles) are required to execute each line of code.

\* Example 8.12.3 is similar to the “Real Application” example given on page 55 of the KIM-1 USER MANUAL. (MOS Technology Inc.)

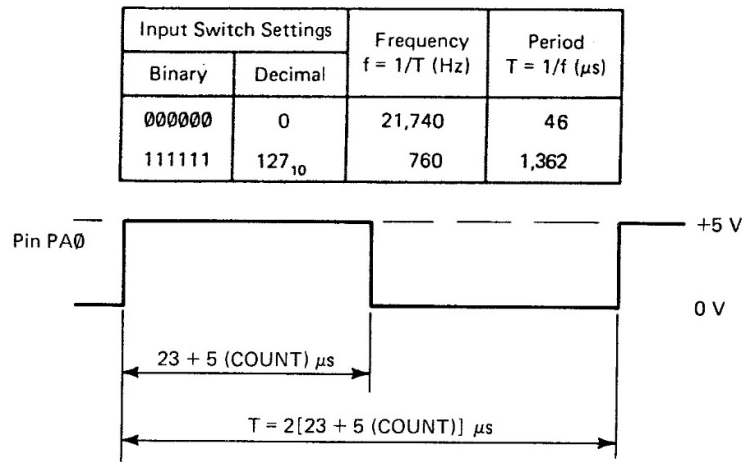
INSTRUCTION				MACHINE CODE				Remarks
Memory Map	Byte 1	Byte 2	Byte 3	MACHINE CYCLES	LABEL	OP CODE	OPERAND	
0200	A9	01		2	INIT	LDA	#\$01	Define I/O. PA0=Output=1. PA1-PA7=Input=0.
0202	8D	01	17	4		STA	PADD	PADD = Port A Data Direction Register
0205	EE	00	17	6	START	INC	PAD	Toggle PA0 pin only on Port A.
0208	AD	00	17	4	READ	LDA	PAD	Read switch settings into accumulator.
020B	49	FF		2		EOR	#\$FF	Complement switch value because closed = 1.
020D	4A			2		LSR	A	Shift accumulator one bit to right.
020E	AA			2		TAX		Transfer COUNT into X-Index Register.
020F	CA			2	DELAY	DEX		Loop to delay by amount specified by
0210	10	FD		3,2	BPL	DELAY		COUNT stored in the X-Index Register.
0212	30	F1		3		BMI	START	Go back to START.
PADD = \$1701								Defines absolute address of Port A Data Direction Register.
PAD = \$1700								Defines absolute address of Port A Data Register as specified by the KIM Monitor.

**Figure 8.12-5** Assembly language, Machine code, and MACHINE CYCLES for a variable frequency square wave generator (Example 8.12.3)

**Execution of Example 8.12.3.** Enter the program as usual, but before executing the program, load the NMI vector locations 17FA and 17FB with 00 and 1C respectively, so that the **[ST]** key will function properly. The program begins in memory location 0200, so load this memory location, and press the **[GO]** key. The display will remain dark and the output square wave will be available at pin 14 of the Applications Connector. This output may now be connected to an audio amplifier to serve as a tone generator; or it may be applied to the horizontal input of an oscilloscope and used to calibrate the horizontal sweep of the oscilloscope.

**The Output Square Wave.** The square wave output is shown in Figure 8.12-6 along with equations for the frequency and period of the square wave as a function of the switch settings. When numerical values are substituted, the resulting frequencies and periods are:

Input Switch Settings		Frequency ( $f = 1/T$ ) Hz.	Period $T = 1/f$ $\mu\text{sec.}$
Binary	Decimal		
0000000	0	21,740	46
111111	127 <sub>10</sub>	760	1362



**Figure 8.12-6** Output square wave for Example 8.12.3. Note that the period is  $T$ , the frequency is  $1/T$ , and COUNT is determined by the input switch setting with  $0 \leq \text{COUNT} \leq 127_{10}$

## PROBLEMS

1. Use the KIM keyboard and display to inspect the contents of memory locations (in hex) 0000, 0400, 1800, and FFFF. Record the contents as observed. Refer to Figure 8.4.1, the memory map for the KIM-1 microcomputer, and the 6530 Monitor listings to explain your observations.
2. List the contents of the following addresses: 1C00 through 1C0A, 1746, 0000 through 0010, and 1780 through 1790.
3. Sketch and label all MCS6502  $\mu$ P registers. Briefly describe their functions.
4. List the addresses where the following can be found in KIM: ROM, RWM, TIMERS, and I/O.
5. Use KIM keyboard and data keys to clear memory locations \$00EF through \$00FE. If this is not possible, explain.
6. Where are the internal registers (A, X, Y, S, P and Z) stored during a single-step operation?

7. Explain why it is necessary to load memory locations 17FA and 17FB before the KIM-1 microcomputer may be used in the single-step mode. Refer to 6530-002 listing between lines 600 and 620 and correlate with the use of the above two memory locations.
8. Refer to Table 8.7-2 which presents the single step execution of Example 8.1 and fill in the table below for each step.

PC	P	A	X	Y	S
----	---	---	---	---	---

9. Study the program below and understand how it works. (There is at least one error.) The intention of the program is to set the contents of locations \$0000 through \$00FF to all ones (\$FF).

ADDR	CODE	PROGRAM
0200	AE FF	LDX #\$FF
0202	18	CLA
0203	95 00	STA 0, X
0205	CA	DEX
0206	D0 FB	BNE LOOP
0208	4C 4F 1C	JMP \$1C4F

10. If you use the KIM-1 microcomputer in the binary mode, what result would you expect if you added 34 and 27? What would you expect in the decimal mode? Use the solution to Example 8.1 and verify your answers using the KIM-1 microcomputer.
11. Modify the source code for Example 8.1 which is given in Table 8.3-1 so as to check for possible overflow. What would you suggest doing in the event of an overflow?
12. Modify the source code for Example 8.1 which is given in Table 8.3-1 so as to display the sum instead of "jumping to itself" at the end of the program. Display the sum in the four leftmost locations of the display. You will need to study the Monitor Listing for the 6530s in order to do this.
13. Modify the source code for Example 8.1 which is given in Table 8.3-1 for the case where the two positive numbers are each three bytes long. Your program should check for overflow and allow enough memory locations to store the sum. Use the decimal mode. Repeat for binary mode.
14. Write a program for execution on the KIM-1 microcomputer which will output all of the printable ASCII characters to a printer such as the Teletype Model ASR33. In the next line printed, move

all of the characters one character to the right and bring the character shifted off to the right back to the leftmost position. (In other words, circular shift right.) Repeat printing new lines until all characters have been printed in all possible positions. (This is a program which printer technicians find very useful when making printer adjustments.)

15. Punch a paper tape for Example 8.2 and compare to the punched paper tape shown in Figure 8.10-6. Explain any differences.
16. Write a program for execution on the KIM-1 which will display the characters as they are “read in” from a magnetic tape recorder. Display the character which has just been read in the leftmost position of the display and then shift the characters to the right to make room for each new character. (The characters will appear at a rate which may be too fast to permit reading every character reliably.)
17. Write a program for execution on the KIM-1 which will generate a very long string of sync characters, an ASCII SYN character (\$16). Record approximately five minutes of only sync characters on a magnetic tape and then play them back using the program written in Problem 16. Vary the volume control and tone control (if your recorder has one) and determine the range for which synchronization is maintained. Determine the optimum settings for your tape recorder.
18. Modify the program given in Table 8.12-1 for Example 8.12.2 so that the two characters E E blink on and off at the rate of about once per second. Although the timers in the 6530s might be used, write a program which will accomplish the same thing. (Not all microcomputers have timers.)
19. Modify the program given in Table 8.12-1 for Example 8.12.2 so that six characters (all of which may be different) will be displayed on the KIM-1 display. (Caution: Be sure to turn off a character before proceeding to turn on the next character.) In order to demonstrate that your program works, display the six characters E, R, R, O, R, and a blank.
20. Modify the program given in Table 8.12-1 for Example 8.12.2 so that more than six characters may be displayed at the rate of six at a time and then shifted one character to the right and displayed again in an endless loop.

21. Add steps to your program for Problem 20 so that the display blinks on and off at the rate of about once per second with on and off times being equal.
22. The solution to Example 8.12.3 in Figure 8.12-5 provides a frequency range from 760 to 21,740 Hz. What changes would be necessary to provide frequencies up to 50,000 Hz? What about for frequencies below 760 Hz?
23. Modify the program given in Figure 8.12-5 for Example 8.12.3 so that the positive portion of the square wave represents only 10% of the period instead of 50%.



# B

## **OPERATING PRINCIPLES OF THE KIM-1 MONITOR AND ON-BOARD I/O HARDWARE**

### **B.1 INTRODUCTION**

In Chapter 8, detailed methods for operating the KIM-1 system and associated peripheral units are described. This appendix will present some of the principles underlying these methods of operation. There are two principal reasons for such a presentation. First, the KIM-1 is an excellent example of a microprocessor-based system, and the underlying design principles present a useful case study for one who aspires to do such design himself. Second, the KIM-1 user will find an understanding of these principles very helpful in actual operation. For example, the user who wishes to use the on-board peripherals, such as the keyboard and the display for nonmonitor functions, will find an understanding of monitor subroutines an invaluable aid.

Figure 4.2-1 is a flow chart of the KIM-1 monitor program. It will serve as an important reference for most of the discussion in this Appendix.

Another important reference is the KIM-1 monitor listing itself, which is included herein as Section B.8.

### **B.2 SINGLE STEP OPERATION OF THE KIM MICROCOMPUTER**

The normal mode of operation for any digital computer, including the KIM, is that in which instructions are executed sequentially at high speed. Automatic high-speed sequencing accounts for most of the unique capability of such systems. It is often desirable, however, to execute a sequence of instructions at manual speed, either to examine step-by-step program execution and thereby verify program integrity,

or to diagnose reasons for program malfunction, i.e., to perform "debugging." The KIM system permits manual sequencing via the keyboard switch labeled SST (single step).

The 6502 CPU chip used in the KIM system has several built-in features which make single-step operation convenient to perform. One such feature is the control output designated as *SYNC*. The 6502 produces a pulse from this terminal during each machine cycle which, if completed, constitutes the fetch of an instruction op code from memory. Since such a cycle signals the end of an instruction execution, the ability to stop the computer when *SYNC* occurs would accomplish the single-step operation desired (assuming, of course, that the system can be restarted).

One other feature of the 6502 which could be used to accomplish single-step operation is the *RDY* (READY) control input. If, during the PHASE ONE clock of any machine cycle when a memory read is being performed, the *RDY* input is pulled low, the completion of the cycle will be inhibited at PHASE TWO clock time. (A memory write cycle will not be inhibited, but the subsequent read cycle will be.) The principal use of the *RDY* control is in managing transfers between the CPU and slow memory or I/O devices which may not respond to read requests with sufficient speed to keep up with the 6502 clock. Such devices can, by properly energizing the *RDY* input, force the CPU to wait until it is ready to deliver the requested data. A second use of this input is implemented by inverting and feeding back the *SYNC* pulse to the *RDY* input. This will cause the system to stop each time a *SYNC* pulse occurs. Therefore the system stops after the execution of each instruction. (Actually the *SYNC* pulse must be stored in a separate flip-flop. The flip-flop will inhibit system operation until it is cleared, at which time the system will fetch and execute the next instruction. Refer to Section 3.1, page 124, of the MOS HARDWARE MANUAL, reference 2, for a detailed description of one mechanization of this scheme.)

The KIM system does not use the *RDY* control to achieve single-step operation. Recall that one of the principal motivations for establishing single-step operation was to permit a leisurely examination of the result of executing each instruction. The KIM monitor affords a powerful mechanism for performing such an examination because the KIM system implements a procedure whereby, after each step of the program under study is executed, system control is passed back to this

monitor. The monitor is then used to examine register contents, data values, etc. To understand the actual mechanization built into the KIM system, refer to Figure B.1. Gate U26, a two-input NAND gate, has as its inputs SYNC from the 6502 and a signal labeled K7. For the time being, assume that K7 is high or at the "1" level. If the single-step switch SST is closed, the occurrence of SYNC will pull the NMI input low, constituting a nonmaskable interrupt input to the 6502. Service of this interrupt is in the form of execution of the monitor routine which lights the display and scans the keyboard. This mode reflects the transfer of control to the monitor as described above. Pushing the GO key on the keyboard will cause other monitor routines (see the discussion of the KIM keyboard below) to be called. These routines will fetch the next user instruction, giving the desired result of stepping through one more step. At this point the process will be repeated.

The perceptive reader will have detected a logical flaw in the above operation. He will ask, "Since the execution of the monitor program is conceptually identical to the execution of the user program, why doesn't the system halt after the execution of each step in the monitor program?" This is the point where the role of the K7 signal must be considered. Examination of the logic diagram will confirm that the K7 signal is produced when an address in *page 7* (address range 1C00 to 1CFF) is present on the KIM address bus. Since the 74145 decoder which generates K7 produces active-low outputs, when such an address is present, the K7 signal is at a low level. Thus the NAND gate driving the NMI input will not produce its active low output when a page 7 address is on the address bus, even though the SYNC signal denoting completion of instruction execution is generated. Since the keyboard/display monitor is located in page 7, it will be executed in uninterrupted fashion. This also has the effect of preventing one from stepping through any of the monitor routines located in page 7.

One final point about single-step operation in the KIM should be mentioned. As discussed more fully in Section B.3 on RESET and STOP operation in the KIM, proper execution of the STOP routine in response to an NMI request demands that the user, as a necessary step in system initialization, load the proper interrupt vectors into the system. Since single-step operation depends on this routine, such loading is necessary to properly enable this mode of operation. In plain English, this means that locations 17FA and 17FB must be loaded with address 1C00, the beginning location of the STOP monitor routine.

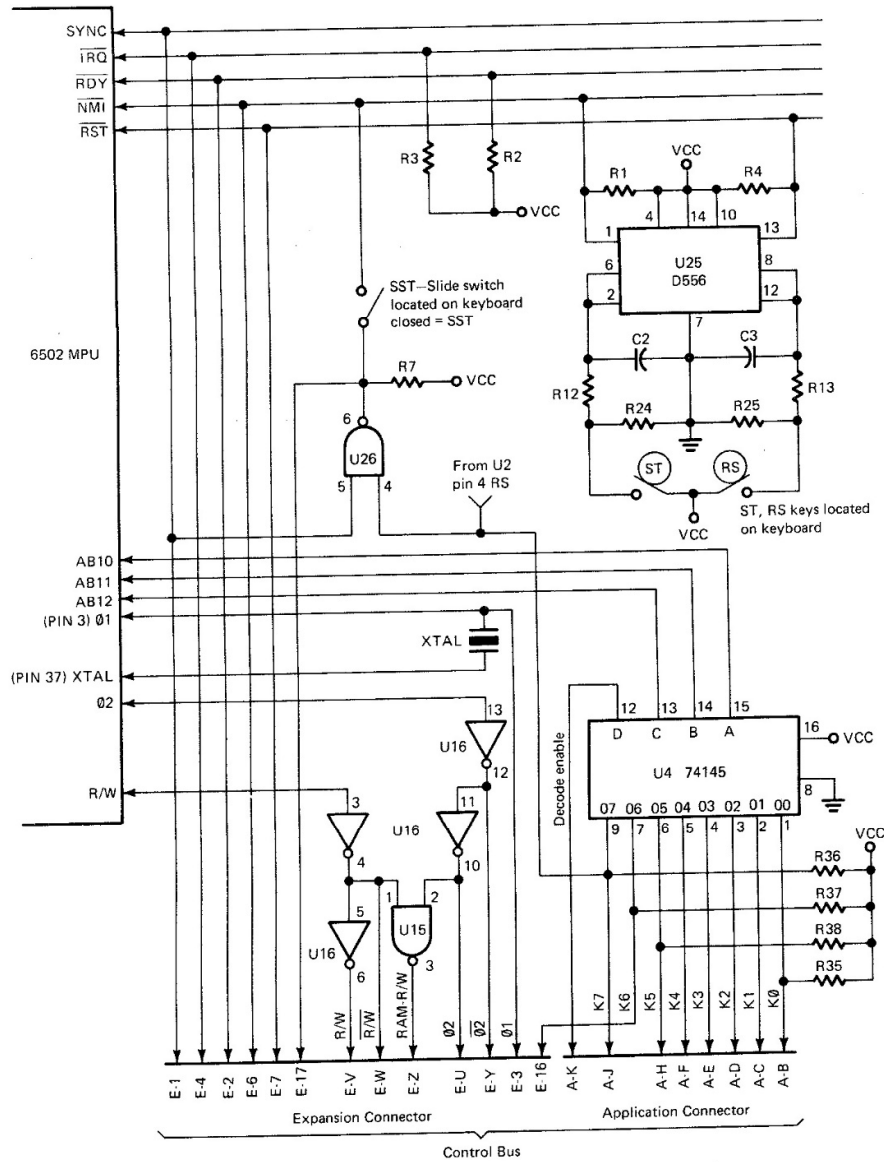


Figure B.1 KIM-1 control and timing circuitry

### B.3 RESET AND STOP OPERATIONS

These two operations bear considerable similarity to one another and will therefore be discussed together. Each of them is initiated by depressing a keyboard key. Referring to the logic diagram of Figure B.1 it may be observed that pushing the ST (STOP) key introduces a positive-going voltage onto pins 2 and 6 of U25, a dual timer. This circuit's characteristics are such that this input will produce a negative-going transition at pin 1, which is delivered to the NMI input of the CPU. Before examining the effect of this input, let us digress and point out the principal function of the dual timer circuit. Its characteristics are such that, even if the ST keystroke manifests *switch bounce* (which is typically the case), the timer output will respond only to the *initial positive transition* at its inputs, effectively debouncing the switch.

Let us return now to an examination of the sequence of events which occur when the NMI input is energized. This constitutes a *nonmaskable interrupt* request and results in a somewhat feverish interval of activity by the CPU. The steps which constitute this activity are:

1. When the NMI request arrives, the current contents of the program counter and the status (P) register are pushed onto the system stack. This allows return to the interrupted routine when the interrupt has been serviced.
2. The contents of memory locations 1FFA and 1FFB are fetched from memory and placed in the program counter. This is an automatic step performed by the 6502 hardware in response to the NMI request. Locations 1FFA and 1FFB are ROM locations, which collectively contain the address 1C1C. (The actual addresses are FFFA and FFFB, but, in the KIM-1, the high-order three address bits are ignored in an unexpanded system.)
3. Normal processor activity resumes. This activity consists of placing the program counter contents on the address bus and performing a memory read. The contents of location 1C1C are thus fetched as an instruction op code.
4. Location 1C1C, also a ROM location, contains 6C, a *jump indirect* op code. The next two locations, 1C1D and 1C1E, are thus accessed by the CPU, producing address 17FA, also preprogrammed in ROM. The nature of the jump indirect instruction is such that this address (along with the next consecutive location 17FB) contains the actual address of the next instruction to be executed. Since 17FA and 17FB are RAM locations,

they must be previously loaded by the system user with the address of the first step in the NMI service routine he wishes to call. The STOP routine provided with the KIM monitor begins at location 1C00, and, if this is the desired routine, that value should have been loaded into 17FA and 17FB during system initialization.

5. The CPU begins execution of the STOP routine. Its details will not be discussed, but it generally consists of a continuous scan of the keyboard combined with a continuous display of the last memory location referenced by the user program along with the data stored at that location. (See the monitor listing of Section B.8 for details.)

This rather lengthy list of activities can be confusing and some care is necessary to sort out some of the terminology encountered. The various elements of 6502 literature refer to the contents of addresses FFFA and FFFB (which translate into 1FFA and 1FFB in limited address space of the unexpanded KIM) collectively as an *interrupt vector*, since they point into memory toward the requested service routine. In the particular case of the KIM, however, the desire to permit the system user to choose not only the monitor STOP routine but any other routine of his choosing as a response to an NMI request results in several of the extra steps outlined above. Since the 17FA, 17FB locations contain the actual address of the first location of the desired service routine, the contents of these locations are referred to in the KIM literature as the NMI interrupt vectors.

The RST (RESET) key on the KIM results in a sequence of activity which closely resembles that discussed above. The other half of the U25 dual timer debounces the switch, applying a request input to the RST input of the 6502. Once again, a sequence of activity ensues, involving fetching interrupt vectors from locations 1FFC and 1FFD. These locations contain the address 1C22, which, when sent to memory as a request for op code, yields not a jump indirect op code, but the actual first instruction in the RESET routine, which is a system initialization routine.

It is interesting to reflect on the fundamental difference between the NMI and RST responses to service requests. The most obvious difference is that the user has some flexibility in the use of the NMI routine, since by loading interrupt vectors of his choice into the appropriate RAM locations, he can control the response to such requests. It may seem that this flexibility is missing relative to the RST routine because the KIM designers saw no need for substitution of a user-

defined response in place of the one provided by the system monitor. The actual reason for the difference is, however, more fundamental. Note that the resources of the monitor program must be used to load the appropriate interrupt vectors into the NMI locations. In order to call upon these resources, the system must enter a state in which a scan of the keyboard is performed. Once this state is entered, the keys can be used to call monitor routines to perform desired activities. Placing the system into this state requires the use of the system initialization routine. If this routine also depended upon the preloading of interrupt vectors, note that, somewhat embarrassingly, there would be no way to perform this preloading. The fixed interrupt vectors for RESET are therefore necessary in order to initiate proper system activity. This general procedure is one of several types most often described as *bootstrap procedures*.

## B.4 OPERATION OF THE KIM KEYBOARD AND DISPLAY

Chapter 8 describes in some detail the operational manner in which the user can interact with the KIM system via the keyboard and display units. Beyond considerations of simply being able to operate the KIM system, however, it is of considerable educational value to examine in some detail the methods by which these system elements are interfaced to the KIM system, because they constitute excellent examples of peripheral units and their interface to a microprocessor system.

As the management of the display and keyboard interfaces is discussed, reference will be made to the routines which comprise the KIM operating system. The reader will find it useful to refer to the listing of these programs (included in Section B.8) as the discussion proceeds, not only to understand the actual KIM software mechanization, but also to learn by example how such software is written. The KIM designers achieved considerable efficiency by sharing one I/O port for the dual functions of display output and keyboard input and by sharing keyboard scan routines between manual keyboard and Teletype interfaces. This achievement of efficiency is, unfortunately, at the expense of simplicity of understanding in some cases, particularly to one who is examining such methods for the first time. An attempt will be made in this discussion to avoid some of the detail of the implementation which tends to obscure a basic understanding of the system. For example, in the discussion, the display drive mechanization will be separated from the keyboard scan mechanization.

### B.4.1 Display Drive in the KIM System

Figure B.2 shows the principal elements in the display drive system. The diagram is simplified by omitting the electrical interface elements. The logical information flow is correct. The peripheral interface adaptor labeled 6530-002 is the key I/O element in this system. The display itself consists of six seven-segment display elements and, as configured, may conveniently be thought of as a two-dimensional array. Control over which *character* is displayed is provided by the A port of the 6530. Lines PA0 through PA6 are routed (through display drivers not shown in the diagram) to the seven LED segments (A through G, as shown in the figure) of each display unit. The other display selection dimension is activated by a 74145 BCD decoder as shown. The active output of this decoder (active-low) is a function of the four-bit code at the ABCD inputs. Any one of 10 binary codes (regular BCD decoding) produces an output from one of the decoder terminals. Thus, to drive a particular display element, the CPU must place the appropriate four-bit code in the SBD (Side B Data Register) positions 1 through 4. If these 4 bit positions are programmed as outputs (by appropriate loading of the Port B Data Register, PBDD), then this code is present on the PB1 through PB4 lines, and will appropriately drive the display array.

Output positions 0 through 3 of the decoder of Figure 4.3 are not used in the display drive. They function as part of the keyboard scan system and as part of the Teletype/keyboard selection control.

Table B.1 indicates the data which must be loaded into the B port Data Register (SBD) to appropriately select any one of the display elements.

**TABLE B.1 KIM Display Character Position Selection Codes**

<i>To Select Character Position (Left-to-Right)</i>	<i>Active Decoder Output Must Be</i>	<i>Decoder Inputs Must Be D C B A</i>	<i>Port B Data Register Contents Must Be* (In Hex)</i>
1	4	0 1 0 0	08
2	5	0 1 0 1	0A
3	6	0 1 1 0	0C
4	7	0 1 1 1	0E
5	8	1 0 0 0	10
6	9	1 0 0 1	12



\*Since only bits 1 through 4 of the Port B data register affect the decoder, there are a number of other hex values which could be used in each of these positions. The only requirement is that bits 1 through 4 must be identical to the pattern for the values shown.

Since only output positions 1 through 4 affect the display, there are a number of redundant combinations which may be used. The hex values shown in the table are those which assume that 0s are to be used in output positions 0, 5, 6, and 7. (Actually, the 6530-002 has no PB6 output.) In the KIM, these other outputs are used for other interfacing functions and it may be that in some cases some value other than 0 might be desired.

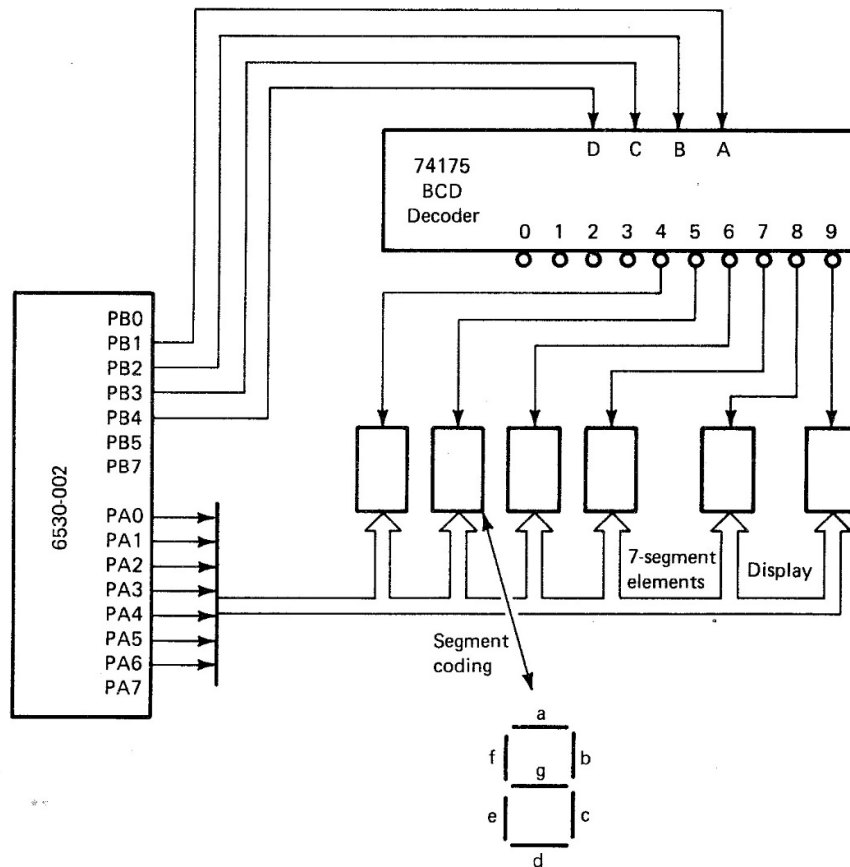


Figure B.2 Simplified block diagram of KIM-1 display drive interface

Table B.2 is a listing of some of the possible 128 characters which can be produced by a seven-segment display. In each case, the ABCDEFG element combination which produces the character is expressed in terms of the hex value which, when produced by the PA0-6 output bus, will light the appropriate segments. Once again, since the PA7 output does not participate in the display drive, it is assigned the value 0 in the table. Since this terminal is used by the KIM system as a serial Teletype input port, its actual output assignment is immaterial.

In summary, to display a desired character in one of the display positions, the following steps must be performed:

1. Program the peripheral Data Direction Registers PADD and PBDD to configure PA0-6 and PB1-4 as outputs.
2. Load the B side Data Register (SBD) with a value appropriate for the selection of the desired display element, as given in Table B.1.
3. Load the A side Data Register (SAD) with a value appropriate to the desired character, as given by Table B.2.





































A program for performing the above steps is quite simple to write, but is very limited in application, since only one display position is affected by it. A more complex program is used by the KIM monitor, and is designated as SCAND (for Scan Display).

#### **B.4.2 Scand, The KIM Display Drive Program**

This program is written as a subroutine and begins on line 1057, page 25, of the KIM monitor listing (Section B.8). The first three steps of the program fetch a byte (two characters) to be displayed and stored in INH, the location reserved for the high byte of the input buffer. The next two steps program the A side of the 6530 port as an output. Next index register X is loaded with 09 to serve as a pointer to the rightmost digit position of the display. Index register Y is loaded with 03 preparatory to counting the number of characters to be displayed (6 characters, three bytes). Then routine SCAND1 is entered.

SCAND1 starts by fetching the first byte to be displayed, shifts it right four positions to isolate a single character, and then jumps to subroutine CONVVD (Convert Display) which starts at line 1085. CONVVD searches a conversion table (line 1202) similar in function to Table B.2. It places the appropriate data value found in the table on the PA0-6 outputs, moves the preset value of index register X (set to 09 previously) to SBD, selecting the rightmost digit position of the display, and exe-

TABLE B.2 KIM Display Character Codes

<i>Character</i>	<i>Symbol Displayed</i>	<i>Segments Lit</i>	<i>Hex Code</i>
1		BC	06
2		ABDEG	5B
3		ABCDG	4F
4		BCFG	66
5		ACDGF	6D
6		ACDEFG	7D
7		ABC	07
8		ABCDEFG	7F
9		ABCDFG	6F
0		ABCDEF	3F
A		ABCEFG	77
B		CDEFG	7C
C		ADEF	39
D		BCDEG	5E
E		ADEFG	79
F		AEFG	71
G		?	?
H		BCEFG	76
I		BC	06
J		BCDE	1E
K		?	?
L		DEF	38
M		?	?
N		?	?
O		CDEG	5C
P		ABEFG	73
Q		?	?
R		EG	50
S		ACDFG	6D
T		EFG	70
U		CDE	1C
V		?	?
W		?	?
X		?	?
Y		BEFG	72
Z		?	?

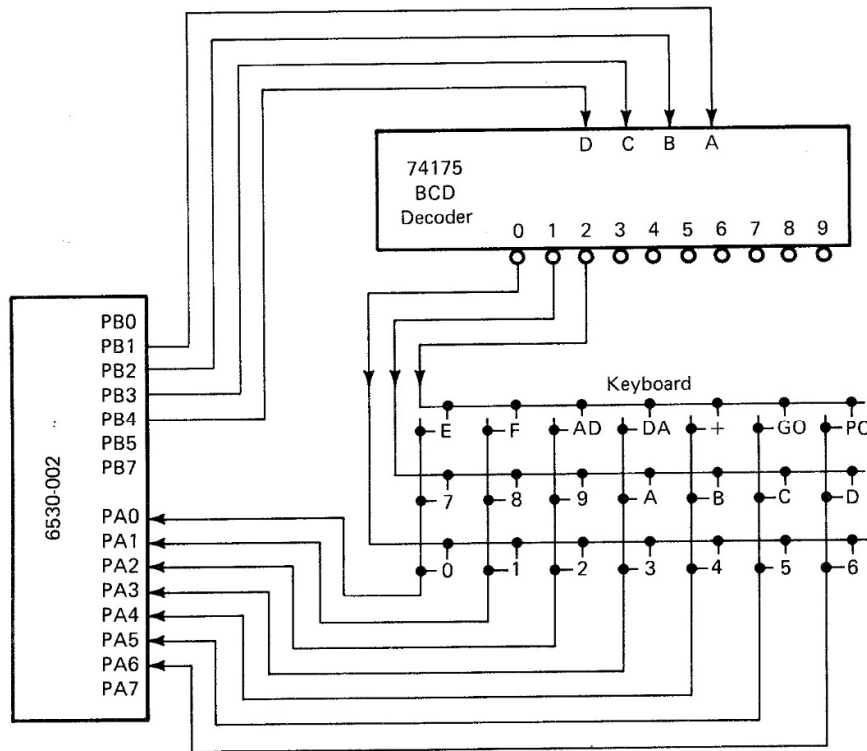
cutes a delay routine which lights the display for 500 instruction cycles. CONVD finally increments index X by 2 (to prepare a pointer to the next leftmost digit position the next time around the loop), restores the value in index register Y, and returns to the calling routine SCAND1. This routine refetches the same byte, masks out the most significant 4 bits to isolate the *other* character in the byte and calls CONVD again to display the second character. Upon return, index Y is decremented, and, if it is not zero, (zero signifies that all characters have been displayed) routine SCAND1 is reentered to display the next two characters.

When all six characters have been displayed, signified by Y having been decremented to 0, the display is turned off by forcing SBD bits 1 through 4 to 0, so that no display element is selected. Peripheral port A is next reconfigured as an input and a jump to AK (get any key) is performed, initiating a scan of the keyboard.

### B.4.3 Operation of the KIM Keyboard

The keyboard mounted on the KIM circuit board, like the display, is in a very real sense nothing more than a peripheral unit which interacts with the CPU via the 6530-002 peripheral interface unit. The function of the various keys in providing system control will not be discussed here because Chapter 8 covers this in great detail. However, a somewhat detailed examination of the control of this keyboard will be made, to promote understanding of the more general problem of peripheral control, in this case, the control of an input device.

Figure B.3 is a simplified logic diagram which illustrates the principal components which are involved in KIM keyboard interaction. The 74145 decoder is the same one previously discussed relative to display management. In that case the active decoder outputs were those labeled 4 through 9, but, as mentioned there, outputs 0 through 3 are the important ones in connection with keyboard interaction. Actually, output 3 is reserved for use as an indicator that the Teletype (rather than the KIM keyboard/display) is the selected controlling device, see Section B.6. Outputs 0 through 2 are connected to rows 0 through 2 of the KIM keyboard, as shown in Figure B.3. Basic keyboard operation consists of sequentially selecting rows 0, 1, and 2 of the keyboard via an operating program. Any depressed key in a selected row will couple this active-low selection signal (through the vertical keyboard bus line connected to that key) back into one of the port A inputs (assuming that port A is programmed to serve as an input). After the sequential scan of the three rows is complete, the information gathered from the port by the CPU



**Figure B.3** Simplified block diagram of KIM-1 keyboard scan system

during the scan can be interpreted and the depressed key identified. The CPU can then branch to a routine designed to service that key. Note that the programmer has complete liberty in assigning each key its role in the system, since the software he provides to service each key-stroke is whatever he chooses it to be.

#### B.4.4 AK (Get Any Key) Keyboard Service Routine

To examine one implementation of a keyboard scan routine, we will examine the principal steps found in the KIM keyboard service routine. Recall that, upon completion of the display management routine (SCANDS) the KIM monitor software, after reconfiguring the port A I/O interface as an input, performs a subroutine jump to a routine called AK (any key), which may be found in the KIM monitor listing beginning on line 1037.

The first pair of steps in AK set index register Y and index register X values to 03 and 01, respectively. The Y value is used to count through the 3 rows of the display and the X value is used, via the BCD decoder and port B output, to sequentially select the three rows. The accumulator is loaded with all 1s (FF hex), and row 0 selected by storing the X value in SBD. The X value is then incremented by 2 to prepare for selection of row 2 during the next pass. Any key depressed on row 0 will send a low-level signal (a "0") back to the corresponding port A input line. Thus, when the next monitor instruction forms the logical AND of SAD contents with the accumulator, the corresponding accumulator bit will become 0. After this operation, Y is decremented and tested for zero, with a branch back to the scan routine if Y is not zero. Thus three such scans, one for each row will be performed. Any key depression will change the original FF value in the accumulator to some other value.

The next three program steps are included to set up the Teletype selection for potential future use and to take care of the possibility that the Teletype, even though not being used, may have transferred data to the KIM.

A final EOR (exclusive-or) instruction complements the entire accumulator. Thus, if no keys were depressed during the scan, an all-zero accumulator state will be reached. If one or more keys were depressed, the accumulator will contain a nonzero value.

After the AK routine is finished, a RTS (Return from Subroutine) instruction is executed. This transfers control back to the SCAND routine. Reexamination of this routine, however, discloses that the JSR (Jump to Subroutine) to AK was its last step. However, SCAND, as it turns out, is reached via JSR instructions itself, so the actual return will be to the routine which called SCAND in the first place. There are several such calls in the monitor, all on page 17 of the listing. In particular, note the pair of calls on lines 662 and 664. A return from SCAND the first time with a nonzero accumulator value will result in a second call of SCAND, again followed by a test for zero accumulator. Once again the problem of unreliable logical outputs from mechanical switch contacts is being addressed. The KIM CPU refuses to believe the answer given by just one pass of SCAND, so it calls it one more time just to make sure the key is still depressed during the second scan. If it is, the routine designated GETK (line 667) is entered. This routine has as its function the determination of what key was pressed, information AK does not yield.

#### **B.4.5 GETK and GETKEY Monitor Routines**

GETK begins with a jump to subroutine GETKEY, beginning on line 1114 of the monitor listing. GETKEY is a moderately involved routine and its details will not be expounded here. Figure B.4 is a flow chart of GETKEY. Suffice it to say that when it is entered, it scans the KIM keyboard in a manner similar to that done during the AK routine, except that much more processing of the data is performed. The result of this processing is that, when a return from GETKEY occurs, a code is left in the accumulator which uniquely specifies the key that was depressed. Table B.3 shows the association between the keyboard keys and the codes produced by GETKEY. Note that a code with a hex value greater than 15 represents a so-called "illegal" code, and is an indicator that the key was not steadily depressed during the execution of the scan, or that two keys or more were depressed. Such an illegality is rare, because it is extremely difficult to create a keystroke so short as to not appear rather slow to the KIM monitor program. Also, two keys would have to be depressed with near-perfect simultaneity to be detected as an illegal activity by the GETKEY routine.

Returning to the flow of events engendered by the original keystroke, return from GETKEY is back to GETK. The key code having been placed in A, the next few steps of GETK are simply comparisons with code values corresponding to NO KEY, PC, ADDRESS MODE, DATA MODE, STEP (+), and GO key depressions. In each of these cases a jump to a subroutine for handling that particular activity is performed, after which a return to normal keyboard scan will be made. If none of the above codes are found, the code must be a data value held in the lower 4 bit positions of A. In this case, two principal activities must be performed by routine DATA.

First, a determination as to whether the keystroke was intended to affect Data or Address portions of the display is made. The routine accomplishes this by examining the MODE value (held at the reserved address 00FF). If ADDRESS mode is selected, the routine loads the display buffer area with the proper characters, shifts the characters around in such a way as to reflect a right-to-left insertion of new characters to the display, and fetches the data at the resulting memory address. This data is loaded into the display buffer area so that when a return is made to the SCAND routine, the display will indicate the new address and the data stored there. If DATA mode is selected, the address presently held in the display buffer will be used to read memory, shift the low-order data

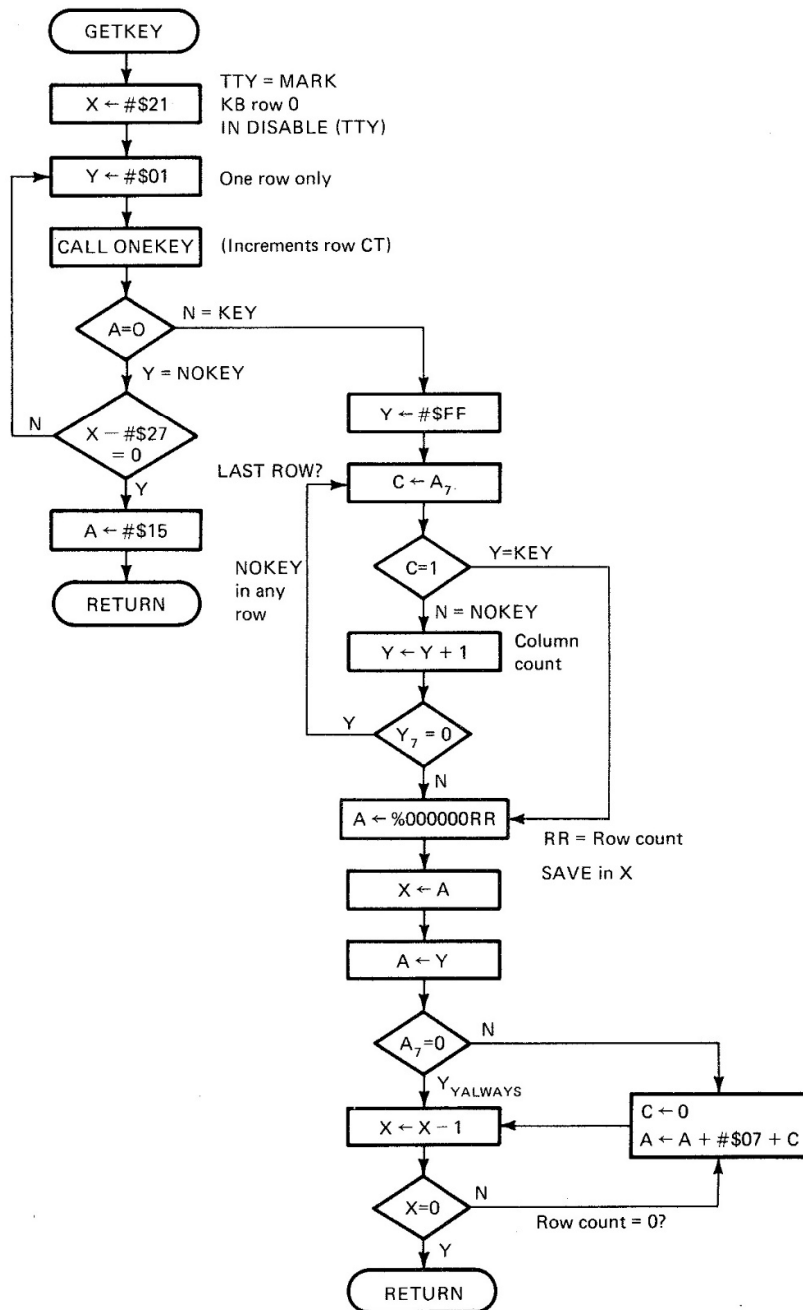


Figure B.4 Flow chart - GETKEY monitor routine



TABLE B.3 KIM Monitor Keystroke Codes

<i>Key Depressed</i>	<i>Code Placed in A by GETKEY (hex)</i>
NONE	15
PC	14
AD	10
DA	11
+	12
GO	13
1	01
2	02
3	03
4	04
5	05
6	06
7	07
8	08
9	09
0	00
A	A0
B	0B
C	0C
D	0D
E	0E
F	0F
Illegal	> 15

character left one hex position and insert the new value (obtained by GETKEY) in the vacated position. The modified value is then rewritten into memory and a return to SCAND (via START, line 636) is performed.

In reading the KIM monitor listing, it is easy to become confused by some program activities which appear to be meaningless or at least obscure. This is a common reaction when trying to obtain an understanding of another programmer's thought processes via a listing of his program. Novice programmers should remember this when writing their own programs, and should try to follow the following two rules:

1. Be very liberal in your use of comments in your programs.

2. Try to avoid being too clever. It may cost a little extra code to do things in a straightforward manner as opposed to developing short but intricate routines. Be forewarned that such “cleverness” will not only be confusing to another worker trying to understand your program listing, but to yourself when you are trying to use or debug the routine a couple of days after having written it.

## B.5 USE OF KIM SUBROUTINES

The discussions of the previous sections involved examinations of some of the activity managed by the KIM monitor. In looking over this operation, you may have noticed that much use is made in the KIM monitor of subroutines performing desired operations. This rather fortunate implementation, one which is a good model for programmers in general, makes available to the KIM user a variety of routines for use outside the normal monitor usage. Table B.4 is a listing of those subroutines included in the portion of the KIM monitor provided for management of display, keyboard, and Teletype. The user can call these routines by simply using the JSR instruction and providing the starting location of the desired routine. They may be thought of as a rather limited “library” available to the system user. The only difference between these subroutines and those you might write for yourself lies in their behavior during the single-step operation of the KIM. Since they all lie in page 7 of memory, you cannot step through them. This does *not* mean that you cannot use them, however, in programs you might like to step through. The only effect will be that in stepping through a JSR instruction which calls one of these routines, you will find that the entire subroutine will be executed before the KIM stops and waits for the next instruction in sequence, the one following the JSR. Note that, in such a case, you are using the subroutine for your own purposes, and the KIM monitor is using the same subroutine, without any conflict between the two uses.

This concept of use of a subroutine library is a very important one. You will find in your future use of microcomputers that the establishment of such a library for use by yourself or others will greatly enhance your ability to use your system without the necessity to develop all software “from the ground up” each time a new and different problem or application is to be examined. You will also find the use of subroutines developed by other users to be an important aid.

TABLE B.4 KIM Monitor Subroutines

<i>Subroutine Name</i>	<i>Starting Address</i>	<i>Description</i>	<i>Comments</i>
INITS	1E88	Initializes system, performing the following steps: 1. Set keyboard to address mode. 2. Program 6530 A port terminals as inputs. 3. Program the lower 6 6530 B port terminals as outputs, and the upper bit as an input (for TTY mode). 4. Loads 03 into B port data register (setting up test for TTY mode). 5. Clear decimal mode of KIM arithmetic unit.	
INIT1	1E8C	Same as INITS, except does not affect keyboard mode.	
SCAND	1F19	Drives KIM display, displaying 4-character address and data stored in memory at that address. Branches to AK routine at completion to test for occurrence of KIM keyboard depression. Returns to calling program after AK.	Address to be displayed stored at POINTL, POINTH.
SCANDS	1F1F	Identical to SCAND, except simply displays data at POINTL, POINTH, INH.	6 hex digits to be displayed must be stored at POINTL, POINTH, AND INH.
AK	1EFE	Scans KIM keyboard to see if any key is depressed. Loads A with 0 if no key, with some nonzero number if key.	PA port must be programmed as input.
GETKEY	1F6A	Scans KIM keyboard, producing code (in A) identifying key depressed (see Table B.4).	PA port must be programmed as input.
INCPT	1F63	Increments the 16-bit value in POINTL, POINTH.	

TABLE B.4 (Continued)

<i>Subroutine Name</i>	<i>Starting Address</i>	<i>Description</i>	<i>Comments</i>
OPEN	1FCC	Moves data from INL to POINTL and from INH to POINTH.	
CONVD	1F48	<ol style="list-style-type: none"> <li>1. Converts hex digit in A to code suitable for driving KIM display element.</li> <li>2. Selects display element with contents of X (see section on driving KIM display).</li> <li>3. Drive display element for about 0.5 milliseconds.</li> <li>4. Increments X by 2, points to next rightmost display element.</li> </ol>	<p>Hex digit to be displayed in lower half of A, with upper half cleared. X contains code selecting display element. X will be modified (incremented by 2) by the routine.</p>
GETCH	1E5A	Get one character from Teletype, placing it in A.	X unchanged, Y=FF on return.
GETBYT	1F9D	Get two hex characters from Teletype (as one byte) and pack into INH and INL (See PACK).	X unchanged, Y=0 on return.
PACK	1FAC	Shift character in A into the INH, INL string. Loses the high-order character of INH. If A does not contain a hex digit, return with no execution.	Clears A on return (if transfer successful).
CHK	1F91	Computes check sum (used with paper tape routine).	
PRTST	1FD5	Prints (via Teletype) a string of ASCII characters from TOP + X to TOP (TOP is a table prestored in ROM).	X must contain pointer into table.
CRLF	1E2F	Prints (via Teletype) a carriage return and line feed.	
OUTCH	1EA0	Prints (via Teletype) one character (in A).	A contains character to be printed, X unchanged, Y=FF on return.

OUTSP	1E9E	Prints (via Teletype) a space. X unchanged, Y=FF on return.
PRTPNT	1E1E	Prints (via Teletype) the contents of POINTL and POINTH as 4 hex characters.
PRTBYT	1E3B	Prints (via Teletype) a byte (in A contains byte at start A) as two hex characters. and at return.
HEXTA	1E4C	Extracts lower 4 bits of A, prints (via Teletype) as hex character.
DELAY	1ED4	Delay one Teletype bit time. Delay depends on baud rate constant (CNTL30, CNTLH30) stored during initialization.
DEHALF	1EEB	Like DELAY, except delays one-half bit time.

---

## B.6 KIM KEYBOARD/TELETYPE SELECTION CONTROL

The KIM system permits the user to select as his primary I/O device either the integral keyboard/display or a console device such as an ASR33 Teletype. It is instructive to examine the means by which the KIM system goes about implementing the choice of option for I/O.

If one reads only the bare operating procedures for the KIM, he finds that the steps involved in setting up the TTY option include wiring of a connector jumper or a switch which is used to select the option. This may suggest that some wiring of hardware is being accomplished, but in fact the selection is performed almost entirely in software. Refer to Figure B.5 which shows a logic drawing of the 6530-002 peripheral interface unit and the switch wired in to permit implementation of the TTY option. Note that output number 03 is tied through this switch to PA0 of the 6530. PA0 should be configured as an input. During the initialization program executed during a RESET operation (a small portion of which is shown in both listing and flow chart form in Figure 4.7), the Port B data register (SBD) is loaded with 07, producing an active-low output at the 03 terminal of the BCD decoder. If the TTY mode switch is closed, this signal will be coupled into the PA0 input. When the program shown executes a BIT (bit test) instruction, (comparing the contents of PA to 01) and then follows this with a BNE (branch on nonzero), the branch will not be taken, since the zero

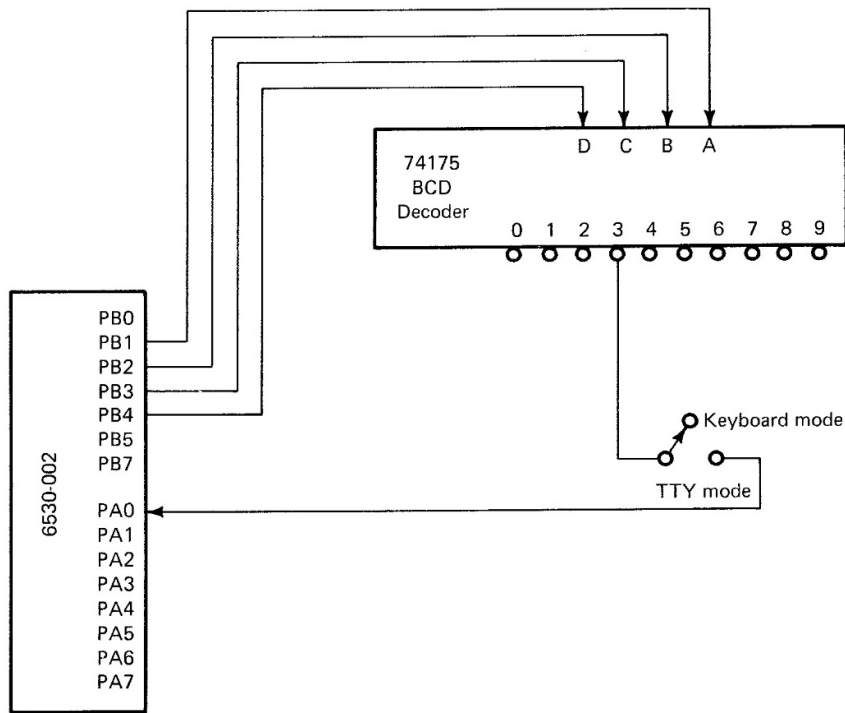


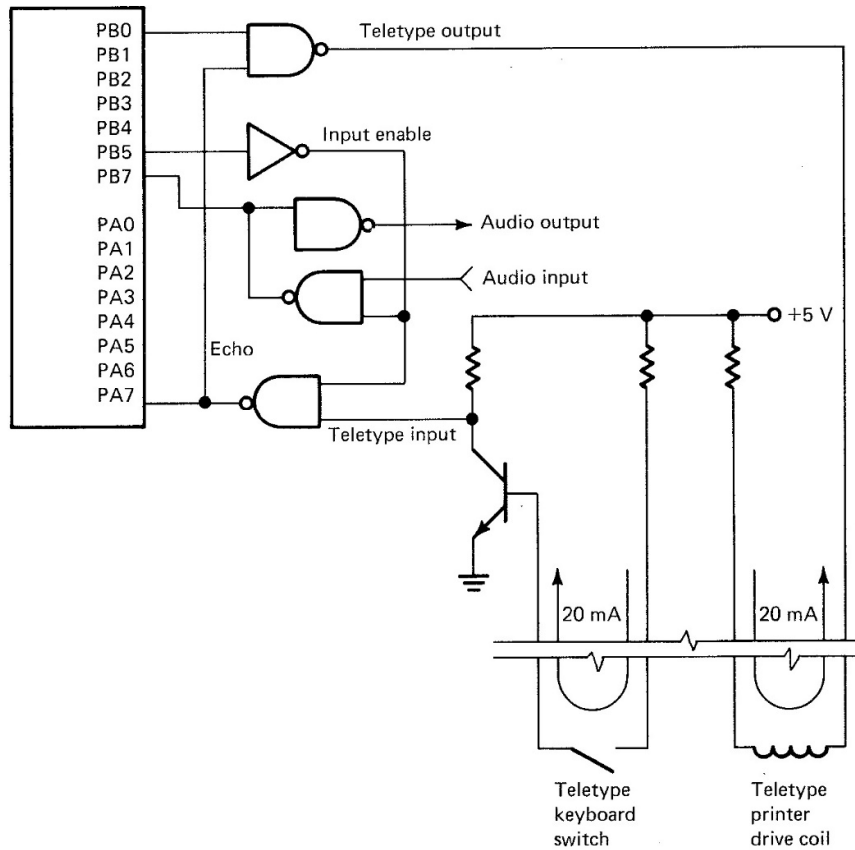
Figure B.5 KIM-1 keyboard/Teletype mode selection circuitry

flag will have been cleared by the bit test operation. Thus the program sequence will not be broken, and the monitor will enter a Teletype handling software mode. If, on the other hand, the BIT, BNE sequence resulted in a zero flag set, indicating that PA0 was 1, a branch to monitor routine START would result, corresponding to keyboard/display selection. (Note that this means that an unconnected input to PA0 is interpreted as a high level or a "1").

There are several branch points in the monitor program which require testing for TTY mode. In each case, a procedure similar to the one described is executed by the software. (See, for example, lines 614, 637, and 659 of the KIM monitor listing.)

## B.7 TELETYPE I/O INITIALIZATION

Figure B.6 is a diagram which shows the interface circuitry required to properly connect the 6530-002 to a 20-milliampere current-loop



**Figure B.6** Teletype current-loop interface circuitry

interface such as is found in a standard ASR33 Teletype. (Also shown is a portion of the audio cassette interface, which will not be discussed here.)

Chapter 5 discusses asynchronous communication principles in some detail, so it is assumed that basic Teletype operation is understood. From the diagram it can be seen that 6530 initialization to support this interface requires that:

1. PB0 be programmed as an output.
2. A high or "1" value be sent to PB0.
3. PA7 be programmed as an input.
4. PB5 be programmed as an output.

5. PB5 be programmed as an output.
6. A low or "0" value be sent to PB5.

The KIM-1 monitor established the 6530-002 interface during an initialization sequence called INITS (lines 966-977). The reader should verify that the steps listed here are performed during this routine.

## B.8 KIM-1 MONITOR PROGRAM LISTING

This section consists of a reproduction of the KIM-1 Monitor listing, included with permission of MOS Technology, Inc.

CARD #	LOC	CODE	CARD
3			666666 555555 333333 000000
4			6 5 3 0 0
5			6 5 3 0 0
6			666666 555555 333333 0 0 0
7			6 6 5 3 0 0
8			6 6 5 3 0 0
9			666666 555555 333333 000000
10			
11			
12			
13			000000 000000 333333
14			0 0 0 0 3
15			----- 0 0 0 0 3
16			----- 0 0 0 0 333333
17			----- 0 0 0 0 3
18			0 0 0 0 3
19			000000 000000 333333
20			
21			
22			
23			
24			
25			COPYRIGHT
26			MOS TECHNOLOGY, INC
27			DATE OCT 18 1975 REV D
28			
29			
30			
31			6530-003 IS AN AUDIO CASSETT TAPE
32			RECORDER EXTENSION OF THE BASIC
33			KIM MONITOR
34			
35			IT FEATURES TWO BASIC ROUTINES
36			LOADT-LOAD MEM FROM AUDIO TAPE
37			DUMPT-STOR MEM ONTO AUDIO TAPE
38			
39			LOADT
40			ID=00 IGNORE ID
41			ID=FF IGN. ID USE SA FOR START ADDR
42			ID=01-FE IGN.ID USE ADDR ON TAPE
43			
44			DUMPT
45			ID=00 SHOULD NOT BE USED
46			ID=FF SHOULD NOT BE USED
47			ID=01-FE NORMAL ID RANGE
48			SAL LSB STARTING ADDRESS
49			SAH MSB
50			EAL LSB ENDING ADDRESS
51			EAH MSB
52			



CARD #	LOC	CODE	CARD
54			;
55			;
56			EQUATES
57			;
58			SET UP FOR 6530-002 I/O
59			;
60			SAD     =\$1740             6530 A DATA
61			PADD   =\$1741             6530 A DATA DIRECTION
62			SBD     =\$1742             6530 B DATA
63			PBDD    =\$1743             6530 B DATA DIRECTION
64			CLK1T   =\$1744             DIV BY 1 TIME
65			CLK8T   =\$1745             DIV BY 8 TIME
66			CLK64T  =\$1746             DIV BY 64 TIME
67			CLKKT   =\$1747             DIV BY 1024 TIME
68			CLKRDI   =\$1747            READ TIME OUT BIT
69			CLKRDT   =\$1746            READ TIME
70			;
71			;
72	000F		++=\$00EF
73			MPU REG. SAVX AREA IN PAGE 0
74			;
75	00F0		PCL     ++=+1 PROGRAM CNT LOW
76	00F1		PCH     ++=+1 PROGRAM CNT HI
77	00F2		PREG    ++=+1 CURRENT STATUS REG.
78	00F3		SPUSER  ++=+1 CURRENT STACK POINT
79	00F4		ACC     ++=+1 ACCUMULATOR
80	00F5		XREG    ++=+1 X INDEX
81			YREG    ++=+1 Y INDEX
82			;
83			KIM FIXED AREA IN PAGE 0
84			;
85	00F6		CHKHI   ++=+1
86	00F7		CHKSUM  ++=+1
87	00F8		INL     ++=+1 INPUT BUFFER
88	00F9		INH     ++=+1 INPUT BUFFER
89	00FA		POINTL  ++=+1 LSB OF OPEN CELL
90	00FB		POINTH  ++=+1 MSB OF OPEN CELL
91	00FC		TEMP    ++=+1
92	00FD		TMPX    ++=+1
93	00FE		CHAR    ++=+1
94	00FF		MODE    ++=+1
95			;
96			KIM FIXED AREA IN PAGE 23
97			;
98			++=\$17E7
99	17E7		CHKL    ++=+1
100	17E8		CHKH    ++=+1             CHKSUM
101	17E9		SAVX    ++=+3
102	17EA		VEB     ++=+6             VOLATILE EXECUTION BLOCK
103	17EB		CNTL30  ++=+1             TTY DELAY
104	17EC		CNTH30  ++=+1             TTY DELAY
105	17ED		TIMH    ++=+1
106	17EE		SAL     ++=+1             LOW STARTING ADDRESS
107	17EF		SAH     ++=+1             HI STARTING ADDRESS
108	17F0		EAL     ++=+1             LOW ENDING ADDRESS
109	17F1		EAH     ++=+1             HI ENDING ADDRESS
110	17F2		ID     ++=+1
111			;
112			INTERRUPT VECTORS
113			;
114	17FA		NMIV    ++=+2             STOP VECTOR (STOP=1C00)
115	17FB		RSTV    ++=+2             RST VECTOR
116	17FC		IRQV    ++=+2             IRQ VECTOR (BRK= 1C00)
117			;
118	1800		++=\$1800
119			;
120			INIT VOLATILE EXECUTION BLOCK
121			DUMP MEM TO TAPE
122			;
123	1800	A9 AD	DUMPT LDA    =\$AD         LOAD ABSOLUTE INST

CARD #	LOC	CODE	CARD		
122	1802	8D EC 17		STA	VEB
123	1805	20 32 19		JSR	INTVEB
124					
125	1808	A9 27		LDA	#\$27
126	180A	8D 42 17		STA	SBD
127	180D	A9 BF		LDA	#\$BF
128	180F	8D 43 17		STA	PBDD
129					
130	1812	A2 64		LDX	#\$64
131	1814	A9 16	DUMPT1	LDA	#\$16
132	1816	20 7A 19		JSR	OUTCHT
133	1819	CA		DEX	
134	181A	D0 F8		BNE	DUMPT1
135					
136					
137	181C	A9 2A		LDA	#\$*
138	181E	20 7A 19		JSR	OUTCHT
139					
140	1821	AD F9 17		LDA	ID
141	1824	20 61 19		JSR	OUTBT
142					
143	1827	AD F5 17		LDA	SAL
144	182A	20 5E 19		JSR	OUTBTC
145	182D	AD F6 17		LDA	SAH
146	1830	20 5E 19		JSR	OUTBTC
147					
148	1833	AD ED 17	DUMPT2	LDA	VEB+1
149	1836	CD F7 17		CMP	EAL
150	1839	AD EE 17		LDA	VEB+2
151	183C	ED F8 17		SBC	EAH
152	183F	90 24		BCC	DUMPT4
153					
154	1841	A9 2F		LDA	#\$//
155	1843	20 7A 19		JSR	OUTCHT
156	1846	AD E7 17		LDA	CHKL
157	1849	20 61 19		JSR	OUTBT
158	184C	AD E8 17		LDA	CHKH
159	184F	20 61 19		JSR	OUTBT
160					
161					
162	1852	A2 02		LDX	#\$02
163	1854	A9 04	DUMPT3	LDA	#\$04
164	1856	20 7A 19		JSR	OUTCHT
165	1859	CA		DEX	
166	185A	D0 F8		BNE	DUMPT3
167					
168	185C	A9 00		LDA	#\$00
169	185E	85 FA		STA	POINTL
170	1860	85 FB		STA	POINTH
171	1862	4C 4F 1C		JMP	START
172					
173	1865	20 EC 17	DUMPT4	JSR	VEB
174	1868	20 5E 19		JSR	OUTBTC
175					
176	186B	20 EA 19		JSR	INCVEB
177	186E	4C 33 18		JMP	DUMPT2
178					
179					
180					
181					
182	1871	0F 19	TAB	WORD	LOAD12
183	1873	A9 8D	LOADT	LDA	#\$8D
184	1875	8D EC 17		STA	VEB
185	1878	20 32 19		JSR	INTVEB
186					

TURN OFF DATAIN PB5

CONVERT PB7 TO OUTPUT

100 CHARS  
SYN CHAR'S

START CHAR

OUTPUT ID

OUTPUT STARTING  
ADDRESSCHECK FOR LAST  
DATA BYTE

OUTPUT END OF DATA CHR

LAST BYTE HAS BEEN  
OUT PUT NOW OUTPUT  
CHKSUM2 CHAR'S  
EOT CHARDISPLAY 0000  
FOR NORMAL EXIT

DATA BYTE OUTPUT

LOAD MEMORY FROM TAPE

INIT VOLATILE EXECUTION  
BLOCK WITH STA ABS.

CARD #	LOC	CODE	CARD		
187	187B	A9 4C	LDA	#\$4C	JUMP TYPE RTRN
188	187D	8D EF 17	STA	VEB+3	
189	1880	AD 71 18	LDA	TAB	
190	1883	8D F0 17	STA	VEB+4	
191	1886	AD 72 18	LDA	TAB+1	
192	1889	8D F1 17	STA	VEB+5	
193					
194	188C	A9 07	LDA	#\$07	RESET P85=0 (DATA IN)
195	188E	8D 42 17	STA	SBD	
196					
197	1891	A9 FF	SYNC LDA	#\$FF	CLEAR SAVX FOR SYNC AREA
198	1893	8D E9 17	STA	SAVX	
199					
200	1896	20 41 1A	SYNC1 JSR	RDBIT	GET A BIT
201	1899	4E E9 17	LSR	SAVX	SHIFT BIT INTO CHAR
202	189C	0D E9 17	ORA	SAVX	
203	189F	8D E9 17	STA	SAVX	
204	18A2	AD E9 17	LDA	SAVX	GET NEW CHAR
205	18A5	C9 16	CMP	#\$16	SYN CHAR
206	18A7	D0 ED	BNE	SYNC1	
207					
208	18A9	A2 0A	LDX	#\$0A	TEST FOR 10 SYN CHARS
209	18AB	20 24 1A	SYNC2 JSR	RDCHT	
210	18AE	C9 16	CMP	#\$16	
211	18B0	D0 DF	BNE	SYNC	IF NOT 10 CHAR RE-SYNC
212	18B2	CA	DEX		
213	18B3	D0 F6	BNE	SYNC2	
214					
215					
216	18B5	20 24 1A	LOADT4 JSR	RDCHT	LOOK FOR START OF
217	18B8	C9 2A	CMP	#'/	DATA CHAR
218	18BA	F0 06	BEQ	LOAD11	
219	18BC	C9 16	CMP	#\$16	IF NOT * SHOULD BE SYN
220	18BE	D0 D1	BNE	SYNC	
221	18C0	F0 F3	BEQ	LOADT4	
222					
223	18C2	20 F3 19	LOAD11 JSR	RDBYT	READ ID FROM TAPE
224	18C5	CD F9 17	CMP	ID	COMPARE WITH REQUESTED ID
225	18C8	F0 0D	BEQ	LOADT5	
226	18CA	AD F9 17	LDA	ID	DEFAULT 00 READ RECORD
227	18CD	C9 00	CMP	#\$00	ANYWAY
228	18CF	F0 06	BEQ	LOADT5	
229	18D1	C9 FF	CMP	#\$FF	DEFAULT FF IGNOR SA ON
230	18D3	F0 17	BEQ	LOADT6	TAPE
231	18D5	D0 9C	BNE	LOADT	
232					
233	18D7	20 F3 19	LOADT5 JSR	RDBYT	GET SA FROM TAPE
234	18DA	20 4C 19	JSR	CHKT	
235	18DD	8D ED 17	STA	VEB+1	SAVX IN VEB+1,2
236	18E0	20 F3 19	JSR	RDBYT	
237	18E3	20 4C 19	JSR	CHKT	
238	18E6	8D EE 17	STA	VEB+2	
239	18E9	4C F8 18	JMP	LOADT7	
240					
241	18EC	20 F3 19	LOADT6 JSR	RDBYT	GET SA BUT IGNORE
242	18EF	20 4C 19	JSR	CHKT	
243	18F2	20 F3 19	JSR	RDBYT	
244	18F5	20 4C 19	JSR	CHKT	
245					
246					
247	18F8	A2 02	LOADT7 LDX	#\$02	GET 2 CHARS
248	18FA	20 24 1A	LOAD13 JSR	RDCHT	GET CHAR(X)
249	18FD	C9 2F	CMP	#'/	LOOK FOR LAST CHAR
250	18FF	F0 14	BEQ	LOADT8	
251	1901	20 00 1A	JSR	PACKT	CONVERT TO HEX

CARD #	LOC	CODE	CARD		
252	1904	D0 23	BNE	LOADT9	Y=1 NON-HEX CHAR
253	1906	CA	DEX		
254	1907	D0 F1	BNE	LOAD13	
255					
256	1909	20 4C 19	JSR	CHKT	COMPUTE CHECKSUM
257	190C	4C EC 17	JMP	VEB	SAVX DATA IN MEMORY
258	190F	20 FA 19	LOAD12 JSR	INCVEB	INCREMENT DATA POINTER
259	1912	4C F8 18	JMP	LOADT7	
260					
261	1915	20 F3 19	LOADT8 JSR	RDBYT	END OF DATA COMPARE CHKSUM
262	1918	CD E7 17	CMP	CHKL	
263	191B	D0 0C	BNE	LOADT9	
264	191D	20 F3 19	JSR	RDBYT	
265	1920	CD E8 17	CMP	CHKH	
266	1923	D0 04	BNE	LOADT9	
267	1925	A9 00	LDA	#00	NORMAL EXIT
268	1927	F0 02	BEQ	LOAD10	
269					
270	1929	A9 FF	LOADT9 LDA	#FF	ERROR EXIT
271	192B	85 FA	LOAD10 STA	POINTL	
272	192D	85 FB	STA	POINTH	
273	192F	4C 4F 1C	JMP	START	
274					
276					
277				SUBROUTINES FOLLOW	
278					
279				SUB TO MOVE SA TO VEB+1+2	
280					
281	1932	AD F5 17	INTVEB LDA	SAL	
282	1935	8D ED 17	STA	VEB+1	
283	1938	AD F6 17	LDA	SAH	
284	193B	8D EE 17	STA	VEB+2	
285	193E	A9 60	LDA	#60	RTS INST
286	1940	8D EF 17	STA	VEB+3	
287	1943	A9 00	LDA	#00	CLEAR CHKSUM AREA
288	1945	8D E7 17	STA	CHKL	
289	1948	8D E8 17	STA	CHKH	
290	194B	60	RTS		
291					
292				COMPUTE CHKSUM FOR TAPE LOAD	
293				RTN USES Y TO SAVX A	
294					
295	194C	A8	CHKT	TAY	
296	194D	18	CLC		
297	194E	6D E7 17	ADC	CHKL	
298	1951	8D E7 17	STA	CHKL	
299	1954	AD E8 17	LDA	CHKH	
300	1957	69 00	ADC	#00	
301	1959	8D E8 17	STA	CHKH	
302	195C	98	TYA		
303	195D	60	RTS		
304					
305				OUTPUT ONE BYTE USE Y	
306				TO SAVX BYTE	
307					
308	195E	20 4C 19	OUTBTC JSR	CHKT	COMP CHKSUM
309	1961	A8	OUTBT TAY		SAVX DATA BYTE
310	1962	4A	LSR	A	SHIFT OFF LSD
311	1963	4A	LSR	A	
312	1964	4A	LSR	A	
313	1965	4A	LSR	A	
314	1966	20 6F 19	JSR	HEXOUT	OUT PUT MSD
315	1969	98	TYA		
316	196A	20 6F 19	JSR	HEXOUT	OUT PUT LSD
317	196D	98	TYA		

```

CARD # LOC      CODE      CARD
318 196E 60      RTS
319
320
321      ;      CONVERT LSD OF A TO ASCII
322      ;      AND OUTPUT TO TAPE
323 196F 29 0F    HEXOUT AND  #$0F
324 1971 09 0A    CMP  #$0A
325 1973 18      CLC
326 1974 30 02    BMI  HEX1
327 1976 69 07    ADC  #$07
328 1978 69 30    HEX1  ADC  #$30
329
330      ;      OUTPUT TO TAPE ONE ASCII
331      ;      CHAR USE SUB'S ONE + ZRO
332
333 197A 8E E9 17  OUTCHT STX  SAVX
334 197D 8C EA 17    STY  SAVX+1
335 1980 A0 08      LDY  #$08      START BIT
336 1982 20 9E 19  CHT1  JSR  ONE      GET DATA BIT
337 1985 4A        LSR  A
338 1986 B0 06      BCS  CHT2
339 1988 20 9E 19  JSR  ONE      DATA BIT=1
340 198B 4C 91 19  JMP  CHT3
341 198E 20 C4 19  CHT2  JSR  ZRO      DATA BIT=0
342 1991 20 C4 19  CHT3  JSR  ZRO
343 1994 88      DEY
344 1995 D0 EB      BNE  CHT1
345 1997 AE E9 17    LDY  SAVX
346 199A AC EA 17    LDY  SAVX+1
347 199D 60      RTS
348
349
350      ;      OUTPUT 1 TO TAPE
351      ;      9 PULSES 138 MICROSEC EACH
352
353 199E A2 09      ONE   LDX  #$09
354 19A0 48      PHA
355 19A1 2C 47 17  ONE1  BIT  CLKRDI      SAVX A
356 19A4 10 FB      BPL  ONE1      WAIT FOR TIME OUT
357 19A6 A9 7E      LDA  #126
358 19A8 8D 44 17  STA  CLK1T
359 19AB A9 47      LDA  #$A7
360 19AD 8D 42 17  STA  SBD      SET  PB7=1
361 19B0 2C 47 17  ONE2  BIT  CLKRDI
362 19B3 10 FB      BPL  ONE2
363 19B5 A9 7E      LDA  #126
364 19B7 8D 44 17  STA  CLK1T
365 19BA A9 27      LDA  #$27
366 19BC 8D 42 17  STA  SBD      RESET PB7=0
367 19BF CA      DEX
368 19C0 D0 DF      BNE  ONE1
369 19C2 68      PLA
370 19C3 60      RTS
371
372
373      ;      OUTPUT 0 TO TAPE
374      ;      6 PULSES 207 MICROSEC EACH
375
376 19C4 A2 06      ZRO   LDX  #$06
377 19C6 48      PHA
378 19C7 2C 47 17  ZRO1  BIT  CLKRDI      SAVX A
379 19CA 10 FB      BPL  ZRO1
380 19CC A9 03      LDA  #195
381 19CE 8D 44 17  STA  CLK1T
382 19D1 A9 47      LDA  #$A7

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```

CARD # LOC      CODE      CARD
383 1903 80 42 17      STA  SBD          SET P87=1
384 1906 80 47 17 2902 BIT  CLKRDI
385 1909 10 FB          BPL  2902
386 1908 A9 03          LDA  #195
387 1900 80 44 17      STA  CLK1T
388 1950 A9 27          LDA  #27
389 19E2 80 42 17      STA  SBD          RESET P87=0
390 19E5 0A            DEX
391 19E6 00 DF          BNE  2901
392 19E8 68            PLA          RESTORE A
393 19E9 60            RTS
394
395 ;
396 ;      SUB TO INC WEB+1,2
397 ;
397 19EA EE ED 17  INCVB INC  WEB+1
398 19ED 00 03          BNE  INCVE1
399 19EF EE EE 17      INC  WEB+2
400 19F2 60            INCVB1 RTS
401 ;
402 ;      SUB TO READ BYTE FROM TAPE
403 ;
404 19F3 20 24 1A  RDBYT JSR  RDCHT
405 19F6 20 00 1A          JSR  PACKT
406 19F9 20 24 1A  RDBYT2 JSR  RDCHT
407 19FC 20 00 1A          JSR  PACKT
408 19FF 60            RTS
409 ;
410 ;      PACK A=ASCII INTO SAVX
411 ;      AS HEX DATA
412 ;
413 1A00 09 30      PACKT CMP  #30
414 1A02 30 1E          BMI  PACKT3
415 1A04 09 47      CMP  #47
416 1A06 10 1A      BPL  PACKT3
417 1A08 09 40      CMP  #40
418 1A0A 30 03          BMI  PACKT1
419 1A0C 18          CLC
420 1A0D 69 09      ADC  #09
421 1A0F 2A          PACKT1 ROL  A
422 1A10 2A          ROL  A
423 1A11 2A          ROL  A
424 1A12 2A          ROL  A
425 1A13 A0 04      LDY  #04
426 1A15 2A          PACKT2 ROL  A
427 1A16 2E EA 17      ROL  SAVX
428 1A19 98          DEY
429 1A1A 00 F9      BNE  PACKT2
430 1A1C A0 EA 17      LDA  SAVX
431 1A1F A0 00      LDY  #00
432 1A21 60            RTS          Y=0 VALID HEX CHAR
433 1A22 08          PACKT3 INY          Y=0 VALID HEX
434 1A23 60            RTS          Y=1 NOT HEX
435 ;
436 ;      GET 1 CHAR FROM TAPE AND RETURN
437 ;      WITH CHAR IN A  USE SAVX+1 TO ASM CHAR
438 ;
439 1A24 3E EB 17  RDCHT STX  SAVX+2
440 1A27 A2 08      LDX  #08          READ 8 BITS
441 1A29 20 41 1A  RDCHT1 JSR  RDBIT          GET NEXT DATA BIT
442 1A2C 4E EA 17      LSR  SAVX+1          RIGHT SHIFT CHAR
443 1A2F 00 EA 17      ORA  SAVX+1          OR IN SIGN BIT
444 1A32 20 EA 17      STA  SAVX+1          REPLACE CHAR
445 1A35 0A          DEX
446 1A36 00 F1      BNE  RDCHT1
447 ;

```

CARD #	LOC	CODE	CARD			
448	1A38	AD EA 17	LDA	SAVX+1	MOVE CHAR INTO A	
449	1A3E	2A	RDL	A	SHIFT OFF PARITY	
450	1A3C	4A	LSR	A		
451	1A3D	4E FB 17	LDR	SAVX+2		
452	1A40	50	RTS			
453						
454						
455					THIS SUB GETS ONE BIT FROM	
456					TAPE AND RETURNS IT IN SIGN OF A	
457	1A41	2C 42 17	RDBIT	BIT	SBD	WAIT FOR END OF START BIT
458	1A44	10 FB	BPL	RDBIT		
459	1A46	4D 46 17	LDA	CLKRDT	GET START BIT TIME	
460	1A49	A0 FF	LDY	#FF	A=256-T1	
461	1A4B	8C 46 17	STY	CLK64T	SET UP TIMER	
462						
463	1A4E	A0 14	LDY	#14		
464	1A50	98	RDBIT3	DEY	DELAY 100 MICROSEC	
465	1A51	D0 FD	BNE	RDBIT3		
466						
467	1A53	2C 42 17	RDBIT2	BIT	SBD	
468	1A56	30 FB	BMI	RDBIT2	WAIT FOR NEXT START BIT	
469						
470	1A58	38	SEC			
471	1A59	ED 46 17	SBC	CLKRDT	(256-T1)-(256-T2)=T2-T1	
472	1A5C	A0 FF	LDY	#FF		
473	1A5E	8C 46 17	STY	CLK64T	SET UP TIMER FOR NEXT BIT	
474						
475	1A61	A0 07	LDY	#07		
476	1A63	28	RDBIT4	DEY	DELAY 50 MICROSEC	
477	1A64	D0 FD	BNE	RDBIT4		
478						
479	1A66	49 FF	EOR	#FF	COMPLEMENT SIGN OF A	
480	1A68	29 80	AND	#80	MASK ALL EXCEPT SIGN	
481	1A6A	60	RTS			
482						
484					DIAGNOSTICS	
485					MEMORY	
486					PLLCAL	
487						
488						
489						
490					PLLCAL OUTPUT 166 MICROSEC	
491					PULSE STRING	
492						
493	1A6B	A9 27	PLLCAL	LDA	#27	
494	1A6D	8D 42 17	STA	SBD	TURN OFF DATIN PB5=1	
495	1A70	A9 BF	LDA	#BF	CONVERT PB7 TO OUTPUT	
496	1A72	8D 43 17	STA	P8DD		
497						
498	1A75	2C 47 17	PLL1	BIT	CLKRDI	
499	1A78	10 FB	BPL	PLL1		
500	1A7A	A9 9A	LDA	#154	WAIT 166 MICRO SEC	
501	1A7C	8D 44 17	STA	CLK1T		
502	1A7F	A9 A7	LDA	#A7	OUTPUT PB7=1	
503	1A81	8D 42 17	STA	SBD		
504						
505	1A84	2C 47 17	PLL2	BIT	CLKRDI	
506	1A87	10 FB	BPL	PLL2		
507	1A89	A9 9A	LDA	#154		
508	1A8B	8D 44 17	STA	CLK1T		
509	1A8E	A9 27	LDA	#27	PB7=0	
510	1A90	8D 42 17	STA	SBD		
511	1A93	4C 75 1A	JMP	PLL1		
512						
513						

```

CARD # LOC      CODE      CARD
514          ;          INTERRUPTS PAGE 27
515          ;
516 1A96        ;          ***$0164  RESERVED FOR TEST
517 1BFA 6B 1A  NMIP27 .WORD PLLCAL
518 1BFC 6B 1A  RSTP27 .WORD PLLCAL
519 1BFE 6B 1A  IRQP27 .WORD PLLCAL
520          ;
522          ;
523          ;
524          ;
525          ;
526          ;          666666 555555 333333 000000
527          ;          6      5      3      0      0
528          ;          6      5      3      0      0
529          ;          666666 555555 333333 0      0
530          ;          6      6      5      3      0      0
531          ;          6      6      5      3      0      0
532          ;          666666 555555 333333 000000
533          ;
534          ;
535          ;
536          ;          000000 000000 222222
537          ;          0      0      0      0      2
538          ;          ----- 0      0      0      0      2
539          ;          ----- 0      0      0      0      222222
540          ;          ----- 0      0      0      0      2
541          ;          0      0      0      0      2
542          ;          000000 000000 222222
543          ;
545          ;
546          ;
547          ;
548          ;          COPYRIGHT
549          ;          MOS TECHNOLOGY INC.
550          ;          DATE OCT 13 1975  REV E
551          ;
552          ;          KIM :TTY INTERFACE
553          ;          :KEYBOARD INTERFACE
554          ;          :7 SEG 6 DIGIT DISPLAY
555          ;
556          ;
557          ;          TTY CMDS:
558          ;          G  GOEXEC
559          ;          CR OPEN NEXT CELL
560          ;          LF OPEN PREV. CELL
561          ;          .  MODIFY OPEN CELL
562          ;          SP OPEN NEW CELL
563          ;          L  LOAD (OBJECT FORMAT)
564          ;          Q  DUMP FROM OPEN CELL ADDR TO HI LIMIT
565          ;          RD RUB OUT - RETURN TO START (KIM)
566          ;          ((ALL ILLEGAL CHAR ARE IGNORED))
567          ;
568          ;          KEYBOARD CMDS:
569          ;          ADDR SETS MODE TO MODIFY CELL ADDRESS
570          ;          DATA SETS MODE TO MODIFY DATA IN OPEN CELL
571          ;          STEP INCREMENTS TO NEXT CELL
572          ;          RST SYSTEM RESET
573          ;          RUN GOEXEC
574          ;          STOP $1000 CAN BE LOADED INTO NMIV TO
575          ;          USE STOP FEATURE
576          ;          PC DISPLAY PC
577          ;
578          ;          CLOCK IS NOT DISABLED IN SIGMA 1
579          ;
580          ;

```



CARD #	LOC	CODE	CARD		
581			;		
582			;		
584	1C00		♦=\$1C00		
585			;		
586			;		
587	1C00	85 F3	SAVE STA ACC	KIM ENTRY VIA STOP (NMI)	
588	1C02	68	PLA	OR BRK (IRQ)	
589	1C03	85 F1	STA PREG		
590	1C05	68	SAVE1 PLA	KIM ENTRY VIA JSR (A LOST)	
591	1C06	85 EF	STA PCL		
592	1C08	85 FA	STA POINTL		
593	1C0A	68	PLA		
594	1C0B	85 F0	STA PCH		
595	1C0D	85 FB	STA POINTH		
596	1C0F	84 F5	SAVE2 STY YREG		
597	1C11	86 F4	STX XREG		
598	1C13	BA	TSX		
599	1C14	86 F2	STX SPUSER		
600	1C16	20 88 1E	JSR INITS		
601	1C19	4C 4F 1C	JMP START		
602			;		
603	1C1C	6C FA 17	NMIT JMP (NMIV)	NON-MASKABLE INTERRUPT TRAP	
604	1C1F	6C FE 17	IRQT JMP (IRQV)	INTERRUPT TRAP	
605			;		
606	1C22	A2 FF	RST LDX \$FF	KIM ENTRY VIA RST	
607	1C24	9A	TXS		
608	1C25	86 F2	STX SPUSER		
609	1C27	20 88 1E	JSR INITS		
610			;		
611			;		
612	1C2A	A9 FF	DETCPS LDA \$FF	COUNT START BIT	
613	1C2C	8D F3 17	STA CNTH30	ZERO CNTH30	
614	1C2F	A9 01	LDA \$01	MASK HI ORDER BITS	
615	1C31	2C 40 17	DET1 BII SAD	TEST	
616	1C34	D0 19	BNE START	KEYBD SSM TEST	
617	1C36	30 F9	BMI DET1	START BIT TEST	
618	1C38	A9 FC	LDA \$FC		
619	1C3A	13	DET3 CLC	THIS LOOP COUNTS	
620	1C3B	69 01	ADC \$01	THE START BIT TIME	
621	1C3D	30 03	BCC DET2		
622	1C3F	EE F3 17	INC CNTH30		
623	1C42	AC 40 17	DET2 LDY SAD	CHECK FOR END OF START BIT	
624	1C45	10 F3	BPL DET3		
625	1C47	3D F2 17	STA CNTL30		
626	1C4A	A2 08	LDX \$08		
627	1C4C	20 6A 1E	JSR GET5	GET REST OF THE CHAR	
628			;	TEST CHAR HERE	
629			;		
630			;		
631			;		
632			;		
633			;		
634			;		
635			MAKE TTY/KB SELECTION		
636	1C4F	20 8C 1E	START JSR INIT1		
637	1C52	A9 01	LDA \$01		
638	1C54	2C 40 17	BIT SAD		
639	1C57	D0 1E	BNE TTYKB		
640	1C59	20 2F 1E	JSR CRLF	PRT CR LF	
641	1C5C	A2 0A	LDX \$0A	TYPE OUT KIM	
642	1C5E	20 31 1E	JSR PRTST		
643	1C61	4C AF 1D	JMP SHOW1		
644			;		
645	1C64	A9 00	CLEAR LDA \$00		
646	1C66	85 F8	STA INL	CLEAR INPUT BUFFER	

CARD #	LOC	CODE	CARD	INSTR	DATA	COMMENT
647	1068	95 F9		STA	INH	
648	106A	20 5A 1E	READ	JSR	GETCH	GET CHAR
649	106D	C9 01		CMP	#\$01	
650	106F	F0 06		BEQ	TTYKB	
651	1071	20 AC 1F		JSR	PACK	
652	1074	4C DB 1D		JMP	SCAN	
653						
654						
655						MAIN ROUTINE FOR KEY BOARD
656						AND DISPLAY
657	1077	20 19 1F	TTYKB	JSR	SCAND	IF A=0 NO KEY
658	107A	D0 D3		BNE	START	
659	107C	A9 01	TTYKB1	LDA	#\$01	
660	107E	2C 40 17		BIT	SAD	
661	1081	F0 CC		BEQ	START	
662	1083	20 19 1F		JSR	SCAND	
663	1086	F0 F4		BEQ	TTYKB1	
664	1088	20 19 1F		JSR	SCAND	
665	108B	F0 EF		BEQ	TTYKB1	
666						
667	108D	20 6A 1F	GETK	JSR	GETKEY	
668	1090	C9 15		CMP	#\$15	
669	1092	10 BB		BPL	START	
670	1094	C9 14		CMP	#\$14	
671	1096	F0 44		BEQ	PCCMD	DISPLAY PC
672	1098	C9 10		CMP	#\$10	ADDR MODE=1
673	109A	F0 2C		BEQ	ADDRM	
674	109C	C9 11		CMP	#\$11	DATA MODE=1
675	109E	F0 2C		BEQ	DATAM	
676	10A0	C9 12		CMP	#\$12	STEP
677	10A2	F0 2F		BEQ	STEP	
678	10A4	C9 13		CMP	#\$13	RUN
679	10A6	F0 31		BEQ	GOV	
680	10A8	0A	DATA	ASL	A	SHIFT CHAR INTO HIGH
681	10A9	0A		ASL	A	ORDER NIBBLE
682	10AA	0A		ASL	A	
683	10AB	0A		ASL	A	
684	10AC	85 FC		STA	TEMP	STORE IN TEMP
685	10AE	A2 04		LDX	#\$04	
686	10B0	A4 FF	DATA1	LDY	MODE	TEST MODE 1=ADDR
687	10B2	D0 0A		BNE	ADDR	MODE=0 DATA
688	10B4	B1 FA		LDA	(POINTL),Y	GET DATA
689	10B6	06 FC		ASL	TEMP	SHIFT CHAR
690	10B8	2A		ROL	A	SHIFT DATA
691	10B9	91 FA		STA	(POINTL),Y	STORE OUT DATA
692	10BB	4C C3 1C		JMP	DATA2	
693						
694	10BE	0A	ADDR	ASL	A	SHIFT CHAR
695	10BF	26 FA		ROL	POINTL	SHIFT ADDR
696	10C1	26 FB		ROL	POINTH	SHIFT ADDR HI
697	10C3	CA	DATA2	DEX		
698	10C4	D0 EA		BNE	DATA1	DO 4 TIMES
699	10C6	F0 08		BEQ	DATAM2	EXIT HERE
700						
701	10C8	A9 01	ADDRM	LDA	#\$01	
702	10CA	D0 02		BNE	DATAM1	
703						
704	10CC	A9 00	DATAM	LDA	#\$00	
705	10CE	85 FF	DATAM1	STA	MODE	
706	10D0	4C 4F 1C	DATAM2	JMP	START	
707						
708	10D3	20 63 1F	STEP	JSR	INCPT	
709	10D6	4C 4F 1C		JMP	START	
710						
711	10D9	4C C8 1D	GOV	JMP	GOEXEC	

CARD #	LOC	CODE	CARD
712			;
713			;
714			;
715			;
716			;
717	10D0	A5 EF	PCCMD LDA PCL
718	10DE	85 FA	STA POINTL
719	10E0	A5 F0	LDA PCH
720	10E2	85 FB	STA POINTH
721	10E4	4C 4F 1C	JMP START
722			;
723			;
724			;
725	10E7	20 5A 1E	LOAD JSR GETCH LOOK FOR FIRST CHAR
726	10EA	C9 3E	CMP #\$3E SMICOLON
727	10EC	D0 F9	BNE LOAD
728	10EE	A9 00	LOADS LDA #\$00
729	10F0	85 F7	STA CHKSUM
730	10F2	85 F6	STA CHKHI
731			;
732	10F4	20 9D 1F	JSR GETBYT GET BYTE CNT
733	10F7	AA	TAX SAVE IN X INDEX
734	10F8	20 91 1F	JSR CHK COMPUTE CHKSUM
735			;
736	10FB	20 9D 1F	JSR GETBYT GET ADDRESS HI
737	10FE	95 FB	STA POINTH
738	1000	20 91 1F	JSR CHK
739	1003	20 9D 1F	JSR GETBYT GET ADDRESS LO
740	1006	95 FA	STA POINTL
741	1008	20 91 1F	JSR CHK
742			;
743	100B	8A	TXA IF CNT=0 DONT
744	100C	F0 0F	BEQ LOAD3 GET ANY DATA
745			;
746	100E	20 9D 1F	LOAD2 JSR GETBYT GET DATA
747	1011	91 FA	STA (POINTL),Y STORE DATA
748	1013	20 91 1F	JSR CHK
749	1016	20 63 1F	JSR INCPY NEXT ADDRESS
750	1019	CA	DEX
751	101A	D0 F2	BNE LOAD2
752	101C	E8	INX
753			;
754	101D	20 9D 1F	LOAD3 JSR GETBYT
755	1020	C5 F6	CMP CHKHI
756	1022	D0 17	BNE LOADE1
757	1024	20 9D 1F	JSR GETBYT
758	1027	C5 F7	CMP CHKSUM
759	1029	D0 13	BNE LOADER
760			;
761	102B	8A	TXA
762	102C	D0 B9	BNE LOAD
763			;
764	102E	A2 0C	LOAD7 LDX #\$0C X-OFF KIM
765	1030	A9 27	LOAD8 LDA #\$27
766	1032	8D 42 17	STA SBD
767	1035	20 31 1E	JSR PRST
768	1038	4C 4F 1C	JMP START
769			;
770	103B	20 9D 1F	LOADE1 JSR GETBYT DUMMY
771	103E	A2 11	LOADER LDX #\$11 X-OFF ERR KIM
772	1040	D0 EE	BNE LOAD8
773			;
774			;
775			;
776			;

CARD #	LOC	CODE	CARD			
777						
778	1D42	A9 00	DUMP	LDA	#\$00	
779	1D44	85 F8		STA	INH	
780	1D46	85 F9		STA	INH	CLEAR RECORD COUNT
781	1D48	A9 00	DUMP0	LDA	#\$00	
782	1D4A	85 F6		STA	CHKHI	CLEAR CHKSUM
783	1D4C	85 F7		STA	CHKSUM	
784						
785	1D4E	20 2F 1E	DUMP1	JSR	CRLF	PRINT CR LF
786	1D51	A9 3B		LDA	#\$3B	PRINT SMICOLON
787	1D53	20 A0 1E		JSR	OUTCH	
788	1D56	A5 FA		LDA	POINTL	TEST POINT GT OR ET
789	1D58	CD F7 17		CMR	EAL	HI LIMIT GO TO EXIT
790	1D5B	A5 FB		LDA	POINTH	
791	1D5D	ED F8 17		SBC	EAH	
792	1D60	90 18		BCC	DUMP4	
793						
794	1D62	A9 00		LDA	#\$00	PRINT LAST RECORD
795	1D64	20 3B 1E		JSR	PRTRYT	0 BYTES
796	1D67	20 CC 1F		JSR	OPEN	
797	1D6A	20 1E 1E		JSR	PRTPNT	
798						
799	1D6D	A5 F6		LDA	CHKHI	PRINT CHKSUM
800	1D6F	20 3B 1E		JSR	PRTRYT	FOR LAST RECORD
801	1D72	A5 F7		LDA	CHKSUM	
802	1D74	20 3B 1E		JSR	PRTRYT	
803	1D77	4C 64 1C		JMP	CLEAR	
804						
805	1D7A	A9 13	DUMP4	LDA	#\$13	PRINT 24 BYTE CNT
806	1D7C	A9		TAX		SAVE AS INDEX
807	1D7D	20 3B 1E		JSR	PRTRYT	
808	1D80	20 91 1F		JSR	CHK	
809	1D83	20 1E 1E		JSR	PRTPNT	
810						
811	1D86	A0 00	DUMP2	LDY	#\$00	PRINT 24 BYTES
812	1D88	B1 FA		LDA	(POINTL),Y	GET DATA
813	1D8A	20 3B 1E		JSR	PRTRYT	PRINT DATA
814	1D8D	20 91 1F		JSR	CHK	COMP CHKSUM
815	1D90	20 63 1F		JSR	INCP	INCREMENT POINT
816	1D93	CA		DEX		
817	1D94	D0 F0		BNE	DUMP2	
818						
819	1D96	A5 F6		LDA	CHKHI	PRINT CHKSUM
820	1D98	20 3B 1E		JSR	PRTRYT	
821	1D9B	A5 F7		LDA	CHKSUM	
822	1D9D	20 3B 1E		JSR	PRTRYT	
823	1DA0	E6 F8		INC	INH	INCREMENT RECORD CNT
824	1DA2	D0 02		BNE	DUMP3	
825	1DA4	E6 F9		INC	INH	
826	1DA6	4C 48 1D	DUMP3	JMP	DUMP0	
827						
828	1DA9	20 CC 1F	SPACE	JSR	OPEN	OPEN NEW CELL
829	1DAC	20 2F 1E	SHOW	JSR	CRLF	PRINT CR LF
830	1DAF	20 1E 1E	SHOW1	JSR	PRTPNT	
831	1DB2	20 9E 1E		JSR	OUTSP	PRT SPACE
832	1DB5	A0 00		LDY	#\$00	PRINT DATA SPECIFIED
833	1DB7	B1 FA		LDA	(POINTL),Y	BY POINT AD = LDA EXT
834	1DB9	20 3B 1E		JSR	PRTRYT	
835	1DBC	20 9E 1E		JSR	OUTSP	PRT SPACE
836	1DBF	4C 64 1C		JMP	CLEAR	
837						
838	1DC2	20 63 1F	RTRN	JSR	INCP	OPEN NEXT CELL
839	1DC5	4C AC 1D		JMP	SHOW	
840						
841	1DC8	A6 F2	GOEXEC	LTX	SPUSER	

CARD #	LOC	CODE	CARD		
842	1DC8	3A		TXS	
843	1DCB	A5 FB		LDA	POINTH
844	1DCD	48		PHA	PROGRAM RUNS FROM
845	1DCE	A5 FA		LDA	OPEN CELL ADDRESS
846	1DD0	48		PHA	
847	1DD1	A5 F1		LDA	PREG
848	1DD3	48		PHA	
849	1DD4	A6 F4		LDX	XREG
850	1DD6	A4 F5		LDY	YREG
851	1DD8	A5 F3		LDA	ACC
852	1DDA	40		RTI	
853					
854	1DD8	C9 20	;	SCAN	CMP
855	1DD0	F0 0A			BEQ
856	1DDF	C9 7F			CMP
857	1DE1	F0 18			BEQ
858	1DE3	C9 0D			CMP
859	1DE5	F0 DB			BEQ
860	1DE7	C9 0A			CMP
861	1DE9	F0 1C			BEQ
862	1DEB	C9 2E			CMP
863	1DED	F0 26			BEQ
864	1DEF	C9 47			CMP
865	1DF1	F0 D5			BEQ
866	1DF3	C9 51			CMP
867	1DF5	F0 0A			BEQ
868	1DF7	C9 4C			CMP
869	1DF9	F0 09			BEQ
870	1DFB	4C 6A 1C			JMP
871					
872	1DFE	4C 4F 1C	;	STV	JMP
873	1E01	4C 42 1D		DUMPV	JMP
874	1E04	4C E7 1C		LOADV	JMP
875					
876	1E07	38		FEED	SEC
877	1E08	A5 FA			LDA
878	1E0A	E9 01			SBC
879	1E0C	95 FA			STA
880	1E0E	B0 02			BCS
881	1E10	C6 FB			DEC
882	1E12	4C AC 1D		FEED1	JMP
883					
884	1E15	A0 00		MODIFY	LDY
885	1E17	A5 F8			LDA
886	1E19	91 FA			STA
887	1E1B	4C C2 1D			JMP
888					
889					
890					
891					
892					
893					
894					
895					
896					
897	1E1E	A5 FB		PRINT	LDA
898	1E20	20 3B 1E			JSR
899	1E23	20 91 1F			JSR
900	1E26	A5 FA			LDA
901	1E28	20 3B 1E			JSR
902	1E2B	20 91 1F			JSR
903	1E2E	60			RTS
904					
905					
906					
907					

CARD #	LOC	CODE	CARD			
908	1E2F	A2 07	ORLF	LDX	#007	
909	1E31	8D 05 1F	PRTST	LDA	TOP,X	
910	1E34	20 A0 1E		JSR	OUTCH	
911	1E37	0A		DEX		
912	1E38	10 F7		BPL	PRTST	STOP ON INDEX ZERO
913	1E3A	60	PRT1	RTS		
914			:			
915			:		PRINT 1 HEX BYTE AS TWO ASCII CHAR'S	
916			:			
917	1E38	85 FC	PRTBYT	STA	TEMP	
918	1E3D	4A		LSR	A	SHIFT CHAR RIGHT 4 BITS
919	1E3E	4A		LSR	A	
920	1E3F	4A		LSR	A	
921	1E40	4A		LSR	A	
922	1E41	20 40 1E		JSR	HEXTA	CONVERT TO HEX AND PRINT
923	1E44	A5 FC		LDA	TEMP	GET OTHER HALF
924	1E46	20 40 1E		JSR	HEXTA	CONVERT TO HEX AND PRINT
925	1E49	A5 FC		LDA	TEMP	RESTORE BYTE IN A AND RETURN
926	1E4B	60		RTS		
927			:			
928	1E4C	29 0F	HEXTA	AND	#00F	MASK HI 4 BITS
929	1E4E	09 0A		CMR	#00A	
930	1E50	18		CLC		
931	1E51	30 02		BMI	HEXTA1	
932	1E52	69 07		ADC	#007	ALPHA HEX
933	1E55	69 30	HEXTA1	ADC	#030	DEC HEX
934	1E57	40 A0 1E		JMP	OUTCH	PRINT CHAR
935			:			
936			:		GET 1 CHAR FROM ITT	
937			:		RETURN FROM SUB WITH CHAR IN A	
938			:		X IS PRESERVED AND Y RETURNED = FF	
939			:			
940	1E5A	86 FD	GETCH	STX	TMPL	SAVE X REG
941	1E5C	A2 08		LDX	#008	SET UP 8 BIT CNT
942	1E5E	99 01		LDA	#001	
943	1E60	20 40 17	GET1	BIT	SAD	
944	1E63	D0 22		BNE	GET6	
945	1E65	30 F9		BMI	GET1	WAIT FOR START BIT
946	1E67	20 D4 1E		JSR	DELAY	DELAY 1 BIT
947	1E6A	20 EB 1E	GET5	JSR	DEHALF	DELAY 1/2 BIT TIME
948	1E6D	A0 40 17	GET2	LDA	SAD	GET 8 BITS
949	1E70	29 80		AND	#080	MASK OFF LOW ORDER BITS
950	1E72	46 FE		LSR	CHAR	SHIFT RIGHT CHARACTER
951	1E74	05 FE		ORA	CHAR	
952	1E76	85 FE		STA	CHAR	
953	1E78	20 D4 1E		JSR	DELAY	DELAY 1 BIT TIME
954	1E7B	0A		DEX		
955	1E7D	D0 EF		BNE	GET2	GET NEXT CHAR
956	1E7E	20 EB 1E		JSR	DEHALF	EXIT THIS RTN
957			:			
958	1E81	A6 FD		LDX	TMPL	
959	1E83	A5 FE		LDA	CHAR	
960	1E85	2A		ROL	A	SHIFT OFF PARITY
961	1E86	4A		LSR	A	
962	1E87	60	GET6	RTS		
963			:			
964			:		INITIALIZATION FOR SIGMA	
965			:			
966	1E88	A2 01	INITS	LDX	#001	SET KB MODE TO ADDR
967	1E8A	86 FF		STX	MODE	
968			:			
969	1E8C	A2 00	INIT1	LDX	#000	
970	1E8E	8E 41 17		STX	PADD	FOR SIGMA USE SADD
971	1E91	A2 3F		LDX	#03F	
972	1E93	8E 43 17		STX	PBDD	FOR SIGMA USE SBDD

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CARD # LOC      CODE      CARD
973 1E96 A2 07      LDX  #$07      ENABLE DATA IN
974 1E98 8E 42 17    STX  SBD      OUTPUT
975 1E9B 08        CLD
976 1E9C 78        SEI
977 1E9D 60        RTS
978
979 ;
980 ; PRINT 1 CHAR CHAR=A
981 ; X IS PRESERVED Y RETURNED = FF
982 ; OUTSP PRINTS 1 SPACE
983
983 1E9E A9 20      OUTSP LDA  #$20
984 1EA0 85 FE      OUTCH STA  CHAR
985 1EA2 86 FD      STX  TMPX
986 1EA4 20 D4 1E    JSR  DELAY      10/11 BIT CODE SYNC
987 1EA7 AD 42 17    LDA  SBD      START BIT
988 1EAA 29 FE      AND  #$FE
989 1EAC 9D 42 17    STA  SBD
990 1EAF 20 D4 1E    JSR  DELAY
991 1EB2 A2 08      LDX  #$08
992 1EB4 AD 42 17    OUT1 LDA  SBD      DATA BIT
993 1EB7 29 FE      AND  #$FE
994 1EB9 46 FE      LSR  CHAR
995 1EBB 69 00      ROR  #$00
996 1EBD 8D 42 17    STA  SBD
997 1EC0 20 D4 1E    JSR  DELAY
998 1EC3 CA        DEX
999 1EC4 D0 EE      BNE  OUT1
1000 1EC6 AD 42 17   LDA  SBD      STOP BIT
1001 1EC9 09 01     ORA  #$01
1002 1ECB 8D 42 17   STA  SBD
1003 1ECE 20 D4 1E   JSR  DELAY      STOP BIT
1004 1ED1 A6 FD      LDX  TMPX      RESTORE INDEX
1005 1ED3 60        RTS
1006
1007 ;
1008 ; DELAY 1 BIT TIME
1009 ; AS DETERMEND BY DETOPS
1010
1010 1ED4 AD F3 17   DELAY LDA  CNTH30      THIS LOOP SIMULATES THE
1011 1ED7 8D F4 17   STA  TIMH      DETOPS SECTION AND WILL DELAY
1012 1EDA AD F2 17   LDA  CNTH30      1 BIT TIME
1013 1EDD 38        DE2  SEC
1014 1EDE E9 01     DE4  SBC  #$01
1015 1EE0 80 03     DE4  BCS  DE3
1016 1EE2 CE F4 17   DE3  DEC  TIMH
1017 1EE5 AC F4 17   DE3  LDY  TIMH
1018 1EE8 10 F3     DE3  BPL  DE2
1019 1EEA 60        RTS
1020
1021 ;
1022 ; DELAY HALF BIT TIME
1023 ; DOUBLE RIGHT SHIFT OF DELAY
1024 ; CONSTANT FOR A DIV BY 2
1025
1022 1EEB AD F3 17   DEHALF LDA CNTH30
1023 1EEE 8D F4 17   STA  TIMH
1024 1EF1 AD F2 17   LDA  CNTH30
1025 1EF4 4A        LSR  A
1026 1EF5 4E F4 17   LSR  TIMH
1027 1EF8 90 E3     BCC  DE2
1028 1EFA 09 30     ORA  #$80
1029 1EFC B0 E0     BCS  DE4
1030
1031 ;
1032 ; SUB TO DETERMINE IF KEY IS
1033 ; DEPRESSED OR COMBIDION OF SSM
1034 ; KEY NOT DEP OR TTY MODE      A = 0
1035 ; KEY DEP OR KB MODE          A NOT ZERO
1036 ;
1037 1EFE A0 03      AK  LDY  #$03      3 ROWS

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CARD #	LOC	CODE	CARD			
1038	1F00	A2 01		LDX	#101	DIGIT 0
1039						
1040	1F02	A9 FF	ONEKEY	LDA	#1FF	
1041	1F04	8E 42 17	AK1	STX	SBD	OUTPUT DIGIT
1042	1F07	E8		INX		GET NXT DIGIT
1043	1F08	E8		INX		
1044	1F09	2D 40 17		AND	SAD	INPUT SEGMENTS
1045	1F0C	88		DEY		
1046	1F0D	D0 F5		BNE	AK1	
1047						
1048	1F0F	A0 07		LDY	#107	
1049	1F11	8C 42 17		STY	SBD	
1050						
1051	1F14	09 80		DRA	#180	
1052	1F16	49 FF		EDR	#1FF	
1053	1F18	60		RTS		
1054						
1055				SUB		OUTPUT TO 7 SEGMENT DISPLAY
1056						
1057	1F19	A0 00	SCAND	LDY	#100	GET DATA SPECIFIED
1058	1F1B	B1 FA		LDA	(POINTL),Y	BY POINT
1059	1F1D	85 F9		STA	INH	SET UP DISPLAY BUFFER
1060	1F1F	A9 7F	SCANDS	LDA	#17F	CHANGE SEG
1061	1F21	8D 41 17		STA	PADD	TO OUTPUT
1062						
1063	1F24	A2 09		LDX	#109	INIT DIGIT NUMBER
1064	1F26	A0 03		LDY	#103	OUTPUT 3 BYTES
1065						
1066	1F28	B9 F8 00	SCAND1	LDA	INL,Y	GET BYTE
1067	1F2B	4A		LSR	A	GET MSD
1068	1F2C	4A		LSR	A	
1069	1F2D	4A		LSR	A	
1070	1F2E	4A		LSR	A	
1071	1F2F	2D 48 1F		JSR	CONVD	OUTPUT CHAR
1072	1F32	B9 F8 00		LDA	INL,Y	GET BYTE AGAIN
1073	1F35	29 0F		AND	#10F	GET LSD
1074	1F37	2D 48 1F		JSR	CONVD	OUTPUT CHAR
1075	1F3A	88		DEY		SET UP FOR NXT BYTE
1076	1F3B	D0 EB		BNE	SCAND1	
1077	1F3D	8E 42 17		STX	SBD	ALL DIGITS OFF
1078	1F40	A9 00		LDA	#100	CHANGE SEG
1079	1F42	8D 41 17		STA	PADD	TO INPUTS
1080	1F45	4C FE 1E		JMP	AK	GET ANY KEY
1081						
1082						
1083						
1084						
1085	1F48	84 FC	CONVD	STY	TEMP	SAVE Y
1086	1F4A	A8		TAY		USE CHAR AS INDEX
1087	1F4B	B9 E7 1F		LDA	TABLE,Y	LOOK UP CONVERSION
1088	1F4E	A0 00		LDY	#100	TURN OFF SEGMENTS
1089	1F50	9C 40 17		STY	SAD	
1090	1F53	8E 42 17		STX	SBD	OUTPUT DIGIT ENABLE
1091	1F56	8D 40 17		STA	SAD	OUT PUT SEGMENTS
1092						
1093	1F59	A0 7F		LDY	#17F	DELAY 500 CYCLES APPROX.
1094	1F5B	88	CONVD1	DEY		
1095	1F5C	D0 FD		BNE	CONVD1	
1096						
1097	1F5E	E8		INX		GET NEXT DIGIT NUM
1098	1F5F	E8		INX		ADD 2
1099	1F60	A4 FC		LDY	TEMP	RESTORE Y
1100	1F62	60		RTS		
1101						
1102						
				SUB		TO INCREMENT POINT.



```

CARD # LOC      CODE      CARD
1103          ;
1104 1F63 E6 FA  INCPT INC  POINTL
1105 1F65 D0 02          BNE  INCPT2
1106 1F67 E6 FB          INC  POINTH
1107 1F69 60          INCPT2 RTS
1108          ;
1109          ; GET KEY FROM KEY BOARD
1110          ; RETURN WITH A=KEY VALUE
1111          ; A GT. 15 THEN ILLEGAL OR NO KEY
1112          ;
1113          ;
1114 1F6A A2 21  GETKEY LDX  #$21      START AT DIGIT 0
1115 1F6C A0 01  GETKEY LDY  #$01      GET 1 ROW
1116 1F6E 20 02 1F JSR  ONEKEY
1117 1F71 D0 07          BNE  KEYIN    A=0 NO KEY
1118 1F73 E0 27          CPX  #$27      TEST FOR DIGT 2
1119 1F75 D0 F5          BNE  GETKEY5
1120 1F77 A9 15          LDA  #$15      15=NO KEY
1121 1F79 60          RTS
1122 1F7A A0 FF  KEYIN LDY  #$FF
1123 1F7C 0A      KEYIN1 ASL  A          SHIFT LEFT
1124 1F7D B0 03          BCS  KEYIN2    UNTIL Y=KEY NUM
1125 1F7F 08          INY
1126 1F80 10 FA      BPL  KEYIN1
1127 1F82 3A      KEYIN2 TXA
1128 1F83 29 0F          AND  #$0F      MASK MSD
1129 1F85 4A          LSR  A          DIV BY 2
1130 1F86 AA          TAX
1131 1F87 98          TYA
1132 1F88 10 03      BPL  KEYIN4
1133 1F8A 18      KEYIN3 CLC
1134 1F8B 69 07          ADC  #$07      MULT (X-1) TIMES A
1135 1F8D 0A      KEYIN4 DEX
1136 1F8E D0 FA      BNE  KEYIN3
1137 1F90 60          RTS
1138          ;
1139          ; SUB TO COMPUTE CHECK SUM
1140          ;
1141 1F91 18      CHK  CLC
1142 1F92 65 F7          ADC  CHKSUM
1143 1F94 85 F7          STA  CHKSUM
1144 1F96 A5 F6          LDA  CHKHI
1145 1F98 69 00          ADC  #$00
1146 1F9A 85 F6          STA  CHKHI
1147 1F9C 60          RTS
1148          ;
1149          ; GET 2 HEX CHAR'S AND PACK
1150          ; INTO INL AND INH
1151          ; X PRESERVED Y RETURNED = 0
1152          ; NON HEX CHAR WILL BE LOADED AS NEAREST HEX EQU
1153          ;
1154 1F9D 20 5A 1E  GETEYT JSR  GETCH
1155 1FA0 20 AC 1F      JSR  PACK
1156 1FA3 20 5A 1E      JSR  GETCH
1157 1FA6 20 AC 1F      JSR  PACK
1158 1FA9 A5 F8          LDA  INL
1159 1FAB 60          RTS
1160          ;
1161          ; SHIFT CHAR IN A INTO
1162          ; INL AND INH
1163          ;
1164 1FAC 09 30      PACK CMP  #130    CHECK FOR HEX
1165 1FAE 30 18      BMI  UPDAT2
1166 1FB0 09 47      CMP  #147    NOT HEX EXIT
1167 1FB2 10 17      BPL  UPDAT2

```

```

CARD # LOC CODE CARD
1168 1FB4 09 40 HEXNUM CMP #140 CONVERT TO HEX
1169 1FB6 30 03 BMI UPDATE
1170 1FB8 18 HEXALP CLC
1171 1FB9 69 09 ADD #109
1172 1FBB 2A UPDATE ROL A
1173 1FBC 2A ROL A
1174 1FBD 2A ROL A
1175 1FBE 2A ROL A
1176 1FBF A0 04 LDY #104 SHIFT INTO I/O BUFFER
1177 1FC1 2A UPDAT1 ROL A
1178 1FC2 26 F8 ROL INL
1179 1FC4 26 F9 ROL INH
1180 1FC6 88 DEY
1181 1FC7 D0 F8 BNE UPDAT1
1182 1FC9 A9 00 LDA #100 A=0 IF HEX NUM
1183 1FCB 60 UPDAT2 RTS
1184 ;
1185 1FC0 A5 F8 OPEN LDA INL MOVE I/O BUFFER TO POINT
1186 1FCE 85 FA STA POINTL
1187 1FD0 A5 F9 LDA INH TRANSFER INH- POINTH
1188 1FD2 85 FB STA POINTH
1189 1FD4 60 RTS
1190 ;
1191 ;
1192 ; END OF SUBROUTINES
1193 ;
1194 ;
1195 ; TABLES
1196 ;
1197 1FD5 00 TOP .BYTE $00,$00,$00,$00,$00,$00,$0A,$0D,'MIK'
1197 1FD6 00
1197 1FD7 00
1197 1FD8 00
1197 1FD9 00
1197 1FDA 00
1197 1FDB 0A
1197 1FDC 0D
1197 1FDD 4D 49 4B .BYTE /', $13, 'RRE', /', $13
1198 1FE0 20
1198 1FE1 13
1198 1FE2 52 52 45
1198 1FE5 20
1198 1FE6 13
1199 ;
1200 ; TABLE HEX TO 7 SEGMENT
1201 ; 0 1 2 3 4 5 6 7
1202 1FE7 BF TABLE .BYTE $BF,$86,$DB,$CF,$E6,$ED,$FD,$87
1202 1FE8 86
1202 1FE9 DB
1202 1FEA CF
1202 1FEB E6
1202 1FEC ED
1202 1FED FD
1202 1FEE 87
1203 ;
1204 1FEF FF .BYTE $FF,$EF,$F7,$FC,$B9,$DE,$F9,$F1
1204 1FF0 EF
1204 1FF1 F7
1204 1FF2 FC
1204 1FF3 B9
1204 1FF4 DE
1204 1FF5 F9
1204 1FF6 F1
1206 ;
1207 ;
1208 ;

```

```

CARD # LOC      CODE      CARD
1209          ;
1210          ;          INTERRUPT VECTORS
1211          ;
1212 1FF7          ◆=$1FFA
1213 1FFA 10 10    NMIENT .WORD NMIT
1214 1FFC 22 10    RSTENT .WORD RST
1215 1FFE 1F 10    IPQENT .WORD IPQT
1216          .END

```

```

END OF MOS/TECHNOLOGY 650X ASSEMBLY VERSION 4
NUMBER OF ERRORS = 0, NUMBER OF WARNINGS = 0

```

## SYMBOL TABLE

SYMBOL	VALUE	LINE	DEFINED	CROSS-REFERENCES
ACC	00F3	76	587	851
ADDR	10BE	694	687	
ADDPM	10C8	701	673	
AK	1EFE	1037	1080	
AK1	1F04	1041	1046	
CHAR	00FE	90	950	951 952 959 984 994
CHK	1F91	1141	734 738	741 748 808 914 899 902
CHKH	17E8	97	158 265	289 299 301
CHKHI	00F6	82	730 755	782 799 819 1144 1146
CHKL	17E7	96	156 262	288 297 298
CHKSUM	00F7	83	729 758	783 801 821 1142 1143
CHKT	194C	295	234 237	242 244 256 308
CHT1	1982	336	344	
CHT2	198E	341	338	
CHT3	1991	342	340	
CLEAR	1064	645	803 836	
CLKKT	1747	65	◆◆◆◆	
CLKRDI	1747	66	355 361	378 384 498 505
CLKRDT	1746	67	459 471	
CLK1T	1744	62	358 364	381 387 501 508
CLK64T	1746	64	461 473	
CLK8T	1745	63	◆◆◆◆	
CNTH30	17F3	101	613 622	1010 1022
CNTH30	17F2	100	625 1012	1024
CONVD	1F48	1085	1071 1074	
CONVD1	1F5B	1094	1095	
CRLF	1E2F	908	640 785	829
DATA	1CA8	680	◆◆◆◆	
DATAM	10CC	704	675	
DATAM1	10CE	705	702	
DATAM2	1CD0	706	699	
DATA1	1CB0	686	698	
DATA2	10C3	697	692	
DEHALF	1EEB	1022	947 956	
DELAY	1ED4	1010	946 953	986 990 997 1003
DETCPS	1C2A	612	◆◆◆◆	
DET1	1C31	615	617	
DET2	1C42	623	621	
DET3	1C3A	619	624	
DE2	1EDD	1013	1018 1027	
DE3	1EE5	1017	1015	
DE4	1EDE	1014	1029	
DUMP	1D42	778	873	
DUMPT	1800	121	◆◆◆◆	
DUMPT1	1814	131	134	
DUMPT2	1833	148	177	
DUMPT3	1854	163	166	

SYMBOL	VALUE	LINE	DEFINED	CROSS-REFERENCES
DUMPT4	1865	173	152	
DUMPV	1E01	873	867	
DUMP0	1D48	781	826	
DUMP1	1D4E	785	♦♦♦♦	
DUMP2	1D86	811	817	
DUMP3	1D46	826	824	
DUMP4	1D7A	805	792	
EAH	17F8	106	151	791
EAL	17F7	105	149	789
FEED	1E07	876	861	
FEED1	1E12	982	880	
GETBYT	1F9D	1154	732	736 739 746 754 757 770
GETCH	1E5A	940	648	725 1154 1156
GETK	1C8D	667	♦♦♦♦	
GETKEY	1F6A	1114	667	
GETKES	1F6C	1115	1119	
GET1	1E60	943	945	
GET2	1E6D	948	955	
GET5	1E6A	947	627	
GET6	1E87	962	944	
GOEXEC	1DC8	841	711	865
GOV	1CD9	711	679	
HEXALP	1FB8	1170	♦♦♦♦	
HEXNUM	1FB4	1163	♦♦♦♦	
HEXOUT	196F	323	314	316
HEXTA	1E4C	928	922	924
HEXTA1	1E55	933	931	
HEX1	1979	328	326	
ID	17F9	107	140	224 226
INCPT	1F63	1104	708	749 815 938
INCPT2	1F69	1107	1105	
INCVEB	19EA	397	176	258
INCVE1	19F2	400	398	
INH	00F9	85	647	780 825 1059 1179 1187
INIT3	1E88	966	600	609
INIT1	1E8C	969	636	
INL	00F8	84	646	779 823 885 1066 1072 1158 1178 1185
INTVEB	1932	281	123	185
IRGENT	1FFE	1215	♦♦♦♦	
IRGP27	1BFE	519	♦♦♦♦	
IRPT	1C1F	604	1215	
IRQV	17FE	113	604	
KEYIN	1F7A	1122	1117	
KEYIN1	1F7C	1123	1126	
KEYIN2	1F82	1127	1124	
KEYIN3	1F8A	1133	1136	
KEYIN4	1F8D	1135	1132	
LOAD	1CE7	725	727	762 874
LOADER	1D3E	771	759	
LOADE1	1D3B	770	756	
LOADS	1CEE	728	♦♦♦♦	
LOADT	1873	183	231	
LOADT4	18B5	216	221	
LOADT5	18D7	233	225	228
LOADT6	18EC	241	230	
LOADT7	18F8	247	239	259
LOADT8	1915	261	250	
LOADT9	1929	270	252	263 266
LOADV	1E04	874	869	
LOAD10	192B	271	268	
LOAD11	18C2	223	218	
LOAD12	190F	258	182	
LOAD13	18FA	248	254	
LOAD2	1D0E	746	751	

SYMBOL	VALUE	LINE	DEFINED	CROSS-REFERENCES															
LOAD3	1D1D	754	744																
LOAD7	1D2E	764	♦♦♦♦																
LOAD8	1D30	765	772																
MODE	00FF	91	686	705	967														
MODIFY	1E15	884	863																
NMIENT	1FFA	1213	♦♦♦♦																
NMIP27	1BFA	517	♦♦♦♦																
NMIT	1C1C	603	1213																
NMIV	17FA	111	603																
ONE	199E	353	336	339															
ONEKEY	1F02	1040	1116																
ONE1	19A1	355	356	368															
ONE2	19B0	361	362																
OPEN	1F0C	1185	796	828															
OUTBT	1961	309	141	157	159														
OUTBTC	195E	308	144	146	174														
OUTCH	1EA0	984	787	910	934														
OUTCHT	197A	333	132	138	155	164													
OUTSP	1E9E	983	831	835															
OUT1	1EB4	992	999																
PACK	1FAC	1164	651	1155	1157														
PACKT	1A00	413	251	405	407														
PACKT1	1A0F	421	418																
PACKT2	1A15	426	429																
PACKT3	1A22	433	414	416															
PADD	1741	59	970	1061	1079														
PBDD	1743	61	128	496	972														
PCCMD	1C0C	717	671																
PCH	00F0	73	594	719															
PCL	00EF	72	591	717															
PLLCAL	1A6B	493	517	518	519														
PLL1	1A75	498	499	511															
PLL2	1A84	505	506																
POINTH	00FB	87	170	272	595	696	720	737	790	843	881	897							
			1106	1188															
POINTL	00FA	86	169	271	592	688	691	695	718	740	747	788							
			812	833	845	877	879	886	900	1058	1104	1186							
PREG	00F1	74	589	847															
PRTBYT	1E3B	917	795	800	802	807	813	820	822	834	898	901							
PRTPNT	1E1E	897	797	809	830														
PRTST	1E31	909	642	767	912														
PRT1	1E3A	913	♦♦♦♦																
RDBIT	1A41	457	200	441	458														
RDBIT2	1A53	467	468																
RDBIT3	1A50	464	465																
RDBIT4	1A63	476	477																
RDBYT	19F3	404	223	233	236	241	243	261	264										
RDBYT2	19F9	406	♦♦♦♦																
RDCHT	1A24	439	209	216	248	404	406												
RDCHT1	1A29	441	446																
READ	1C6A	648	870																
RST	1C22	606	1214																
RSTENT	1FFC	1214	♦♦♦♦																
RSTP27	1BFC	518	♦♦♦♦																
RSTV	17FC	112	♦♦♦♦																
RTRN	1D02	838	859	887															
SAD	1740	58	615	623	638	660	943	948	1044	1089	1091								
SAH	17F6	104	145	283															
SAL	17F5	103	143	281															
SAVE	1C00	587	♦♦♦♦																
SAVE1	1C05	590	♦♦♦♦																
SAVE2	1C0F	596	♦♦♦♦																
SAVX	17E9	98	198	201	202	203	204	233	334	345	346	427							
			430	439	442	443	444	448	451										

SYMBOL	VALUE	LINE	DEFINED	CROSS-REFERENCES																	
SBD	1742	60	126	195	360	366	383	389	457	467	494	503									
			510	766	974	987	989	992	996	1000	1002	1041									
			1049	1077	1090																
SCAN	1DD8	854	652																		
SCAND	1F19	1057	657	662	664																
SCANDS	1F1F	1060	♦♦♦♦																		
SCAND1	1F28	1066	1076																		
SHOW	1DAC	829	839	882																	
SHOW1	1DAF	830	643																		
SPACE	1D39	828	855																		
SPUSER	00F2	75	599	608	841																
START	1C4F	636	171	273	601	616	658	661	669	706	709	721									
			768	872																	
STEP	1CD3	708	677																		
STV	1DFE	872	857																		
SYNC	1891	197	211	220																	
SYNC1	1896	200	206																		
SYNC2	18AB	209	213																		
TAB	1871	182	189	191																	
TABLE	1FE7	1202	1087																		
TEMP	00FC	88	684	689	917	923	925	1085	1099												
TIMH	17F4	102	1011	1016	1017	1023	1026														
TMFX	00FD	89	940	958	985	1004															
TOP	1FD5	1197	909																		
TTYKB	1C77	657	639	650																	
TTYKB1	1C7C	659	663	665																	
UPDATE	1F8B	1172	1169																		
UPDAT1	1FC1	1177	1181																		
UPDAT2	1FCB	1183	1165	1167																	
VEB	17EC	99	122	148	150	173	184	188	190	192	235	238									
			257	282	284	286	397	399													
XREG	00F4	77	597	849																	
YREG	00F5	78	596	850																	
ZPD	19C4	376	341	342																	
ZPD1	19C7	378	379	391																	
ZPD2	19D6	384	385																		

## INSTRUCTION COUNT

ADC	13
AND	9
ASL	7
BCC	4
BOS	5
BEO	26
BIT	12
BMI	9
BNE	44
BPL	15
BRK	0
BVC	0
BVS	0
CLC	8
CLD	1
CLI	0
CLV	0
CMP	38
CPX	1
CPY	0
DEC	2
DEX	14
DEY	8
EOR	2
INC	7
INX	5
INY	2
JMP	31
JSR	115
LDA	108
LDX	29
LDY	25
LSR	22
NOP	0
ORA	6
PHA	5
PHP	0
PLA	5
PLP	0
ROL	18
RTI	1
RTS	28
SBC	5
SEC	3
SED	0
SEI	1
STA	81
STX	14
STY	7
TAX	3
TAY	3
TSX	1
TXA	3
TXS	2
TYA	4

```
# SYMBOLS = 204 (LIMIT = 400)      # BYTES = 1690 (LIMIT = 4096)
# LINES = 1242 (LIMIT = 1500)      # XREFS = 646 (LIMIT = 900)
STOP 0
```